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## SEVENTH WINTER SCHOOL (1979)

## MODEL THEORETIC APPROACH TO TOPOLOGICAL FUNCTORS, II.

by

Jiří Rosický

This paper is a sequel of [6]. Most of results here presented will appear in the forthcoming author's paper [7].

Under a concrete category  $(\mathcal{A}, U)$  we will mean a category  $\mathcal{A}$  equipped with a faithful functor  $U: \mathcal{A} \rightarrow \text{Set}$  satisfying the following two conditions:

- (1) If  $A \in \mathcal{A}$ ,  $X$  is a set and  $f: UA \rightarrow X$  a bijection, then there is  $B \in \mathcal{A}$  and an isomorphism  $g: A \rightarrow B$  such that  $Ug = f$
- (2) If  $A \in \mathcal{A}$  and  $f: A \rightarrow A$  is an isomorphism such that  $Uf$  is the identity, then  $f$  is the identity.

Under a functor  $F: (\mathcal{A}, U) \rightarrow (\mathcal{B}, V)$  between concrete categories we will mean a functor  $F: \mathcal{A} \rightarrow \mathcal{B}$  such that  $V.F = U$ .

A type is given by a class of function symbols and a class of relation symbols. Their arities are arbitrary cardinals. The infinitary first-order language  $L_{\infty, \omega}(\tau)$  of type  $\tau$  includes a proper class  $V$  of variables and besides the usual logical symbols it admits infinitary conjunctions, disjunctions and quantifiers. A class of sentences of  $L_{\infty, \omega}(\tau)$  is called a theory of type  $\tau$ . We denote by  $(\mathcal{A}_\tau, U_\tau)$  or  $(\mathcal{A}_T, U_T)$  the concrete category of all  $\tau$ -structures or  $T$ -models resp. These categories need not be legitimate, i.e. they need not form a class. A theory having a representative set of  $n$ -ary atomic formulae for each cardinal  $n$  will be called normal. If  $T$  is normal, then  $(\mathcal{A}_T, U_T)$  is a legitimate category and even it is strongly fibre-small in the sense of [1].

If  $(\mathcal{A}, U)$  is a concrete category and  $n$  a cardinal, then  $U^n$  will denote the functor  $\text{Set}(n, U-)$ . Subfunctors of  $U^n$  will be called  $n$ -ary relation symbols interpretable in  $(\mathcal{A}, U)$  and natural transformations

$U^n \rightarrow U$   $n$ -ary function symbols interpretable in  $(A, U)$ . It is motivated by the fact that any relation or function symbol of type  $\uparrow$  determines a subfunctor of  $U^n$  or a natural transformation  $U^n \rightarrow U$  resp. Let  $\mathcal{F}_U$  be the collection of all relation and function symbols interpretable in  $(A, U)$ . We emphasize that  $\mathcal{F}_U$  need not be a type because it need not be a class.

Let  $\sim \in \mathcal{F}_U$  be a type. There is a functor  $G_\sim: (A, U) \rightarrow (A_\sim, U_\sim)$  such that if  $A \in \mathcal{A}$ , then the  $\sim$ -structure  $G_\sim(A)$  has the underlying set  $UA$ , the  $n$ -ary relation on  $UA$  corresponding to  $R \in \text{Rel}_n(\sim)$  equals to  $R(A)$  and the  $n$ -ary function  $f: (UA)^n \rightarrow UA$  corresponding to  $f \in \text{Fnt}_n(\sim)$  is the component  $f_A$  of the natural transformation  $f$ . Let  $T_\sim$  be the theory of type  $\sim$  consisting of all sentences which hold in all  $\sim$ -structures  $G_\sim(A)$  for  $A \in \mathcal{A}$ . Clearly we get the functor  $G_\sim: (A, U) \rightarrow (A_{T_\sim}, U_{T_\sim})$ .

We may restrict ourselves in the formation of  $T_\sim$  to some specified kind of sentences. This yields a general method of getting suitable completions or hulls of  $(A, U)$ . E.g. (with size conditions aside), if  $\sim$  consists of all function symbols from  $\mathcal{F}_U$  and  $T \in T_\sim$  of all atomic sentences, then  $T$  is the Linton's equational theory of  $U$  and  $G: (A, U) \rightarrow (A_T, U_T)$  is the equational completion of  $(A, U)$  (see [5]).

If  $A \in \mathcal{A}$ , then  $R_A(X) = \{Uf / f: A \rightarrow X\}$  defines a subfunctor  $R_A$  of  $U^{UA}$ . Let  $\mathcal{T}_U \in \mathcal{F}_U$  be the type consisting of  $R_A$  where  $A$  carries over mutually non-isomorphic objects  $A \in \mathcal{A}$  such that  $UA$  is a cardinal. Then  $G_{\mathcal{T}_U}$  is a full embedding and it is important that whenever  $(A, U)$  is strongly fibre-small, then  $T_{\mathcal{T}_U}$  is normal and  $(A, U)$  isomorphic to  $(A_{T_{\mathcal{T}_U}}, U_{T_{\mathcal{T}_U}})$ .

Further, if  $T$  consist of all universal Horn sentences without equality (their specification follows) from  $T_{\mathcal{T}_U}$ , then  $(A_T, U_T)$  is the Mac Nille completion of  $(A, U)$  (in the sense of Herrlich [3]). It

proves the conjecture from [6].

**Theorem:** A concrete category  $(\mathcal{A}, U)$  is (absolutely) topological iff it is isomorphic to the category of models of a relational normal universal Horn theory  $T$  without equality of some type  $\tau$ .

Relational means that  $\tau$  contains relation symbols only and universal Horn theory without equality consists of sentences bearing this name, i.e. arising from formulas  $\bigwedge_{i \in I} R_i(x_i) \rightarrow R(x)$ , where  $R_i \in \text{Rel}_{n_i}(\tau)$ ,  $R \in \text{Rel}_n(\tau)$ ,  $x_i \in V^{n_i}$  and  $x \in V^n$ , by universal quantification of all their variables.

Similarly, using  $\tau_U$  and a suitable kind of sentences one can treat (epi-monosource)-topological categories (in the sense of [4]) or semi-topological categories (in the sense of [8]). In either case we get a completion playing the role of Mac Neilles one in the event of topological categories.

A relational theory  $T$  of type  $\tau$  will be called reflexive if for any relation symbol  $R$  of  $\tau$   $T \models (\forall x)R(x, x, \dots, x, \dots)$  holds where  $x \in V$ . It is transitive if for any cardinal  $n$  and any  $R \in \text{Rel}_n(\tau)$   $T \models (\forall x) [(\bigwedge_{i \in n} R(x_{i,1}, x_{i,2}, \dots, x_{i,j}, \dots)) \wedge (\bigwedge_{j \in n} R(x_{1,j}, x_{2,j}, \dots, x_{i,j}, \dots))] \rightarrow R(x_{1,1}, x_{2,2}, \dots, x_{i,1}, \dots)]$  holds where  $x = (x_{i,j}) \in V^{n \times n}$ . Motivating is the case of a binary relation symbol  $R$ .

**Proposition:** Let  $T$  be a relational, normal, reflexive and transitive universal Horn theory without equality. Then  $(\mathcal{A}_T, U_T)$  is a cartesian closed topological category.

The author conjectures that this proposition can be converted. Namely, one is tempted to seek for a type  $\sim \in \tau_U$  such that  $(\mathcal{A}_{\sim}, U_{\sim})$  is (in general non-legitimate) cartesian closed topological hull of  $(\mathcal{A}, U)$  and its legitimacy corresponds to strict fibre-smallness of  $(\mathcal{A}, U)$  in the sense of Adámek and Koubek [2] (i.e. model theoretically

recover their theorem).

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