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Travel groupoids

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## TRAVEL GROUPOIDS

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*Abstract.* In this paper, by a *travel groupoid* is meant an ordered pair  $(V, *)$  such that  $V$  is a nonempty set and  $*$  is a binary operation on  $V$  satisfying the following two conditions for all  $u, v \in V$ :

$$(u * v) * u = u;$$
$$\text{if } (u * v) * v = u, \text{ then } u = v.$$

Let  $(V, *)$  be a travel groupoid. It is easy to show that if  $x, y \in V$ , then  $x * y = y$  if and only if  $y * x = x$ . We say that  $(V, *)$  is on a (finite or infinite) graph  $G$  if  $V(G) = V$  and

$$E(G) = \{\{u, v\} : u, v \in V \text{ and } u \neq u * v = v\}.$$

Clearly, every travel groupoid is on exactly one graph. In this paper, some properties of travel groupoids on graphs are studied.

*Keywords:* travel groupoid, graph, path, geodetic graph

*MSC 2000:* 20N02, 05C38, 05C12

By a graph we mean here a (finite or infinite) undirected graph with no multiple edges or loops. We will use the terminology of the book [1] but we extend it also to infinite graphs here. By a geodetic graph we mean a connected graph  $G$  such that there exists exactly one shortest  $u - v$  path in  $G$  for all  $u, v \in V(G)$ .

The letters  $h - n$  will serve for denoting non-negative integers.

## 1. TRAVEL GROUPOIDS

By a *travel groupoid* we will mean an ordered pair  $(V, *)$  such that  $V$  is a nonempty set and  $*$  is a binary operation on  $V$  satisfying the following axioms (t1) and (t2):

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- (t1)  $(u * v) * u = u$  (for all  $u, v \in V$ );  
 (t2) if  $(u * v) * v = u$ , then  $u = v$  (for all  $u, v \in V$ ).

We say that a travel groupoid  $(V, *)$  is *finite* if  $V$  is finite. If  $(V, *)$  is a travel groupoid, then we say that  $*$  is a travel operation on  $V$ . Special kinds of travel groupoids (or travel operations) were introduced in [2] and [3].

The idea of the proof of the next proposition can be found in the proof of Lemma 2 in [3].

**Proposition 1.** *If  $(V, *)$  is a travel groupoid, then  $x * x = x$  (for each  $x \in V$ ).*

**Proof.** Let  $x \in V$ . As follows from (t1),  $(x * x) * x = x$ . This implies that  $((x * x) * x) * x = x * x$ . By (t2),  $x * x = x$ , which completes the proof.  $\square$

**Proposition 2.** *Let  $(V, *)$  be a travel groupoid. Then*

- (1)  $x * y = y$  if and only if  $y * x = x$  (for all  $x, y \in V$ ),  
 (2)  $x * y = x$  if and only if  $x = y$  (for all  $x, y \in V$ )

and

- (3)  $x * (x * y) = x * y$  (for all  $x, y \in V$ ).

**Proof.** Clearly, (1) follows from (t1) and (2) follows from Proposition 1 and (t2).

Consider arbitrary  $x, y \in V$ . By (t1),  $(x * y) * x = x$ . As follows from (1),  $x * (x * y) = x * y$ . Thus (3) holds.  $\square$

Let  $(V, *)$  be a travel groupoid, and let  $G$  be a graph. We say that  $(V, *)$  is *on*  $G$  or that  $G$  *has*  $(V, *)$  if  $V(G) = V$  and

$$E(G) = \{\{u, v\} : u, v \in V \text{ and } u \neq u * v = v\}.$$

As follows from Proposition 2, every travel groupoid is on exactly one graph.

**Proposition 3.** *Let  $(V, *)$  be a travel groupoid on a graph  $G$ , let  $u, v \in V$  and  $u \neq v$ . Then  $u$  and  $u * v$  are adjacent vertices of  $G$ .*

**Proof.** The proposition follows from (1) and (3).  $\square$

Let  $(V, *)$  be a travel groupoid. If  $u, v \in V$ , then we define

$$(4) \quad u *^0 v = u$$

and

$$u *^{i+1} v = (u *^i v) * v \text{ for every } i \geq 0.$$

It is clear that if  $j, k \geq 0$ , then  $(u *^j v) *^k v = u *^{j+k} v$ .

**Proposition 4.** *Let  $(V, *)$  be a travel groupoid, let  $u, v \in V$ , and let  $k \geq 1$ . If  $u *^k v \neq v$ , then  $u *^{k-1} v \neq v$  and the elements*

$$u *^{k-1} v, \quad u *^k v, \quad \text{and} \quad u *^{k+1} v$$

*are pairwise distinct.*

*Proof.* Let  $u *^k v \neq v$ . If  $u *^{k-1} v = v$ , then Proposition 1 implies that  $u *^k v = v$ ; a contradiction. Thus  $u *^{k-1} v \neq v$ . Recall that  $u *^k v = (u *^{k-1} v) * v$ . If  $u *^k v = u *^{k-1} v$ , then, by virtue of (2),  $u *^{k-1} v = v$ ; a contradiction. Thus  $u *^k v \neq u *^{k-1} v$ . If  $u *^k v = u *^{k+1} v$ , then it follows from (2) that  $u *^k v = v$ ; a contradiction. Thus  $u *^k v \neq u *^{k+1} v$ . If  $u *^{k+1} v = u *^{k-1} v$ , then, as follows from (t2),  $u *^{k-1} v = v$ ; a contradiction. Thus  $u *^{k-1} v \neq u *^{k+1} v$ , which completes the proof.  $\square$

**Remark 1.** Let  $(V, *)$  be a travel groupoid, and let  $u, v \in V$ . If there exists  $i \geq 0$  such that  $u *^i v = v$ , then, by virtue of (2),  $u *^{i+1} v = v$ . This implies that there exists at most one  $k \geq 1$  such that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ .

The following theorem motivates the terms “travel groupoids” and “travel operations”.

**Theorem 1.** *Let  $(V, *)$  be a travel groupoid on a graph  $G$ , let  $u, v \in V$ , and let  $k \geq 1$ . Assume that  $u *^{k-1} v \neq v$ . Then the sequence*

$$(5) \quad u *^0 v, \dots, u *^{k-1} v, u *^k v$$

*is a walk in  $G$ . Moreover, if  $u *^k v = v$ , then the sequence (5) is an  $u - v$  path in  $G$ .*

*Proof.* Since  $u *^{k-1} v \neq v$ , it follows from Proposition 4 that  $u *^h v \neq v$  for each  $h$ ,  $0 \leq h \leq k - 1$ . By the definition,  $u *^{h+1} v = (u *^h v) * v$  for each  $h$ ,  $0 \leq h \leq k - 1$ . Thus, by virtue of Proposition 3, the sequence (5) is a walk in  $G$ .

Let  $u *^k v = v$ . Assume that there exist  $i$  and  $j$ ,  $1 \leq i < j \leq k$ , such that  $u *^i v = u *^j v$ . By virtue of Proposition 4,  $j < k$ . Thus  $v = u *^k v = (u *^j v) *^{k-j} v = (u *^i v) *^{k-j} v = u *^{k-(j-i)} v \neq v$ ; a contradiction. We see that the vertices  $u *^0 v, \dots, u *^{k-1} v, u *^k v$  are pairwise distinct. Hence the sequence (5) is an  $u - v$  path in  $G$ , which completes the proof.  $\square$

Let  $G$  be a geodetic graph, and let  $d$  denote the distance function of  $G$ . Put  $V = V(G)$ . It is not difficult to see that if  $u, v \in V$  and  $u \neq v$ , then there exists exactly one vertex  $A_G(u, v)$  such that

$$d(u, A_G(u, v)) = 1 \quad \text{and} \quad d(A_G(u, v), v) = d(u, v) - 1.$$

Define a binary operation  $*$  on  $V$  as follows:

$$x * y = A_G(x, y) \quad \text{if } x \neq y$$

and

$$x * y = x \quad \text{if } x = y$$

for all  $x, y \in V$ . We will say that  $(V, *)$  is the *proper* groupoid of  $G$ .

It is clear that the proper groupoid of a geodetic graph  $G$  is a travel groupoid on  $G$ . Thus every geodetic graph has at least one travel groupoid.

Obviously, every tree is a geodetic graph. (Note that the proper groupoid of a finite tree was characterized in [3]).

**Proposition 5.** *Every finite tree has exactly one travel groupoid.*

*Proof.* Consider an arbitrary finite tree  $T$ . Put  $V = V(T)$ . Let  $(V, *)$  be the proper groupoid of  $T$ . Suppose, to the contrary, that there exists a travel groupoid  $(V, \circ)$  of  $T$  such that  $(V, \circ)$  is different from  $(V, *)$ . Then there exist  $u, v \in V$  such that  $u \circ v \neq u * v$ . By Proposition 3, both vertices  $u * v$  and  $u \circ v$  are adjacent to  $u$  in  $T$ . Since  $u \circ v \neq u * v$ , we see that the vertices  $u \circ v$  and  $u * v$  belong to distinct components of  $T - u$ . Recall that  $(V, *)$  is the proper groupoid of  $T$ . This implies that the vertices  $u * v$  and  $v$  belong to the same component of  $T - u$ . Since  $T$  contains no cycle, then, by virtue of (t1) and (t2), the vertices

$$u \circ v, u \circ^2 v, u \circ^3 v, \dots$$

are pairwise distinct, which contradicts the fact that  $V$  is finite. Thus the proposition is proved.  $\square$

## 2. SIMPLE TRAVEL GROUPOIDS

We say that a travel groupoid  $(V, *)$  is *simple* if it satisfies the following axiom (t3) if  $v * u \neq u$ , then  $u * (v * u) = u * v$  (for all  $u, v \in V$ ).

Note that the travel groupoids discussed in [2] are simple.

**Remark 2.** Let  $(V, *)$  be a simple travel groupoid, and let  $u, v \in V$  such that  $v * u \neq u$ . By (1),  $u * v \neq v$ . Thus, by (t3),  $u * (v * u) = u * v$  and  $v * (u * v) = v * u$ .

The next remark gives an example of a travel groupoid which is not simple.

**Remark 3.** Let  $D$  be a directed cycle with  $|V(D)| = 2n$ , where  $n \geq 2$ . Put  $V = V(D)$ . Clearly, for every  $u \in V$  there exists exactly one vertex, say the vertex  $u'$ , such that  $(u, u')$  is a directed edge in  $D$ . Let  $C$  denote the underlying graph of  $D$ . Obviously,  $C$  is a cycle of length  $2n$ . Let  $d$  denote the distance function of  $C$ . We denote by  $*$  the binary operation on  $V$  defined as follows for all  $v, w \in V$ :

$$v * w = v \text{ if } d(v, w) = 0;$$

$$v * w = v' \text{ if } d(v, w) = n;$$

$$v * w \text{ is the only vertex } t \text{ of } C \text{ with the property that } d(v, t) = 1 \text{ and } d(t, w) = d(v, w) - 1 \text{ if } 0 < d(v, w) < n.$$

It is obvious that  $(V, *)$  is a travel groupoid. Consider arbitrary  $x, y \in V$  such that  $d(x, y) = n$ . Then  $d(x * y, x * (y * x)) = 2$ . Thus  $(V, *)$  is not simple.

**Lemma 1.** Let  $(V, *)$  be a simple travel groupoid, let  $u, v, w \in V$ , and let  $k \geq 1$ . Assume that  $u *^{k-1} w \neq w$  and

$$(6) \quad v * u = w.$$

Then

$$(7) \quad u *^i v = u *^i w$$

and

$$(8) \quad v * (u *^i w) = w$$

for each  $i$ ,  $0 \leq i \leq k$ , and

$$u *^{k-1} v \neq v.$$

*Proof.* We will first prove that (7) and (8) hold for each  $i$ ,  $0 \leq i \leq k$ . We proceed by induction on  $i$ . Let first  $i = 0$ . Obviously,  $u *^0 v = u = u *^0 w$ . By (6),

$v * (u *^0 w) = w$ . Let now  $1 \leq i \leq k$ . By the induction hypothesis,

$$(9) \quad u *^{i-1} v = u *^{i-1} w$$

and

$$(10) \quad v * (u *^{i-1} w) = w.$$

Since  $u *^{k-1} w \neq w$ , Proposition 4 implies that  $u *^{i-1} w \neq w$ . By virtue of (10),  $v * (u *^{i-1} w) = w \neq u *^{i-1} w$ . As follows from (t3) and Remark 2,

$$(11) \quad (u *^{i-1} w) * v = (u *^{i-1} w) * (v * (u *^{i-1} w))$$

and

$$(12) \quad v * (u *^{i-1} w) = v * ((u *^{i-1} w) * v).$$

Obviously,  $u *^i v = (u *^{i-1} v) * v$ . It follows from (9), (11) and (10) that  $(u *^{i-1} v) * v = (u *^{i-1} w) * v = (u *^{i-1} w) * (v * (u *^{i-1} w)) = (u *^{i-1} w) * w$ . Thus  $u *^i v = u *^i w$  and (7) holds.

Next, as follows from (7), (9), (12) and (10),  $v * (u *^i w) = v * (u *^i v) = v * ((u *^{i-1} v) * v) = v * ((u *^{i-1} w) * v) = v * (u *^{i-1} w) = w$ . Thus (8) holds.

We want to prove now that  $u *^{k-1} v \neq v$ . Suppose, to the contrary, that  $u *^{k-1} v = v$ . By (7),  $u *^{k-1} v = u *^{k-1} w$ , and thus  $w \neq u *^{k-1} w = v$ . As follows from (8) and Proposition 1,  $w = v * (u *^{k-1} w) = v * v = v$ , which completes the proof.  $\square$

**Proposition 6.** *Let  $(V, *)$  be a simple travel groupoid, and let  $k \geq 1$ . If  $x, y \in V$ ,  $x *^{k-1} y \neq y$ , and  $x *^k y = y$ , then  $y *^{k-1} x \neq x$  and*

$$(13) \quad y *^j x = x *^{k-j} y$$

for each  $j$ ,  $0 \leq j \leq k$ .

*P r o o f.* We proceed by induction on  $k$ .

Let first  $k = 1$ . Consider arbitrary  $x, y \in V$  such that  $x *^0 y \neq y$  and  $x *^1 y = y$ . Then  $x \neq y$ . We have  $y *^0 x \neq x$ . Obviously, (13) holds for  $j = 0$ . As follows from (1), (13) holds also for  $j = 1$ .

Let now  $k \geq 2$ . Consider arbitrary  $x, y \in V$  such that  $x *^{k-1} y \neq y$  and  $x *^k y = y$ . Since  $x *^{k-1} y \neq y$ , it follows from (2) that  $x \neq y$ . Put  $z = x * y$ . Then  $z *^{k-2} y \neq y$  and  $z *^{k-1} y = y$ . By the induction hypothesis,  $y *^{k-2} z \neq z$  and

$$(14) \quad y *^j z = z *^{(k-1)-j} y = (x * y) *^{(k-1)-j} y = x *^{k-j} y$$

for each  $j$ ,  $0 \leq j \leq k - 1$ . Since  $y *^{k-2} z \neq z$  and  $x * y = z$ , Lemma 1 implies that

$$(15) \quad y *^j x = y *^j z$$

for each  $j$ ,  $0 \leq j \leq k - 1$ . Combining (14) and (15) for each  $j$ ,  $0 \leq j \leq k - 1$ , we get (13) for each  $j$ ,  $0 \leq j \leq k - 1$ . This means that  $y *^{k-1} x = x * y$ . Since  $x \neq y$ , (2) implies that  $y *^{k-1} x \neq x$ . Moreover, by (t1),  $y *^k x = (y *^{k-1} x) * x = (x * y) * x = x$ . Hence (13) holds also for  $j = k$ , which completes the proof.  $\square$

**Corollary 1.** *Let  $(T, *)$  be a simple travel groupoid, let  $u, v \in V$ , and let  $k \geq 1$ . If  $u *^k v = v$ , then  $v *^k u = u$ .*

**Proof.** The case of  $u = v$  is obvious. Let  $u \neq v$ . Then  $u *^0 v \neq v$ . Since  $u *^k v = v$ , we see that there exists  $i$ ,  $1 \leq i \leq k$ , such that  $u *^{i-1} v \neq v$  and  $u *^i v = v$ . By virtue of Proposition 6,  $v *^i u = u *^0 v = u$  and therefore

$$v *^k u = (v *^i u) *^{k-i} u = u *^{k-i} u = u,$$

which completes the proof.  $\square$

**Theorem 2.** *Let  $(V, *)$  be a simple travel groupoid on a graph  $G$ , let  $u, v \in V$  and let  $k \geq 1$ . Assume that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ . Then the sequence*

$$v *^0 u, \dots, v *^{k-1} u, v *^k u$$

*is a  $v - u$  path in  $G$ .*

**Proof.** Combining Theorem 1 and Proposition 6, we get the theorem.  $\square$

The next two lemmas will be used in Section 3.

**Lemma 2.** *Let  $(V, *)$  be a simple travel groupoid, let  $u, v \in V$ , and let  $j \geq 1$ . Assume that  $v *^j u \neq u$ . Then  $u * (v *^j u) = u * v$ .*

**Proof.** We proceed by induction on  $j$ . The case of  $j = 1$  immediately follows from the definition of a simple travel groupoid. Let  $j \geq 2$ . Since  $v *^j u \neq u$ , it follows from Proposition 4 that  $v *^{j-1} u \neq u$ . By the induction hypothesis,

$$(16) \quad u * (v *^{j-1} u) = u * v.$$

Obviously,  $v *^j u = (v *^{j-1} u) * u$ . Since  $(v *^{j-1} u) * u \neq u$ , (t3) implies that

$$(17) \quad u * ((v *^{j-1} u) * u) = u * (v *^{j-1} u).$$

Combining (16) and (17), we get  $u * (v *^j u) = u * ((v *^{j-1} u) * u) = u * v$ , which completes the proof.  $\square$



**Lemma 3.** Let  $(V, *)$  be a simple travel groupoid, let  $u, v \in V$ , and let  $k \geq 2$ . Assume that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ . Then

$$u *^i (v * u) = u *^i v$$

for each  $i$ ,  $0 \leq i \leq k - 1$ .

*Proof.* Proposition 6 implies that  $v *^{k-1} u \neq u$  and  $v *^k u = u$ . Put  $w = v * u$ . Then  $w *^{k-2} u \neq u$  and  $w *^{k-1} u = u$ . By Proposition 6 again,  $u *^{k-2} w \neq w$ . Lemma 1 implies that  $u *^i v = u *^i w = u *^i (v * u)$  for each  $i$ ,  $0 \leq i \leq k - 1$ , which completes the proof.  $\square$

**Remark 4.** Let  $(V, *)$  be a simple travel groupoid on a finite graph  $G$ . It was proved in [2] and [4] that  $G$  is a geodetic graph and  $(V, *)$  is its proper groupoid if and only if  $G$  is connected and  $(V, *)$  satisfies the following axiom

(tg) if  $w * v = v$  and  $u * v \neq u * w$ , then  $w * (u * v) = v$  (for all  $u, v, w \in V$ ).

The assumption that  $G$  is connected can not be deleted. There exists a simple travel groupoid satisfying (tg) on a finite disconnected graph (see Remark 2 in [2]).

### 3. NON-CONFUSING TRAVEL GROUPOIDS

Let  $(V, *)$  be a travel groupoid, and let  $u, v \in V$  such that  $u \neq v$ . By (2),  $u * v \neq u$  and by (t2),  $u *^2 v \neq u$ . If there exists  $i \geq 3$  such that  $u *^i v = u$ , then we say that the ordered pair  $(u, v)$  is a *confusing pair* in  $(V, *)$ .

The next lemma will be used in the next section.

**Lemma 4.** Let  $(V, *)$  be a travel groupoid, let  $u, v \in V$ ,  $u \neq v$ , and let  $i \geq 3$  such that  $u *^i v = u$ . Then there exists  $j$ ,  $3 \leq j \leq i$ , such that  $u *^j v = u$  and the elements

$$u *^0 v, \dots, u *^{j-2} v, \quad \text{and} \quad u *^{j-1} v$$

are pairwise distinct.

*Proof.* Since  $u \neq v$ , (t2) implies that there exists  $j$ ,  $3 \leq j \leq i$  such that  $u *^j v = u$  and all the elements

$$u *^1 v, u *^2 v, \dots, u *^{j-1} v$$

are different from  $u$ . Assume that there exist  $k$  and  $m$ ,  $1 \leq k < m \leq j - 1$ , such that  $u *^k v = u *^m v$ . Then  $m \geq k + 2$ . It is clear that

$$u *^n v \in \{u *^k v, u *^{k+1} v, \dots, u *^{m-1} v\} \quad \text{for all } n \geq m,$$

and therefore  $u *^i v \neq u$ , which is a contradiction. Thus the lemma is proved.  $\square$

**Remark 5.** Let  $(V, *)$  be a travel groupoid on a finite graph  $G$ . It is clear that if  $G$  is not connected, then  $(V, *)$  has a confusing pair.

We say that a travel groupoid  $(V, *)$  is *non-confusing* if there exists no confusing pair in  $(V, *)$ .

**Proposition 7.** *Let  $(V, *)$  be a finite non-confusing travel groupoid, and let  $u, v \in V$  and  $u \neq v$ . Then there exists exactly one  $k \geq 1$  such that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ .*

*Proof.* Define

$$u_i = u *^i v \quad \text{for all } i \geq 0.$$

Suppose, to the contrary, that

$$u_i \neq v \quad \text{for all } i \geq 0.$$

Since  $V$  is finite, there exist  $j$  and  $m$ ,  $0 \leq j < m$ , such that  $u_m = u_j$ . We have

$$u_m = u_j *^{m-j} v = u_j.$$

Thus  $m - j \geq 3$  and  $(u_j, v)$  is a confusing pair in  $(V, *)$ , which is a contradiction. We have prove that there exists  $k \geq 1$  such that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ . By Remark 1,  $k$  is defined uniquely. Thus the theorem is proved.  $\square$

**Theorem 3.** *Let  $(V, *)$  be a finite travel groupoid on a graph  $G$ . Then  $(V, *)$  is non-confusing if and only if the following statement holds for all distinct  $u, v \in V$ : there exists  $k \geq 1$  such that the sequence*

$$u *^0 v, \dots, u *^{k-1} v, u *^k v$$

*is an  $u - v$  path in  $G$ .*

*Proof.* Combining Theorem 1 and Proposition 7, we obtain the theorem.  $\square$

The next remark gives an example of a simple travel groupoid on a finite geodetic graph. This travel groupoid has a confusing pair.

**Remark 6.** Let  $m, n \geq 3$  be odd, and let  $u_0, u_1, \dots, u_{m-1}, v, w_0, w_1, \dots, w_{n-1}$  are pairwise distinct elements. Put

$$U = \{u_0, u_1, \dots, u_{m-1}, v\} \quad \text{and} \quad W = \{w_0, w_1, \dots, w_{n-1}, v\}.$$

Obviously,  $U \cap W = \{v\}$ . Define  $u_m = u_0$  and  $w_n = w_0$ . Let  $G_U$  be the graph with  $V(G_U) = U$  and

$$E(G_U) = \{u_0u_1, \dots, u_{m-2}u_{m-1}, u_{m-1}u_m, u_0v\}.$$

Moreover, let  $G_W$  be the graph with  $V(G_W) = W$  and

$$E(G_W) = \{w_0w_1, \dots, w_{n-2}w_{n-1}, w_{n-1}w_n, w_0v\}.$$

Since  $m$  and  $n$  are odd, we see that both  $G_U$  and  $G_W$  are geodetic graphs. At the end of Section 1, the mapping  $A_G$  was defined for a geodetic graph  $G$ . In the same way, we define the mappings  $A_{G_U}$  and  $A_{G_W}$  for the geodetic graphs  $G_U$  and  $G_W$  respectively.

Put  $V = U \cup W$ . We denote by  $*$  the binary operation on  $V$  defined for all  $x, y \in V$  as follows:

$$x * y = x \text{ if } x = y;$$

$$x * y = A_{G_U}(x, y) \text{ if } x, y \in U \text{ and } x \neq y;$$

$$x * y = A_{G_W}(x, y) \text{ if } x, y \in W \text{ and } x \neq y;$$

$$x * y = u_i \text{ if } x = u_{i-1} \text{ and } y \in W \setminus \{v\} \text{ for each } i, 0 \leq i \leq m - 1;$$

$$x * y = w_j \text{ if } x = w_{j-1} \text{ and } y \in U \setminus \{v\} \text{ for each } j, 0 \leq j \leq n - 1.$$

It is easy to see that  $(V, *)$  is a simple travel groupoid. The ordered pair  $(u_0, w_0)$  is an example of a confusing pair in  $(V, *)$ . Let  $G_0$  denote the graph of  $(V, *)$ . It is easy to see that  $G$  is a geodetic graph.

**Proposition 8.** *Let  $(V, *)$  be a finite simple non-confusing travel groupoid, let  $u, v \in V$ . Then  $(u *^i v) *^i u = u$  for each  $i \geq 0$ .*

*Proof.* The case of  $u = v$  follows immediately from Proposition 1. Assume that  $u \neq v$ . According to Proposition 7, there exists  $k \geq 1$  such that  $u *^{k-1} v \neq v$  and  $u *^k v = v$ . Recall that  $(V, *)$  is simple. If  $i \geq k$ , then  $u *^i v = v$  and, by virtue of Corollary 1,  $(u *^i v) *^i u = v *^i u = u$ . Let  $i < k$ . By Proposition 6,  $v *^k u = u$  and  $v *^{k-i} u = u *^i v$ . Thus  $(u *^i v) *^i u = (v *^{k-i} u) *^i u = v *^k u = u$ , which completes the proof.  $\square$

**Proposition 9.** *Let  $(V, *)$  be a finite simple non-confusing travel groupoid, and let  $u, v, w \in V$  such that  $u \neq v$ . Assume that there exists  $k \geq 1$  such that  $u *^{k-1} w \neq v$  and  $u *^k w = v$ . Then  $u *^{k-1} v \neq v$  and  $u *^k v = v$ .*

*Proof.* We see that  $u \neq w$  (otherwise,  $u *^k w = u \neq v$ ; a contradiction). By Proposition 7, there exists exactly one  $m \geq 1$  such that  $u *^{m-1} w \neq v$  and  $u *^m w = v$ .

If  $k > m$ , then  $u *^{k-1} w = u *^k w$ ; a contradiction. Thus  $k \leq m$ . As follows from Proposition 6,  $w *^{m-1} u \neq u$  and  $w *^{m-j} u = u *^j w$  for each  $j$ ,  $0 \leq j \leq m$ . Thus  $w *^m u = u$ . Since  $u *^k w = v$ , we get  $v = w *^{m-k} u$ . Hence

$$v *^{k-1} u = (w *^{m-k} u) *^{k-1} u \neq u \quad \text{and} \quad v *^k u = (w *^{m-k} u) *^k u = u.$$

If we apply Proposition 6 again, we get  $u *^{k-1} v \neq v$  and  $u *^k v = v$ , which completes the proof.  $\square$

Let  $(V, *)$  be a finite travel groupoid on a graph  $G$ , and let  $x, y \in V$ . Clearly,  $x$  and  $y$  are distinct and non-adjacent vertices of  $G$  if and only if  $x * y \neq y$ .

Let  $(V, *)$  be a simple non-confusing travel groupoid, and let  $x, y \in V$  such that  $x * y \neq y$ . By virtue of Proposition 7, there exists exactly one  $k \geq 2$  such that

$$(18) \quad x *^{k-1} y \neq y \quad \text{and} \quad x *^k y = y.$$

As follows from Proposition 6,

$$(19) \quad y *^{k-1} x \neq x \quad \text{and} \quad y *^k x = x.$$

Put

$$(20) \quad y = x' \quad \text{and} \quad x = y'.$$

Consider arbitrary  $u, v \in V$  such that  $u \in \{x, y\}$ . Assume that there exists  $j \geq 1$  such that  $u *^{j-1} v \neq u'$  and  $u *^j v = u'$ . By virtue of Remark 1 and Proposition 9,  $j = k$ . Moreover, Proposition 9 implies that

$$(21) \quad \text{if } u *^k v = u', \text{ then } u *^{k-1} v \neq u'.$$

By the *xy-strengthening* of  $*$  on  $V$  we mean the binary operation  $\circ$  on  $V$  defined for all  $u, v \in V$  as follows:

$$u \circ v = u *^k v \text{ if } u \in \{x, y\} \text{ and } u *^k v = u';$$

$$u \circ v = u * v \text{ otherwise.}$$

This means that  $w \circ w = w$  for every  $w \in V$ .

**Lemma 5.** *Let  $(V, *)$  be a finite simple non-confusing travel groupoid on a graph  $G$ , let  $x, y \in V$  such that  $x * y \neq y$ , and let  $\circ$  be the  $xy$ -strengthening of  $*$  on  $V$ . Then  $(V, \circ)$  is a simple non-confusing travel groupoid on  $G + xy$ .*

*Proof.* There exists exactly one  $k \geq 2$  such that (18) holds. Moreover, we have (19). Use the convention (20).

We first show that  $(V, \circ)$  satisfies the axioms (t1), (t2), and (t3) and that  $(V, \circ)$  is non-confusing. Consider arbitrary  $r, s \in V$ .

*Verification of (t1).* Put  $t = (r \circ s) \circ r$ . We will show that  $t = r$ . If  $r = s$ , then  $r \circ s = r$  and therefore  $t = r \circ r = r$ . Assume that  $r \neq s$ . Then there exists  $i \in \{1, k\}$  such that  $r \circ s = r *^i s$ . Since  $r \neq s$  and  $(V, *)$  is non-confusing, we get  $r *^i s \neq r$ . There exists  $j \in \{1, k\}$  such that  $t = (r *^i s) \circ r = (r *^i s) *^j r$ . If  $i = 1$ , then (t1) and Proposition 1 imply that  $t = r$ . Assume that  $i = k$ . Then  $r \in \{x, y\}$  and  $r *^i s = r'$ . This implies that  $j = k$ . We get  $t = (r *^k s) *^k r = r$  again.

*Verification of (t2).* Obviously, there exist  $i, j \in \{1, k\}$  such that  $(r \circ s) \circ s = (r *^i s) *^j s = r *^{i+j} s$ . Let  $(r \circ s) \circ s = r$ . Since  $(V, *)$  is non-confusing, we get  $r = s$ .

We see that  $(V, \circ)$  is a travel groupoid.

*Verification of (t3).* Assume that  $s \circ r \neq r$ . As follows from (2),  $s \neq r$ . We will prove that  $r \circ (s \circ r) = r \circ s$ . If  $r, s \in \{x, y\}$ , then  $s *^k r = r$  and therefore  $s \circ r = r$ , which is a contradiction. Thus

$$(22) \quad \text{at most one of } r \text{ and } s \text{ belongs to } \{x, y\}.$$

Since  $(V, *)$  is non-confusing, Proposition 7 implies that there exists  $m \geq 1$  such that  $r *^{m-1} s \neq s$  and  $r *^m s = s$ . Recall that  $k \geq 2$ . Since  $(V, *)$  is simple, it follows from Lemma 3 that

$$(23) \quad \text{if } k < m, \text{ then } r *^k (s * r) = r *^k s.$$

Recall that  $s \circ r \neq r$ . Since  $s \circ r = s *^k r$  or  $s * r$ , Remark 1 implies that  $s * r \neq r$ . By (t3),

$$r * (s * r) = r * s.$$

Let first  $r \in \{x, y\}$  and  $r *^k s = r'$ . Then  $r \circ s = r *^k s = r'$ . By (22),  $s \neq r'$ . Then  $k < m$  and  $s \circ r = s * r$ . It follows from (23) that  $r *^k (s * r) = r'$  and therefore

$$r \circ (s \circ r) = r \circ (s * r) = r *^k (s * r) = r *^k s = r' = r \circ s.$$

Let now  $r \in \{x, y\}$  and  $r *^k s \neq r'$ . Then  $r \circ s = r * s$ . By (22),  $s \neq r'$  and thus  $s \circ r = s * r$ . Assume that  $r *^k (s * r) = r'$ . Then  $r \circ (s * r) = r *^k (s * r)$ . As follows

from (21),  $r *^{k-1} (s * r) \neq r'$ . By virtue of Lemma 1,  $r *^k s = r *^k (s * r) = r'$ , which is a contradiction. Thus  $r \circ (s * r) = r * (s * r)$  and therefore  $r \circ (s \circ r) = r * s = r \circ s$ .

Finally, let  $r \notin \{x, y\}$ . Then  $r \circ s = r * s$  and  $r \circ (s \circ r) = r * (s \circ r)$ . Assume that  $s \in \{x, y\}$  and  $s *^k r = s'$ . Then  $s \circ r = s *^k r$ . Since  $s *^k r \neq r$ , Lemma 2 implies that  $r \circ (s \circ r) = r * (s *^k r) = r * s = r \circ s$ . If

$$s \notin \{x, y\} \quad \text{or} \quad (s \in \{x, y\} \text{ and } s *^k r \neq s'),$$

then  $s \circ r = s * r$  and therefore  $r \circ (s \circ r) = r * (s * r) = r * s = r \circ s$ .

Thus  $(V, \circ)$  is simple.

Assume that  $r \neq s$  and there exists  $i \geq 1$  such that  $r \circ^i s = r$ . Clearly, there exists  $m \geq i$  such that  $r \circ^i s = r *^m s$ . We have that  $r *^m s = r$ , which contradicts the fact that  $(V, *)$  is non-confusing. Thus  $(V, \circ)$  is non-confusing, too.

Recall that  $x \neq y$  and  $x \circ y = y$ . We can see that  $(V, \circ)$  is a simple non-confusing travel groupoid on  $G + xy$ , which completes the proof of the lemma.  $\square$

**Theorem 4.** *For every finite connected graph  $G$  there exists a simple non-confusing travel groupoid on  $G$ .*

**Proof.** Put  $V = V(G)$  and  $\beta(G) = |E(G)| - |V| + 1$ . We proceed by induction on  $\beta(G)$ . Obviously,  $\beta(G) \geq 0$ . Let first  $\beta(G) = 0$ . Then  $G$  is a tree. It is easy to see that its proper groupoid is simple and non-confusing. Let now  $\beta(G) \geq 1$ . Then there exist distinct  $x, y \in V$  such that  $x$  and  $y$  are adjacent in  $G$  and  $G - xy$  is connected. By the induction hypothesis, there exists a simple non-confusing travel groupoid  $(V, *)$  on  $G - xy$ . Lemma 5 implies that there exists a simple non-confusing groupoid on  $G$ , which completes the proof.  $\square$

#### 4. SMOOTH AND SEMI-SMOOTH TRAVEL GROUPOIDS

We say that a travel groupoid  $(V, *)$  is *smooth* if it satisfies the following axiom (t4) if  $u * v = u * w$ , then  $u * (w * v) = u * v$  (for all  $u, v, w \in V$ ).

Moreover, we say that a travel groupoid  $(V, *)$  is *semi-smooth* if it satisfies the following axiom

(t5) if  $u * v = u * w$ , then  $u * (v * w) = u * v$  or  $u * ((v * w) * w) = u * v$  (for all  $u, v, w \in V$ ).

Obviously, every smooth travel groupoid is semi-smooth.

**Proposition 10.** *Every semi-smooth travel groupoid is non-confusing.*

*Proof.* Let  $(V, *)$  be a semi-smooth travel groupoid. Obviously, there exists a graph  $G$  such that  $(V, *)$  is on  $G$ . Suppose, to the contrary, that there exists a confusing pair in  $(V, *)$ . As follows from Lemma 4, there exist  $u, w \in V$  and  $k \geq 3$  such that  $u \neq w$  and  $u *^k w = u$ , and the vertices  $u *^0 w, \dots, u *^{k-2} w$  and  $u *^{k-1} w$  are pairwise distinct. Define

$$u_i = u *^i w \quad \text{for } i = 0, 1, \dots, k.$$

Hence  $u_0 \neq u_1 \neq u_{k-1} \neq u_k = u_0$ . Obviously,  $u_0 * u_1 = u_1$  and  $u_0 * u_k = u_0 * u_0 \neq u_1$ . Moreover,  $u_0 * u_{k-1} = u_k * u_{k-1} = (u_{k-1} * w) * u_{k-1}$  and thus, by (t1),  $u_0 * u_{k-1} = u_{k-1}$ . This implies that there exist  $j$ ,  $0 \leq j \leq k-2$ , such that  $u_0 * u_j = u_1$  and  $u_0 * u_{j+1} \neq u_1 \neq u_0 * u_{j+2}$ . We have  $u_0 * w = u_1 = u_0 * u_j$ ,  $u_0 * (u_j * w) = u_0 * u_{j+1} \neq u_0 * u_j$ , and  $u_0 * ((u_j * w) * w) = u_0 * u_{j+2} \neq u_0 * u_j$ , which contradicts (t5). Thus the proposition is proved.  $\square$

**Proposition 11.** *Every complete bipartite graph has a simple smooth travel groupoid.*

*Proof.* Let  $G$  be complete bipartite graph. Put  $V = V(G)$ . There exist nonempty sets  $U$  and  $U'$  such that  $U \cap U' = \emptyset$ ,  $U \cup U' = V$  and the following statement holds for all distinct  $v, w \in V$ :

$$v \text{ and } w \text{ are adjacent in } G \text{ if and only if } |\{v, w\} \cap U| = 1 = |\{v, w\} \cap U'|.$$

Recall that  $U$  and  $U'$  are nonempty. Choose a vertex  $u \in U$  and a vertex  $u' \in U'$ . We denote by  $*$  the binary operation on  $V$  defined as follows:

- $x * y = x$  if  $x = y$ ;
- $x * y = y$  if  $x$  and  $y$  are adjacent in  $G$ ;
- $x * y = u'$  if  $x, y \in U$  and  $x \neq y$ ;
- $x * y = u$  if  $x, y \in U'$  and  $x \neq y$ .

It can be easily verified that  $(V, *)$  satisfies (t1), (t2), (t3), and (t4). Hence  $(V, *)$  is a simple smooth travel groupoid.  $\square$

Recall that every tree is a geodetic graph and that the proper groupoid of every geodetic graph is a simple travel groupoid.

**Proposition 12.** *The proper groupoid of every tree is a smooth travel groupoid.*

*Proof* is easy.  $\square$

Note that every complete graph is geodetic. Obviously, the proper groupoid of every complete graph is a smooth travel groupoid.

**Theorem 5.** *Let  $G$  be a geodetic graph of diameter two, and let  $(V, *)$  be the proper groupoid on  $G$ . Then  $(V, *)$  is a smooth travel groupoid.*

*Proof.* Clearly,  $(V, *)$  is a simple travel groupoid such that

$$(24) \quad x *^2 y = y \quad \text{for all } x, y \in V.$$

We will prove that  $(V, *)$  is smooth. Suppose, to the contrary, that  $(V, *)$  is not smooth. Then there exist  $u, v, w \in V$  such that  $u * v = u * w$  and

$$(25) \quad u * (v * w) \neq u * v.$$

This implies that  $v \neq v * w \neq w$ . By (24),  $v *^2 w = w$ . Recall that  $(V, *)$  is simple. Since  $v * w \neq w$ , Proposition 6 implies that  $w * v \neq v$ ,  $w *^2 v = v$ , and  $w * v = v * w$ . Thus the sequence

$$v, v * w = w * v, w$$

is a shortest  $v - w$  path in  $G$ .

Put  $t = u * v$ . Then  $t = u * w$ . As follows from (24),  $t * v = v$  and  $t * w = w$ . By (1),  $v * t = t$ . If  $t = v$ , then  $u * w = v$  and therefore  $u *^2 w = v * w \neq w$ ; a contradiction. Thus  $t \neq v$ . Since  $v * t = t$ , we see that  $v$  and  $t$  are adjacent in  $G$ . Let  $t = w$ . Since  $t * v = v$ , we get  $w * v = v$ ; a contradiction. Thus  $t \neq w$ . Since  $t * w = w$ , we see that  $t$  is adjacent to  $w$ . Thus the sequence

$$v, t, w$$

is a shortest  $v - w$  path in  $G$ .

Assume that  $t = v * w$ . Using (3), we see that  $u * (v * w) = u * (u * v) = u * v$ , which contradicts (25). Thus  $u * v \neq v * w$ . We see that  $G$  has two distinct shortest  $v - w$  paths in  $G$ . This means that  $G$  is not a geodetic graph, which is a contradiction. Thus the theorem is proved.  $\square$

We pose two questions.

**Question 1.** Does there exist a geodetic graph  $G$  such that the proper groupoid of  $G$  is not smooth? (If so, does there exist a geodetic groupoid  $G$  such that the proper groupoid of  $G$  is not semi-smooth?)

**Question 2.** Does there exist a connected graph  $G$  such that  $G$  has no smooth travel groupoid? (If so, does there exist a connected graph  $G$  such that  $G$  has no semi-smooth travel groupoid?)



## 5. GRAPHS WITH TRAVEL GROUPOIDS

Recall that, by Theorem 4, every finite connected graph has a simple non-confusing travel groupoid.

**Theorem 6.** *Let  $G$  be a finite graph. Then  $G$  has a travel groupoid if and only if  $G$  is connected or  $G$  is disconnected and no component of  $G$  is a tree.*

*Proof.* Assume that  $G$  is connected or  $G$  is disconnected and no component of  $G$  is a tree. If  $G$  is connected, then, by Theorem 4, there exists a travel groupoid on  $G$ . Let  $G$  be disconnected. Then every component of  $G$  contains a cycle. It is easy to see that there exists a mapping  $f$  of  $V(G)$  into itself such that the following statements hold for every  $u \in V(G)$ :

$$u \text{ and } f(u) \text{ are adjacent vertices in } G$$

and

$$u \neq f(f(u)).$$

By virtue of Theorem 4, every component  $F$  of  $G$  has a travel groupoid, say a travel groupoid  $(V(F), *F)$ . For all  $x, y \in V(G)$ , we define

$$x * y = x *_{H} y \text{ if there exists a component } H \text{ of } G \text{ such that } x, y \in V(H)$$

and

$$x * y = f(x) \text{ if } x \text{ and } y \text{ belong to distinct components of } G.$$

It is easy to see that  $(V(G), *)$  satisfies (t1) and (t2). Hence  $G$  has a travel groupoid.

Conversely, assume that  $G$  is disconnected and at least one component  $T$  of  $G$  is a tree. Suppose, to the contrary, that  $G$  has a travel groupoid, say a travel groupoid  $(V, *)$ , where  $V = V(G)$ . Consider  $u \in V(T)$  and  $v \in V(G) \setminus V(T)$ . Since  $V(T)$  is finite and  $T$  contains no cycle, we see that there exists  $k \geq 1$  such that  $u *^{k+1} v = u *^{k-1} v$ . We have  $((u *^{k-1} v) * v) * v = u *^{k-1} v$ , and thus, by (t2),  $u *^{k-1} v = v$ . Proposition 3 implies that  $u$  and  $v$  belong to the same component of  $G$ , which is a contradiction. Thus the theorem is proved.  $\square$

**Question 3.** Does there exist an infinite graph  $G$  with no finite components such that  $G$  has no travel groupoid?

**Remark 7.** Let  $(V, *)$  be a finite travel groupoid. Put

$$X = \{(u, v, w) : u, v, w \in V \text{ and } v = u * w\}.$$

Then  $(V, X)$  is a signpost system in the sense of [5]. We say that  $(V, X)$  is the signpost system of  $(V, *)$ .

The signpost systems of travel groupoids create a special subclass of the class of all signpost systems. The terms “simple”, “non-confusing” and “smooth” introduced in the present paper for travel groupoids are inspired by the same terms used for signpost systems in [5].

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