

Václav Soukup

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ON TRANSFER FUNCTION MATRIX OF LINEAR CASCADE SYSTEMS

VÁCLAV SOUKUP

Lower triangular transfer function matrix of cascade linear system is investigated. The paper aims to show and prove that unlike for the other multi-input, multi-output systems general fashions of the coprime transfer function matrix fraction representations can be found in this case.

1. CASCADE SYSTEM TRANSFER FUNCTION MATRIX

Many technological, industrial and other processes are characterized by an one way, one line flow of information and energy ([2]). Such cascade processes represent the special case of multi-input, multi-output structure that may be modelled by the block diagram in Fig. 1 where

- P_i denotes the i th stage of the process,
- y_i the output of the i th stage (i th controlled variable), and
- u_i the input of the i th stage (i th control variables);
- $i = 1, \dots, n$.

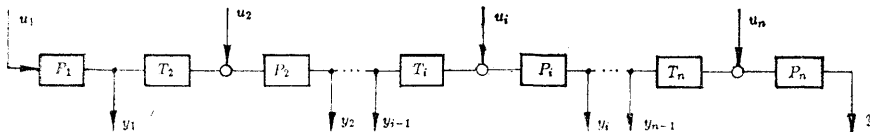


Fig. 1.

Feed properties between parts P_{i-1} and P_i of the process are modelled by elements T_i , $i = 2, \dots, n$.

The paper deals with the regular n -input, n -output cascade only which is affected by no external disturbances.

Assuming continuous-time, linear dynamics of P_i as well as T_i the following

equations can be written (in Laplace transforms):

$$(1) \quad Y_1(s) = P_1(s) U_1(s)$$

and

$$Y_i(s) = P_i(s) [U_i(s) + T_i(s) Y_{i-1}(s)], \quad i = 2, \dots, n.$$

The transfer functions

$$P_i(s) = \frac{b_i(s)}{a_i(s)}, \quad (a_i, b_i) \sim 1; \quad i = 1, \dots, n$$

(2) and

$$T_i(s) = \frac{q_i(s)}{p_i(s)}, \quad (p_i, q_i) \sim 1; \quad i = 2, \dots, n$$

where $a_i(s)$, $b_i(s)$, $p_i(s)$ and $q_i(s)$ are supposed to be polynomials in s . Nevertheless the final results which come in Theorems 1 and 2 of the paper are also applicable if factors $\exp(-\tau_i s)$ corresponding to possible dead times occur in $b_i(s)$ and/or $q_i(s)$.

Note that the standard symbols of polynomial theory ([1]) are used in the paper:

(a, b) for the greatest common divisor (GCD) of a and b ,

$a \sim b$ if a and b are associates, i.e., $a = \alpha b$ where $\alpha = \text{const.}$ (a polynomial of degree 0) and

$b \mid a$ if b is a divisor of a .

Combining the equations (1) we can write in vector-matrix form

$$(3) \quad \mathbf{Y}(s) = \mathbf{G}(s) \mathbf{U}(s)$$

where

$$(4) \quad \mathbf{Y}(s) = [Y_1(s); \dots; Y_n(s)]^T \quad \text{and} \quad \mathbf{U}(s) = [U_1(s); \dots; U_n(s)]^T$$

and the system transfer function matrix (TFM)

$$(5) \quad \mathbf{G}(s) = \begin{bmatrix} G_{11}(s) & & & 0 \\ G_{21}(s) & G_{22}(s) & & \\ \vdots & \vdots & \ddots & \\ G_{n1}(s) & G_{n2}(s) & \dots & G_{nn}(s) \end{bmatrix}$$

is lower triangular ($n \times n$) matrix with the elements

$$G_{ij}(s) = \frac{b_i(s)}{a_i(s)} \quad \text{for } j = i$$

$$G_{ij}(s) = \frac{b_i(s) \dots b_j(s) q_i(s) \dots q_{j+1}(s)}{a_i(s) \dots a_j(s) p_i(s) \dots p_{j+1}(s)} \quad \text{for } j < i$$

$$G_{ij}(s) = 0 \quad \text{for } j > i, \quad i, j = 1, \dots, n.$$

The structure given by Fig. 1 includes usual cases which we encounter in practice. Especially, either

- T_i represents the process part which is affected by no external input and has no own measured output or

— two neighbouring $(i - 1)$ th and i th stages of a process are coupled through a physical transducer T_i to adapt signal y_{i-1} to the actuator of P_i .

Fraction representations of the matrix $\mathbf{G}(s)$ are investigated in the next sections. It is well known that the coprime matrix fraction (CMF) descriptions of a general TFM can be found by numerical ways only (cf. [1]). But numerical algorithms need not be applied in the case of cascade system matrix. It will be shown and proved that \mathbf{G} given by (5) with (6) can be usually transformed into CMF fashion immediately using the elements a_i , b_i , p_i and q_i of single transfer functions in \mathbf{G} .

The reader is assumed to be acquainted with the main fundamentals of polynomial and polynomial matrix theory of dynamic systems ([1], [3]).

2. LEFT COPRIME MATRIX FRACTION REPRESENTATION OF \mathbf{G}

Any $(n \times r)$ TFM \mathbf{G} of a linear, free of dead times, continuous-time system can be always written in the form (cf. [1])

$$(7) \quad \mathbf{G} = \mathbf{A}_L^{-1} \mathbf{B}_L$$

where \mathbf{A}_L and \mathbf{B}_L are $(n \times n)$ and $(n \times r)$ polynomial matrices in s , respectively. The corresponding input-output equation

$$(8) \quad \mathbf{A}_L \mathbf{Y} = \mathbf{B}_L \mathbf{U}$$

is valid.

The matrices \mathbf{A}_L and \mathbf{B}_L represent a left matrix fraction description of \mathbf{G} . Such a representation is left coprime (LCMF) if and only if

$$(9) \quad \mathbf{A}_L = \mathbf{D} \mathbf{F}_L \quad \text{and} \quad \mathbf{B}_L = \mathbf{D} \mathbf{H}_L$$

where the $(n \times n)$ polynomial matrix \mathbf{D} known as the greatest common left divisor (GCLD) of \mathbf{A}_L and \mathbf{B}_L has the property $\det \mathbf{D} \sim 1$, i.e., \mathbf{D} is unimodular. The pairs $\mathbf{A}_L, \mathbf{F}_L$ and $\mathbf{B}_L, \mathbf{H}_L$ are then the pairs of left equivalent polynomial matrices; \mathbf{A}_L and \mathbf{B}_L are called *left coprime matrices*.

Theorem 1. LCMF representation of a cascade system TFM \mathbf{G} standing in (5) with (6) can be written in the form

$$(10) \quad \mathbf{A}_L = \begin{bmatrix} a_1 & & & 0 \\ -b_2 q_2 & a_2 p_2 & & \\ & & \ddots & \\ 0 & & -b_n q_n & a_n p_n \end{bmatrix}$$

and

$$(11) \quad \mathbf{B}_L = \text{diag} [b_1; b_2 p_2; \dots; b_n p_n]$$

if and only if

$$(12) \quad (p_i, b_i b_{i-1}) \sim 1 \quad \text{for any } i = 2, \dots, n.$$

Proof.

A. At first we must prove that the matrices (10) and (11) represent a left matrix fraction description of \mathbf{G} at all.

Using (10) and (11) in the equation (8) we get

$$(13) \quad a_1 Y_1 = b_1 U_1$$

and

$$-b_i q_i Y_{i-1} + a_i p_i Y_i = b_i p_i U_i \quad \text{for } i = 2, \dots, n.$$

Hence

$$(14) \quad Y_1 = \frac{b_1}{a_1} U_1$$

and

$$Y_i = \frac{b_i q_i}{a_i p_i} Y_{i-1} + \frac{b_i}{a_i} U_i \quad \text{for } i = 2, \dots, n.$$

Gradual substitutions Y_i into the $(i + 1)$ th equation in (13) $i = 1, \dots, n - 1$, yield the relation $\mathbf{Y} = \mathbf{G}\mathbf{U}$ with \mathbf{G} standing in (5) with (6). The same result can be obtained if the inverse of \mathbf{A}_L is formed and then $\mathbf{G} = \mathbf{A}_L^{-1} \mathbf{B}_L$ determined. Thus (10) and (11) is a left matrix fraction representation of \mathbf{G} .

B. Secondly, it must be proved that the matrices (10) and (11) are left coprime if and only if the conditions (12) are true.

If: Assume a polynomial matrix \mathbf{D} to be GCLD of \mathbf{A}_L and \mathbf{B}_L and denote $d = \det \mathbf{D}$. The expressions (9) can be written as

$$(15) \quad [\mathbf{a}_1; \dots; \mathbf{a}_n] = \mathbf{D}[\mathbf{f}_1; \dots; \mathbf{f}_n] \quad \text{and} \quad [\mathbf{b}_1; \dots; \mathbf{b}_n] = \mathbf{D}[\mathbf{h}_1; \dots; \mathbf{h}_n]$$

where $\mathbf{a}_i, \mathbf{b}_i, \mathbf{f}_i$ and \mathbf{h}_i are the columns of $\mathbf{A}_L, \mathbf{B}_L, \mathbf{F}_L$ and \mathbf{H}_L , respectively; $i = 1, \dots, n$. Obviously $\mathbf{a}_i = \mathbf{D}\mathbf{f}_i$ and $\mathbf{b}_i = \mathbf{D}\mathbf{h}_i$.

Now the set $\{\mathbf{M}_n\}$ of $(n \times n)$ polynomial matrices \mathbf{M}_n can be considered each of them being constructed as a different combination of \mathbf{a}_i and \mathbf{b}_j ; $i, j = 1, \dots, n$, i.e., \mathbf{M}_n is formed by n mutually different columns which are taken from $2n$ columns. Combining \mathbf{f}_i and \mathbf{h}_j in similar way the set $\{\mathbf{L}_n\}$ of $(n \times n)$ polynomial matrices \mathbf{L}_n is defined. Note that $\mathbf{A}_L, \mathbf{B}_L \in \{\mathbf{M}_n\}$ and $\mathbf{F}_L, \mathbf{H}_L \in \{\mathbf{L}_n\}$. Then

$$(16) \quad \mathbf{M}_n = \mathbf{D}\mathbf{L}_n$$

and hence

$$(17) \quad \det \mathbf{M}_n = d \det \mathbf{L}_n$$

is true for any corresponding pair of $\mathbf{M}_n \in \{\mathbf{M}_n\}$ and $\mathbf{L}_n \in \{\mathbf{L}_n\}$.

Let all nonzero determinants m_n of all matrices \mathbf{M}_n constitute the set $\{m_n\}$ and the GCD of all m_n be denoted by (m_n) .

It follows from (17) that

$$(18) \quad d \mid (m_n).$$

Hence \mathbf{A}_L and \mathbf{B}_L are LCMF of \mathbf{G} if $(m_n) \sim 1$. But it is not easy to survey the finding

of $\{m_n\}$ and (m_n) for general n . That is the set $\{\mathbf{M}_n\}$ contains $(2n)!(n!)^{-2}$ matrices. To overcome this difficulty we start with $n = 2$.

Here we determine

$$\{\mathbf{M}_2\} = \left\{ \begin{bmatrix} a_1 & 0 \\ -b_2q_2 & a_2p_2 \end{bmatrix}, \begin{bmatrix} b_1 & 0 \\ 0 & a_2p_2 \end{bmatrix}, \begin{bmatrix} a_1 & b_1 \\ -b_2q_2 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ b_2p_2 & a_2p_2 \end{bmatrix}, \right. \\ \left. \begin{bmatrix} a_1 & 0 \\ -b_2q_2 & b_2p_2 \end{bmatrix}, \begin{bmatrix} b_1 & 0 \\ 0 & b_2p_2 \end{bmatrix} \right\},$$

$$\{m_2\} = \{a_1a_2p_2, b_1a_2p_2, b_1b_2q_2, a_1b_2p_2, b_1b_2p_2\} \quad \text{and} \quad (m_2) = (p_2, b_2b_1)$$

Then $(p_2, b_2b_1) \sim 1$ ensures $(m_2) \sim 1$.

For $n = 3$ we obtain

$$\{m_3\} = \{a_1a_2p_2a_3p_3, b_1a_2p_2a_3p_3, a_1a_2p_2b_3p_3, b_1a_2p_2b_3p_3, b_1b_2q_2a_3p_3, \\ b_1b_2q_2b_3p_3, b_1b_2q_2b_3q_3, a_1b_2p_2a_3p_3, a_1b_2p_2b_3p_3, a_1b_2p_2b_3q_3, \\ b_1b_2p_2a_3p_3, b_1b_2p_2b_3q_3, b_1b_2p_2b_3p_3\}$$

and

$$(m_3) = (p_2, b_2b_1) \left(p_3, b_3 \frac{b_2}{(b_2, p_2)} (b_2, b_1, p_2) \right).$$

Obviously $(p_2, b_2b_1) \sim 1$ together with $(p_3, b_3b_2) \sim 1$ ensures $(m_3) \sim 1$. Thus far sufficiency of (12) for $n = 2$ and $n = 3$ has been proved.

The structure of determinants m_n for an increasing n must be studied for a general proof.

One can see that

$$(19) \quad m_n = c_1c_2 \dots c_n$$

where either a_i or b_i stands at the position c_i and either a_ip_i or b_ip_i or b_iq_i at the position c_i , $i = 2, \dots, n$. But not all combinations occur in m_n since $c_i = b_iq_i$ can succeed to $c_{i-1} = b_{i-1}p_{i-1}$ or $c_{i-1} = b_{i-1}q_{i-1}$ only.

Then we can decompose

$$(20) \quad \{m_n\} = \{m_{nA}\} \cup \{m_{nB}\} \cup \{m_{nQ}\}$$

where the subset

$$\begin{aligned} \{m_{nA}\} & \text{ contains all } m_n \text{ ended by } c_n = a_np_n, \\ \{m_{nB}\} & \text{ contains all } m_n \text{ ended by } c_n = b_np_n, \quad \text{and} \\ \{m_{nQ}\} & \text{ contains all } m_n \text{ ended by } c_n = b_nq_n. \end{aligned}$$

Let $\lambda\{\cdot\}$ denotes the number of determinants within a set $\{\cdot\}$. Since

$$(21) \quad \begin{aligned} \lambda\{m_{nA}\} &= \lambda\{m_{nB}\} = \lambda\{m_{n-1}\}, \quad \text{and} \\ \lambda\{m_{nQ}\} &= \lambda\{m_{n-1,B}\} + \lambda\{m_{n-1,Q}\} \end{aligned}$$

the recurrent relation

$$\lambda\{m_n\} = 2\lambda\{m_{n-1}\} + \lambda\{m_{n-1,B}\} + \lambda\{m_{n-1,Q}\}$$

is true starting with

$$\lambda\{m_2\} = \lambda\{m_{2A}\} + \lambda\{m_{2B}\} + \lambda\{m_{2Q}\}$$

where

$$\lambda\{m_{2A}\} = \lambda\{m_{2B}\} = 2 \quad \text{and} \quad \lambda\{m_{2Q}\} = 1.$$

Suppose now that $(m_j) \sim 1$ is ensured by the conditions

$$(22) \quad (p_i, b_i b_{i-1}) \sim 1 \quad \text{for any } i = 2, \dots, j < n.$$

Considering $i = j + 1$ we need to show that together with (22) the only additional condition

$$(23) \quad (p_{j+1}, b_{j+1} b_j) \sim 1$$

is sufficient to satisfy $(m_{j+1}) \sim 1$.

According to (20) and (21)

$$\begin{aligned} \{m_{j+1}\} &= \{m_{j+1,A}\} \cup \{m_{j+1,B}\} \cup \{m_{j+1,Q}\} = \\ &= \{m_j a_{j+1} p_{j+1}\} \cup \{m_j b_{j+1} p_{j+1}\} \cup \{m_{jB} b_{j+1} q_{j+1}\} \cup \{m_{jQ} b_{j+1} q_{j+1}\} \end{aligned}$$

and then

$$(24) \quad (m_{j+1}) \sim (p_{j+1}, b_{j+1}(m_{jB}, m_{jQ})).$$

Considering (22) we determine

$$(m_{jB}, m_{jQ}) \sim (b_j p_j, b_{j-1} p_{j-1} b_j q_j, b_{j-1} q_{j-1} b_j q_j) \sim b_j (p_j, b_{j-1}) \sim b_j$$

and hence taking into account (24)

$$(25) \quad (m_{j+1}) \sim (p_{j+1}, b_{j+1} b_j).$$

Consequently $(m_{j+1}) \sim 1$ is ensured by (23) for any $j < n$ if (22) are valid. Hence \mathbf{A}_L and \mathbf{B}_L are left coprime if (12) are true.

Only if: The conditions (12) can be decomposed for any i into two separate relations:

$$(26) \quad (p_i, b_i) \sim 1$$

and

$$(27) \quad (p_i, b_{i-1}) \sim 1.$$

Assume now that \mathbf{A}_L and \mathbf{B}_L standing in (10) and (11), resp., are left coprime but (12) are not valid.

1. If (12) are broken by $(p_i, b_i) \sim 1$ for one $i \in [2, n]$ we can denote

$$(28) \quad p_{ii} = \frac{p_i}{(p_i, b_i)} \quad \text{and} \quad b_{ii} = \frac{b_i}{(p_i, b_i)}.$$

with the GCLD

$$(35) \quad \mathbf{D} = \begin{bmatrix} 1 & & & & & & & \\ 0 & 1 & & & & & & 0 \\ & & \ddots & & & & & \\ & & & 0 & 1 & & & \\ & & & & & x & (p_i, b_{i-1}) & \\ & & & & & 0 & 1 & \\ & 0 & & & & & & \ddots & \\ & & & & & & & & 0 & 1 \end{bmatrix}$$

The polynomials x and y satisfy the equation

$$(36) \quad a_{i-1}p_{i-1}x + (p_i, b_{i-1})y = -b_iq_i$$

which is always solvable seeing that $(p_{i-1}, b_{i-1}) \sim 1$ as well as $(a_{i-1}, b_{i-1}) \sim 1$ is assumed.

Having x, y the remaining polynomials in (33) and (34) are

$$(37) \quad \begin{aligned} v &= b_{i-1,i}q_{i-1}x, \quad \text{and} \\ z &= -b_{i-1,i}p_{i-1}x. \end{aligned}$$

For $i = 2$ we put $p_{i-1} = 1$ and $q_{i-1} = 0$ in (36) and (37); the polynomial v in A_{LO} is omitted.

In virtue of (35) we have $d = \det \mathbf{D} = (p_i, b_{i-1}) \sim 1$ and consequently the matrices (10) and (11) are not left coprime. \square

3. RIGHT COPRIME MATRIX FRACTION REPRESENTATION OF \mathbf{G}

Any $(n \times r)$ TFM \mathbf{G} of a linear, free of dead times, continuous-time system can be always written also in the form (cf. [1])

$$(38) \quad \mathbf{G} = \mathbf{B}_R \mathbf{A}_R^{-1}$$

where \mathbf{A}_R and \mathbf{B}_R are $(r \times r)$ and $(n \times r)$ polynomial matrices in s , respectively.

The matrices \mathbf{A}_R and \mathbf{B}_R represent a right matrix fraction description of \mathbf{G} . Such a representation is right coprime (RCMF) if and only if

$$(39) \quad \mathbf{A}_R = \mathbf{F}_R \mathbf{D} \quad \text{and} \quad \mathbf{B}_R = \mathbf{H}_R \mathbf{D}$$

where the $(r \times r)$ polynomial matrix \mathbf{D} is the greatest common right divisor (GCRD) of \mathbf{A}_R and \mathbf{B}_R and $\det \mathbf{D} \sim 1$, i.e., \mathbf{D} is unimodular. The pairs $\mathbf{A}_R, \mathbf{F}_R$ and $\mathbf{B}_R, \mathbf{H}_R$ are then the pairs of right equivalent polynomial matrices; \mathbf{A}_R and \mathbf{B}_R are called *right coprime matrices*.

Theorem 2. RCMF representation of a cascade system TFM \mathbf{G} standing in (5)

with (6) can be written in the form

$$(40) \quad \mathbf{A}_R = \begin{bmatrix} a_1 p_2 & & & & \\ -b_1 q_2 & a_2 p_3 & & & 0 \\ & & \ddots & & \\ 0 & -b_{n-2} q_{n-1} & a_{n-1} p_n & & \\ & & & -b_{n-1} q_n & a_n \end{bmatrix}$$

and

$$(41) \quad \mathbf{B}_R = \text{diag} [b_1 p_2; b_2 p_3; \dots; b_{n-1} p_n; b_n]$$

if and only if

$$(42) \quad (p_i, b_i b_{i-1}) \sim 1 \quad \text{for any } i = 2, \dots, n.$$

Proof.

A. Considering $r = n$ for a regular cascade system and using an auxiliary $(n \times 1)$ vector signal \mathbf{X} we can express according to (38)

$$(43) \quad \mathbf{Y} = \mathbf{B}_R \mathbf{A}_R^{-1} \mathbf{U} = \mathbf{B}_R \mathbf{X}$$

and hence

$$(44) \quad \mathbf{U} = \mathbf{A}_R \mathbf{X}.$$

If (40) is substituted into (44) the components of \mathbf{U} are

$$(45) \quad U_1 = a_1 p_2 X_1,$$

$$U_i = -b_{i-1} q_i X_{i-1} + a_i p_{i+1} X_i \quad \text{for } i = 2, \dots, n-1$$

and

$$U_n = -b_{n-1} q_n X_{n-1} + a_n X_n.$$

Then by gradual substitutions X_i from i th into $(i+1)$ th equation (45), $i = 1, \dots, n-1$,

$$(46) \quad X_1 = \frac{1}{a_1 p_2} U_1,$$

$$X_i = \sum_{j=1}^{i-1} \frac{b_{i-1} q_i \dots b_j q_{j+1}}{a_i p_{i+1} \dots a_j p_{j+1}} U_j + \frac{1}{a_i p_{i+1}} U_i, \quad i = 2, \dots, n-1,$$

and

$$X_n = \sum_{j=1}^{n-1} \frac{b_{n-1} q_n \dots b_j q_{j+1}}{a_n a_{n-1} p_n \dots a_j p_{j+1}} U_j + \frac{1}{a_n} U_n.$$

Now using \mathbf{B}_R given by (41) in (43) we get

$$(47) \quad Y_i = b_i p_{i+1} X_i \quad \text{for } i = 1, \dots, n-1, \quad \text{and} \\ Y_n = b_n X_n$$

If the equations (46) are substituted into (47) the desired form $\mathbf{Y} = \mathbf{G}\mathbf{U}$ with \mathbf{G} given by (5) and (6) is obtained. Thus (40) and (41) is a right matrix fraction description of \mathbf{G} .

and

$$\mathbf{B}_R = \begin{bmatrix} b_1 p_2 & & & & & \\ 0 & \ddots & & & & \\ & & b_{i-2} p_{i-1} & & & \\ & & & b_{i-1} p_{ii} & & \\ & & & z & b_i p_{i+1} & \\ & & & & & \ddots & \\ & & & & & & b_{n-1} p_n \\ 0 & & & & & & & b_n \end{bmatrix} \quad \mathbf{D} = \mathbf{B}_{RO} \mathbf{D}$$

with the GCRD

$$(53) \quad \mathbf{D} = \begin{bmatrix} 1 & & & & & & & & \\ 0 & \ddots & & & & & & & \\ & & 1 & & & & & & \\ & & & (p_i, b_i) & & & & & \\ & & & x & 1 & & & & \\ & & & & & 0 & \ddots & & \\ & 0 & & & & & & & 1 \\ & & & & & & & & & 0 \\ & & & & & & & & & & 1 \end{bmatrix}$$

The polynomials x, y represent a solution of the equation

$$(54) \quad a_i p_{i+1} x + (p_i, b_i) y = -b_{i-1} q_i$$

which is always solvable since $(p_{i+1}, b_i) \sim 1$ as well as $(a_i, b_i) \sim 1$ is assumed.

Then

$$(55) \quad \begin{aligned} v &= b_{ii} q_{i+1} x, \text{ and} \\ z &= -b_{ii} p_{i+1} x. \end{aligned}$$

For $i = n$ we must put $p_{i+1} = 1$ and $q_{i+1} = 0$ in (54) and (55). The polynomial v in \mathbf{A}_{RO} is omitted.

According to (53) $d = \det \mathbf{D} = (p_i, b_i) \sim 1$ and hence the matrices (40) and (41) are not right coprime.

2. If (42) are broken by $(p_i, b_{i-1}) \sim 1$ for one $i \in [2, n]$ then using the denotations (32) the matrices (40) and (41) can be written in the form

$$\mathbf{A}_R = \begin{bmatrix} a_1 p_2 & & & & & & & & \\ -b_1 q_2 & & & & & & & & 0 \\ & \ddots & & & & & & & \\ & & a_{i-2} p_{i-1} & & & & & & \\ & & -b_{i-2} q_{i-1} & a_{i-1} p_{i,i-1} & & & & & \\ & & -b_{i-1,i} q_i & a_i p_{i+1} & & & & & \\ & & & -b_i q_{i+1} & & & & & \\ & & & & & \ddots & & & \\ 0 & & & & & & a_{n-1} p_n & & \\ & & & & & & & & & -b_{n-1} q_n & a_n \end{bmatrix} \quad \mathbf{D} = \mathbf{A}_{RO} \mathbf{D}$$

and

$$(57) \quad \mathbf{B}_R = \text{diag} [b_1 p_2; \dots; b_{i-2} p_{i-1}; b_{i-1} p_{i,i-1}; b_i p_{i+1}; \dots; b_{n-1} p_n; b_n] \mathbf{D} = \\ = \mathbf{B}_{RO} \mathbf{D}$$

where

$$(58) \quad \mathbf{D} = \text{diag} [1; \dots; 1; (p_i, b_{i-1}); 1; \dots; 1]$$

Since $d = \det \mathbf{D} = (p_i, b_{i-1}) \sim 1$ the matrices (40) and (41) are not right coprime. \square

Note. Polynomial fraction fashions of $P_i(s)$ as well as $T_i(s)$ have been assumed so far. If dead times τ_i are present in a cascade system factors $\exp(-\tau_i s)$ occur in the numerators $b_i(s)$ and/or $q_i(s)$.

But the results of both Theorems 1 and 2 are valid for dead-time systems too, of course, with $\mathbf{A}_L, \mathbf{B}_L$ and $\mathbf{A}_R, \mathbf{B}_R$ being not necessarily polynomial matrices. The algebraic structure containing the elements $t_i(s) \exp(-\tau_i s)$ instead of polynomials can be considered where $t_i(s)$ is a polynomial and $\tau_i \geq 0$. Then going through the steps of presented proofs no special presumptions and operations in this structure are required. It can be appreciated from the physical system viewpoint that $\exp(-\tau_i s)$ and an arbitrary polynomial $c(s)$ have only the common factor equivalent to unit, i.e., $(\exp(-\tau_i s), c) \sim 1$.

4. EXAMPLE

Let the three-stage cascade process (according to Fig. 1) be described by

$$P_1(s) = \frac{b_1(s)}{a_1(s)} = \frac{10}{s(s+2)}, \quad P_2(s) = \frac{b_2(s)}{a_1(s)} = \frac{1}{s+2}, \quad P_3(s) = \frac{b_3(s)}{a_3(s)} = \frac{s+2}{s(s+1)}, \\ T_2(s) = \frac{q_2(s)}{p_2(s)} = \frac{\exp(-s)}{s+0.1}, \quad T_3(s) = \frac{q_3(s)}{p_3(s)} = \frac{1}{s+1}$$

and hence its TFM

$$\mathbf{G}(s) = \begin{bmatrix} \frac{10}{s(s+2)} & 0 & 0 \\ \frac{10 \exp(-s)}{s(s+2)^2(s+0.1)} & \frac{1}{s+2} & 0 \\ \frac{10 \exp(-s)}{s^2(s+2)(s+0.1)(s+1)^2} & \frac{1}{s(s+1)^2} & \frac{s+2}{s(s+1)} \end{bmatrix}$$

Observing that $(p_2, b_2 b_1) \sim 1$ as well as $(p_3, b_3 b_2) \sim 1$ the conditions (12) = (42) are fulfilled.

Using (10) and (11) the LCMF representation of \mathbf{G} is given by

$$\mathbf{A}_L = \begin{bmatrix} s(s+2) & 0 & 0 \\ -\exp(-s) & (s+2)(s+0.1) & 0 \\ 0 & -(s+2) & s(s+1)^2 \end{bmatrix},$$

$$\mathbf{B}_L = \begin{bmatrix} 10 & 0 & 0 \\ 0 & s+0.1 & 0 \\ 0 & 0 & (s+2)(s+1) \end{bmatrix}$$

and according to (40) and (41) the RCMF description of \mathbf{G} is formed by

$$\mathbf{A}_R = \begin{bmatrix} s(s+2)(s+0.1) & 0 & 0 \\ -10 \exp(-s) & (s+2)(s+1) & 0 \\ 0 & -1 & s(s+1) \end{bmatrix},$$

$$\mathbf{B}_R = \begin{bmatrix} 10(s+0.1) & 0 & 0 \\ 0 & s+1 & 0 \\ 0 & 0 & s+2 \end{bmatrix}.$$

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Ing. Václav Soukup, CSc., katedra řídicí techniky elektrotechnické fakulty ČVUT (Department of Automatic Control, Faculty of Electrical Engineering — Czech Technical University), Karlovo nám. 13, 121 35 Praha 2. Czechoslovakia.