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A general upper bound in extremal theory of sequences

MARTIN KLAZAR

Abstract. We investigate the extremal function $f(u, n)$ which, for a given finite sequence u over k symbols, is defined as the maximum length m of a sequence $v = a_1a_2\dots a_m$ of integers such that 1) $1 \leq a_i \leq n$, 2) $a_i = a_j, i \neq j$ implies $|i - j| \geq k$ and 3) v contains no subsequence of the type u . We prove that $f(u, n)$ is very near to be linear in n for any fixed u of length greater than 4, namely that

$$f(u, n) = O(n2^{O(\alpha(n)^{|u|-4})}).$$

Here $|u|$ is the length of u and $\alpha(n)$ is the inverse to the Ackermann function and goes to infinity very slowly. This result extends the estimates in [S] and [ASS] which treat the case $u = abababa\dots$ and is achieved by similar methods.

Keywords: sequence, Davenport-Schinzel sequence, length, upper bound

Classification: 05D99

INTRODUCTION

In the Extremal theory of sequences we investigate the quantity

$$f(u, n) = \max\{|v| \mid u \not\leq v, \|v\| \leq n, v \text{ is } \|u\|\text{-regular}\}.$$

Here u and v are finite sequences of arbitrary symbols, n is a nonnegative integer, $|v|$ stands for the length of v and $\|v\|$ denotes the cardinality of $S(v)$, the set of all symbols that occur in v . If there is a subsequence s in v such that s differs from u only in the names of the symbols we write $u \leq v$ and say that v contains u . For instance $v_1 = 123245131$ contains both $u_1 = xxyy$ and $u_2 = ababa$. A sequence $u = a_1a_2\dots a_m$ is called k -regular if $a_i = a_j, i \neq j$ implies $|i - j| \geq k$. Example: v_1 and u_2 are 2-regular but are not 3-regular and u_1 is not 2-regular. If $u = a_1a_2\dots a_m$ and $a_i = a \in S(u)$ then we shall refer to a_i as to the a -letter.

The function $f(u, n)$ extends in a natural way the function $F = f(ababa, n)$ investigated at first by Davenport and Schinzel in [DS]. They proved the upper bound $F = O(n \log n / \log \log n)$ that was later improved by Szemerédi to $O(n \log^* n)$ ([Sz]). Here $\log^* n$ is the minimum number of iterations of the power function 2^m (starting with $m = 1$) which are needed to get a number greater or equal to n . The question whether $F = O(n)$ ($f(abab, n) = 2n - 1$ trivially) remained open until 1986 when it was answered by Hart and Sharir in [HS] negatively. They showed

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that $F = \Theta(n\alpha(n))$ where $\alpha(n)$ goes to infinity but very slowly (a precise definition of $\alpha(n)$ will be given in the second part of this paper). M. Sharir obtained later

$$f(al_s, n) = O(n\alpha(n)^{O(\alpha(n)^{s-5})})$$

for arbitrary alternating sequence $al_s = ababab \dots$ of the length $s \geq 5$ ([S]). Recently almost tight estimates were derived ([ASS]):

$$\begin{aligned} f(al_s, n) &\leq n \cdot 2^{(\alpha(n))^{\frac{s-5}{2}} \log_2 \alpha(n) + C_s(n)} && \text{for } s \geq 5 \text{ odd} \\ f(al_s, n) &\leq n \cdot 2^{(\alpha(n))^{\frac{s-4}{2}} + C_s(n)} && \text{for } s \geq 6 \text{ even} \\ f(al_s, n) &= \Omega(n \cdot 2^{K_s \cdot (\alpha(n))^{\frac{s-4}{2}} + Q_s(n)}) && \text{for } s \geq 6 \text{ even} \end{aligned}$$

where $K_s = \frac{1}{(\frac{s-4}{2})!}$ and $C_s(n)$ and $Q_s(n)$ are asymptotically smaller than the main terms. For $s = 6$ even, $f(ababab, n) = \Theta(n2^{\alpha(n)})$ ([ASS]). How complex the previous formulae may seem on the first view, one thing is clear: $f(al_s, n)$ is almost almost linear in n for all s .

The first aim of this paper is to show that the same is true for arbitrary sequence u . The second aim is to give a brief and clear idea about the techniques developed by Agarwal, Hart, Sharir and Shor for obtaining almost linear upper bounds on $f(al_s, n)$ to the reader that is not familiar with them.

In the first part we show a simple method that leads to the upper bound $f(u, n) = O(n^2)$ for all u . Then, in the second part, we use a slightly generalized method of [S] to derive the estimate

$$f(u, n) = O(n \cdot 2^{O(\alpha(n)^{|u|-4})}).$$

PART 1

We first define a modification of the function $f(u, n)$ for l -regular sequences:

$$f(u, n, l) = \max\{|v| \mid u \not\leq v, \|v\| \leq n, v \text{ is } l\text{-regular}\}$$

where $l \geq \|u\|$.

Lemma 1.1. a) $f(u, n, l)$ is defined and finite for any $n \geq 1$ and moreover $f(u, n, l) = O(|u| \cdot \|u\| \cdot n^{\|u\|})$.

b) $f(u, n, l) \leq f(u, n, k) \leq (1 + f(u, l - 1, k))f(u, n, l)$ for all $l > k \geq \|u\|, n \geq 1$.

PROOF: ad a) We suppose there is at least one repetition in u , otherwise the function $f(u, n, l)$ is constant. If $n < l$ then $f(u, n, l) = n$. If $n \geq l$ then any l -regular sequence v satisfying $|v| \geq \|u\| \cdot \binom{n}{\|u\|} (|u| - 1) + 1$ must contain u . We split $v = v_1 v_2 \dots v_c w$ so that $|v_i| = \|v_i\| = \|u\|$ and $c = (|u| - 1) \binom{n}{\|u\|} + 1$. According to the Dirichlet Principle there exist $|u|$ indices $1 \leq i_1 < i_2 < \dots < i_{|u|} \leq c$ that $S(v_{i_1}) = S(v_{i_2}) = \dots = S(v_{i_{|u|}})$. Thus $u \leq v_{i_1} v_{i_2} \dots v_{i_{|u|}}$.

ad b) The first inequality is obvious. Suppose $v = a_1 a_2 \dots a_m$ is k -regular, does not contain u and $\|v\| \leq n$. We choose a subsequence v^* of v in this way: we start with $v^* = a_1$ and $i = 1$ and search for the minimum j such that $j > i$ and $v^* a_j$ is l -regular. If such a j exists then we put $v^* = v^* a_j$ and $i = j$ and repeat. Otherwise the algorithm terminates. Obviously $\|v^*\| \leq n$ and v^* is l -regular. Moreover $|v| \leq (1 + f(u, l - 1, k))|v^*|$ because any interval I in v omitted by the previous algorithm satisfies $\|I\| \leq l - 1$. We got the second inequality. \square

Definition 1.2. Let u, v be sequences. We write $u \leq\leq v$ if $u \leq v^*$ for all v^* obtained from v by restricting v to some $\|u\|$ symbols. Thus in this case v contains u in all possible ways.

Lemma 1.3. For any sequence u there exist positive integers m and s such that $u \leq v$ whenever $\|v\| \geq m$ and $al_s \leq\leq v$.

Before proving this lemma we derive the main result of this section.

Theorem 1.4. $f(u, n) = O(n^2)$ for all sequences u . The constant in O depends on u .

PROOF: Let $m = m(u)$ be as in Lemma 1.3. According to Lemma 1.1 b) we have $f(u, n) = f(u, n, \|u\|) \leq (1 + f(u, m - 1, \|u\|))f(u, n, m)$. We estimate $f(u, n, m)$. Suppose v is m -regular, $\|v\| \leq n$, $u \not\leq v$ and $|v| = f(u, n, m)$. It suffices to estimate the number c in the splitting $v = v_1 v_2 \dots v_c w$ where $|v_i| = \|v_i\| = m$ and $|w| \leq m - 1$. Let $s = s(u)$ stand for the second number of Lemma 1.3. For any v_i there exist symbols $a, b \in S(v_i)$ such that v restricted on the symbols $\{a, b\}$ does not contain al_s . Otherwise $u \leq v$ according to Lemma 1.3. But the mapping $F : \{v_1, v_2, \dots, v_c\} \rightarrow \binom{S(v)}{2}$ that maps any v_i on a pair $\{a, b\}$ mentioned above maps only at most $s - 2$ v_i 's on one pair because of the property of the symbols $\{a, b\}$. Thus $c \leq (s - 2)\binom{n}{2}$. Finally

$$f(u, n) \leq (1 + f(u, m - 1, \|u\|))m(c + 1) \leq (1 + f(u, m - 1, \|u\|))m(1 + (s - 2)\binom{n}{2}).$$

Thus

$$f(u, n) = O(n^2).$$

\square

It remains to prove Lemma 1.3. We use the following well known:

Lemma 1.5 (Erdős P., Szekeres G. 1935 [ES]). Any $(n - 1)^2 + 1$ -term sequence (of integers) contains a n -term monotone subsequence.

PROOF OF LEMMA 1.3: We denote by $X(k, l)$ the set of all sequences of the form $y_1 y_2 \dots y_l$ where $y_i = x_1 x_2 \dots x_k$ or $y_i = x_k x_{k-1} \dots x_1$ for k distinct symbols x_1, x_2, \dots, x_k . Thus $|X(k, l)| = 2^l$ and $|u| = kl$ and $\|u\| = k$ for any $u \in X(k, l)$. Since $u \leq w$ for any $w \in X(\|u\|, |u|)$, it suffices to prove the following claim.

Claim. For all positive integers k and l there exist positive integers m and s such that $w \leq v$ for some $w \in X(k, l)$ whenever $\|v\| \geq m$ and $al_s \leq v$.

PROOF OF THE CLAIM: We put $s = 2l$ and $m = k_1$ where $k_l = k$ and $k_{t-1} = 4(t-1)k_t^2 + 3$ for $t = l, l-1, \dots, 2$. Suppose v meets the prescribed conditions. We prove by induction that for all $t = 1, 2, \dots, l$ there exists $w \in X(k_t, t)$ such that $w \leq v$. For $t = 1$ this is obvious. Suppose it is true for $t-1 \geq 1$. We have $w \in X(k_{t-1}, t-1), w \leq v$. We take a fixed w -copy U in v and split v into $t-1$ intervals $v = v_1 v_2 \dots v_{t-1}$ where v_i contains i -th part of U (i.e. y_i). U consists of $k_{t-1}(t-1)$ letters $x_i^j, j = 1 \dots t-1, i = 1 \dots k_{t-1}$ in v, x_i^j occur in $v_j, a < b$ implies that x_a^j precedes x_b^j and $x_1^p = x_1^q, x_2^p = x_2^q, \dots$ or $x_1^p = x_{k_{t-1}}^q, x_2^p = x_{k_{t-1}-1}^q, \dots$ for all p, q . It remains to give names to the symbols — say that x_i^j is $z(i)$ -letter for $i = 1, 2, \dots, k_{t-1}$. There must be other $z(i)$ -letters in v besides those in U ($al_s \leq v$). Let us consider the pairs of symbols $(z(1), z(k_{t-1})), (z(2), z(k_{t-1}-1)), \dots, (z(L), z(k_{t-1}-L+1)), L = \lceil k_{t-1}/2 \rceil - 1$. The Dirichlet Principle implies that there are a set $M \subset \{1, 2, \dots, L\}, |M| \geq \frac{L}{t-1}$ and an index $r \in \{1, 2, \dots, t-1\}$ that $z(i)z(k_{t-1}-i+1)z(i)$ or $z(k_{t-1}-i+1)z(i)z(k_{t-1}-i+1)$ is a 3-term subsequence of v_r for any $i \in M$. We used that $al_s \leq v$ and $s > 2(t-1)$. We can suppose w.l.o.g. $r = 1$. Thus we have 2-term subsequence $z(k_{t-1}-i+1)z(i)$ of v_1 for any $i \in M$ (the opposite order than in U). The $z(L+1)$ -letter x_{L+1}^1 (lies in U) splits v_1 on two intervals $v_1 = v_1' v_1''$. There are at least $|M|/2$ i 's in M such that $z(i)$ -letter occurs in v_1'' or there are $|M|/2$ i 's in M such that $z(k_{t-1}-i+1)$ -letter occurs in v_1' . We obtained t separated areas — namely $v_1', v_1'', v_2, \dots, v_{t-1}$ — in which $z(i)$ -letter occurs for at least $|M|/2$ i 's. From those at least $|M|/2$ i 's we choose according to Lemma 1.5 at least $\sqrt{|M|/2}$ i 's in such a way that we obtain a w' -copy in $v, w' \in X(\lceil \sqrt{|M|/2} \rceil, t)$. We are finished because $\lceil \sqrt{|M|/2} \rceil \geq \lceil \sqrt{L/2(t-1)} \rceil \geq \dots \geq k_t$. \square

Remark 1.6. If we estimate $k_{t-1} = 4(t-1)k_t^2 + 3 \leq t(2k_t)^2$ then it may be easily derived that it suffices to put in Lemma 1.3 $s = 2|u|, m = (4|u| \cdot \|u\|)^{2|u|-1}$.

PART 2

In this section we prove a result far stronger than $f(u, n) = O(n^2)$. At first we give the precise (standard) definition of $\alpha(n)$.

For any function $B : \mathbf{N} \rightarrow \mathbf{N}$ the symbol $B^{(s)}(n)$ denotes $B(B(\dots(B(n))\dots))$ (s times). We define further the functional inverse of B as $B^{-1}(n) = \min\{s \geq 1 \mid B(s) \geq n\}$. For nondecreasing and unbounded B the functional inverse B^{-1} is nondecreasing and unbounded as well. The functions $A_k(n)$ are defined by induction:

$$A_k(1) = 2, A_1(n) = 2n \text{ and } A_k(n) = A_{k-1}^{(n)}(1).$$

Thus $A_2(n) = 2^n, A_3(n) = 2^{2^{\dots^2}}$ n times. The Ackermann function is diagonal function of that schema: $A(n) = A_n(n)$. The function $\alpha(n)$ is defined as $\alpha(n) =$

$A^{-1}(n)$. Apart the hierarchy A_1, A_2, \dots (A_{i+1} grows to infinity much faster than A_i), we have the hierarchy $\alpha_1, \alpha_2, \dots, \alpha_i = A_i^{-1}$ (α_{i+1} grows to infinity much more slowly than α_i). Thus $\alpha_1(n) = \lceil \frac{n}{2} \rceil, \alpha_2(n) = \lceil \log_2 n \rceil, \alpha_3(n) = \log^*(n), \dots$. The function α is far “lazier” than any α_i . It is easy to prove for α_i a recurrent formula $\alpha_{i+1}(n) = \min\{s \geq 1 \mid \alpha_i^{(s)}(n) = 1\}$. Thus

$$(1) \quad \alpha_{i+1}(\alpha_i(m)) = \alpha_{i+1}(m) - 1 \quad \text{for all } i \geq 1, m \geq 3.$$

Further ([ASS])

$$(2) \quad \alpha_{\alpha(n)+1}(n) \leq 4 \quad \text{for all } n \geq 1.$$

A sequence u is called a 1-chain if no symbol occurs repeatedly in u . $Y(k, l)$ denotes the set of all sequences of the form $y_1 y_2 \dots y_l$ where any y_i is a permutation of k fixed symbols x_1, x_2, \dots, x_k . $Y(k, l) \not\leq v$ means that $u \leq v$ for no $u \in Y(k, l)$. We modify a bit the function $\Psi_s(m, n)$ of [S] and introduce the function

$$\Psi_r^s(m, n) = \max\{|v| \mid v \text{ is } r\text{-regular}, \|v\| \leq n, v = v_1 v_2 \dots v_m \\ \text{where any } v_i \text{ is 1-chain and } Y(r, s) \not\leq v\}.$$

We will estimate $f(u, n)$ in four steps. We will proceed induction on s . At first we estimate $\Psi_r^3(m, n)$. Then we derive, supposing we have an upper bound on $\Psi_r^{s-1}(m, n)$, a recurrent inequality for $\Psi_r^s(m, n)$. In the third step using that inequality the upper bound considered in Step 2 is extended on $\Psi_r^s(m, n)$. Finally we estimate $f(u, n)$ by appropriate $\Psi_r^s(m, n)$.

Step 1.

Lemma 2.1. $\Psi_r^3(m, n) \leq 2rn$.

PROOF: Suppose v is r -regular, $\|v\| \leq n$ and $Y(r, 3) \not\leq v$ (we ignore here the first variable in Ψ). We split $v = v_1 v_2 \dots v_c w$ where $|v_i| = \|v_i\| = r$ and $|w| < r$. Any v_i must contain the first letter or the last letter of some symbol (otherwise $u \leq v$ for some $u \in Y(r, 3)$). Thus

$$|v| = cr + |w| \leq (2\|v\| - |w|)r + |w| \leq 2rn.$$

□

Step 2.

Lemma 2.2. Suppose $\Psi_r^{s-1}(m, n) \leq F_{s-1}(m)m + G_{s-1}(m)n$ for $m, n \geq 1$ for some nondecreasing functions $F_{s-1}, G_{s-1} : \mathbf{N} \rightarrow \mathbf{N}$. Then for any partition $m = m_1 + \dots + m_b, m_i \geq 1, 1 < b < m$ there exists a partition $n = n_0 + n_1 + \dots + n_b, n_i \geq 0$ such that

$$(3) \quad \Psi_r^s(m, n) \leq \sum_{i=1}^b \Psi_r^s(m_i, n_i) + 2\Psi_r^s(b, n_0)G_{s-1}(m) + mH_{s-1}(m)$$

where $H_{s-1}(m) = 3(r - 1) + 2F_{s-1}(m) + 2(r - 1)G_{s-1}(m)$.

PROOF: We start with a preliminary consideration. Suppose an r -regular sequence u is splitted into o 1-chains $u = u_1 u_2 \dots u_o$. Then a subsequence v of u need not be r -regular but it suffices to delete at most $(r - 1)(o - 1)$ letters from v and what remains is r -regular. This consideration will be used in this proof and then again in the fourth step.

Let v be r -regular, $\|v\| \leq n$, $Y(r, s) \not\leq v$, v consists of m 1-chains and $|v| = \Psi_r^s(m, n)$. We group 1-chains of v in b layers (the partition $m = m_1 + \dots + m_b$ is given) L_1, L_2, \dots, L_b where L_i consists of m_i 1-chains. Thus $v = L_1 L_2 \dots L_b$. We split any L_i in three subsequences v_i^1, v_i^2 and v_i^3 , v_i^1 consists of those letters that occur only in L_i (i.e. $S(v_i^1) \cap S(L_j) = \emptyset$ for $i \neq j$), v_i^2 consists of those that occur also before L_i and v_i^3 consists of the remaining ones (i.e. do not occur before L_i but occur after L_i). Obviously

$$(4) \quad \Psi_r^s(m, n) = |v| = \sum_{i=1}^b |v_i^1| + \sum_{i=1}^b |v_i^2| + \sum_{i=1}^b |v_i^3|.$$

The upper bound on the first term in (4) is clearly

$$\sum_{i=1}^b (\Psi_r^s(m_i, n_i) + (m_i - 1)(r - 1)) = \sum_{i=1}^b \Psi_r^s(m_i, n_i) + (m - b)(r - 1)$$

where $n_i = \|v_i^1\|$. We come naturally to the partition $n = n_0 + n_1 + \dots + n_b$, n_0 is the number of all symbols figurating in all v_i^2, v_i^3 . Observe that $Y(r, s - 1) \not\leq v_i^2, v_i^3$ for all i . This fact enables us to estimate the remaining two terms in (4). We do it only for the second one, the third one is treated similarly. According to the hypothesis

$$\begin{aligned} \sum_{i=1}^b |v_i^2| &\leq \sum_{i=1}^b (F_{s-1}(m_i)m_i + G_{s-1}(m_i)\|v_i^2\| + (m_i - 1)(r - 1)) \leq \\ &\leq F_{s-1}(m)m + G_{s-1}(m)\sum_{i=1}^b \|v_i^2\| + (m - b)(r - 1). \end{aligned}$$

We transform any v_i^2 to w_i by taking any $a \in S(v_i^2)$ just once (the 1-chain w_i is a subsequence of v_i^2). The sequence $w = w_1 w_2 \dots w_b$ meets (after deleting at most $(b - 1)(r - 1)$ letters) all conditions to be estimated by $\Psi_r^s(b, n_0)$. Thus

$$\sum_{i=1}^b \|v_i^2\| = |w| \leq \Psi_r^s(b, n_0) + (b - 1)(r - 1).$$

We substitute all derived bounds in (4):

$$\begin{aligned} \Psi_r^s(m, n) &\leq \sum_{i=1}^b \Psi_r^s(m_i, n_i) + (m - b)(r - 1) + \\ &\quad + 2[F_{s-1}(m)m + G_{s-1}(m)(\Psi_r^s(b, n_0) + (b - 1)(r - 1)) + (m - b)(r - 1)]. \end{aligned}$$

We got (3). □

Step 3.

Lemma 2.3. *Let F_{s-1}, G_{s-1} and H_{s-1} be as in Lemma 2.2. Then for any $m, n \geq 1, k \geq 2$*

$$(5) \quad \Psi_r^s(m, n) \leq \alpha_k(m)m.H_{s-1}(m).(5G_{s-1}(m))^{k-2} + 2n.(2G_{s-1}(m))^{k-1}.$$

PROOF: For $m \leq 4$ (5) holds because of the trivial inequality $\Psi_r^s(m, n) \leq mn$. We prove (5) induction on k , for k fixed induction on m . We start with $k = 2$. It suffices to verify induction on m the estimate

$$\Psi_r^s(m, n) \leq H_{s-1}(m)\lceil \log_2 m \rceil m + 4G_{s-1}(m)n$$

((5) for $k = 2$) using the inequality

$$\Psi_r^s(m, n) \leq \Psi_r^s(\lfloor \frac{m}{2} \rfloor, n_1) + \Psi_r^s(\lceil \frac{m}{2} \rceil, n_2) + 4G_{s-1}(m)n_0 + mH_{s-1}(m)$$

((3) for $b = 2$). It is left to the reader.

In case $k > 2, m \geq 3$ we put in (3) $b = \lceil \frac{m}{\alpha_{k-1}(m)} \rceil, m_i \leq \lceil \frac{m}{b} \rceil \leq \alpha_{k-1}(m)$. Thus $\alpha_k(m_i) \leq \alpha_k(m) - 1$ (according to (1)) and $b\alpha_{k-1}(b) \leq b\alpha_{k-1}(m) \leq 2m$. We estimate the term $\Psi_r^s(m_i, n_i)$ in (3) by (5) for k, m_i , and the term $\Psi_r^s(b, n_0)$ by (5) for $k - 1, b$. Then

$$\begin{aligned} \Psi_r^s(m, n) &\leq \sum_{i=1}^b (H_{s-1}(m_i)(5G_{s-1}(m_i))^{k-2}\alpha_k(m_i)m_i + 2(2G_{s-1}(m_i))^{k-1}n_i) + \\ &+ (H_{s-1}(b)(5G_{s-1}(b))^{k-3}\alpha_{k-1}(b)b + 2(2G_{s-1}(b))^{k-2}n_0)2G_{s-1}(m) + mH_{s-1}(m) \leq \\ &\leq H_{s-1}(m)(5G_{s-1}(m))^{k-2}(\alpha_k(m) - 1)m + 2(2G_{s-1}(m))^{k-1}(n - n_0) + \\ &+ H_{s-1}(m)((5G_{s-1}(m))^{k-2} - 1)m + 2(2G_{s-1}(m))^{k-1}n_0 + mH_{s-1}(m) \leq \\ &\leq H_{s-1}(m)(5G_{s-1}(m))^{k-2}\alpha_k(m)m + 2(2G_{s-1}(m))^{k-1}n. \end{aligned}$$

□

Lemma 2.4. *For any $s \geq 4$ the inequality*

$$(6) \quad \Psi_r^s(m, n) \leq m(10r)^{\alpha^{s-3}(m)+4\alpha^{s-4}(m)} + n(4r)^{\alpha^{s-3}(m)+2\alpha^{s-4}(m)} \quad m, n \geq 1$$

holds.

PROOF: We consider the functions $\overline{F}_s, \overline{G}_s, s \geq 3$ that are defined by the following recurrent relations (we write \overline{F}_s instead $\overline{F}_s(m), \overline{G}_s$ instead $\overline{G}_s(m)$ and α instead of $\alpha(m)$ for the sake of brevity):

$$\begin{aligned} \overline{F}_3 &= 0, \overline{G}_3 = 2r \\ \overline{F}_s &= 4(3(r - 1) + 2\overline{F}_{s-1} + 2(r - 1)\overline{G}_{s-1})(5\overline{G}_{s-1})^{\alpha-1}, \overline{G}_s = 2(2\overline{G}_{s-1})^\alpha. \end{aligned}$$

Induction on s shows that

$$\Psi_r^s(m, n) \leq \overline{F}_s(m)m + \overline{G}_s(m)n$$

for any $m, n \geq 1, s \geq 3$. Indeed, for $s = 3$ it follows from Step 1 and for general s we obtain this inequality from (5) where we put $k = \alpha(m) + 1$ and use (2). We count explicit upper bounds on both functions. Clearly $\overline{G}_s = 2.4^{\alpha^{s-4} + \alpha^{s-5} + \dots + \alpha} \cdot (4r)^{\alpha^{s-3}}$ for $s \geq 5$ and $\overline{G}_4 = 2(4r)^\alpha$. Hence $\overline{G}_s \leq (4r)^{\alpha^{s-3} + 2\alpha^{s-4}}$ for $s \geq 4$.

Further $\overline{F}_4 = \frac{2}{5}(4r-1-\frac{3}{r})(10r)^\alpha \geq \overline{G}_4$ and therefore $\overline{F}_s \geq \overline{G}_s$ for all $s \geq 4$. Thus $\overline{F}_s \leq 4(3(r-1) + 2r\overline{F}_{s-1})(5\overline{F}_{s-1})^{\alpha-1} \leq 4r(5\overline{F}_{s-1})^\alpha$. If we solve this recurrent relation as an equation then an upper bound on \overline{F}_s is obtained. We start with $\overline{F}_4 \leq 2r(10r)^\alpha$ and derive

$$\overline{F}_s \leq (2r)^{\alpha^{s-4}} \cdot (4r)^{\alpha^{s-5} + \dots + 1} \cdot 5^{\alpha^{s-4} + \dots + \alpha} \cdot (10r)^{\alpha^{s-3}} \leq (10r)^{\alpha^{s-3} + 4\alpha^{s-4}}. \quad \square$$

Step 4.

Lemma 2.5.

$$(7) \quad f(u, n) \leq 2\|u\| \cdot 2^{|u|-4} \cdot n \cdot (10\|u\|)^{2\alpha^{|u|-4}(n) + 8\alpha^{|u|-5}(n)}$$

for any sequence $u, |u| \geq 5$.

PROOF: We will find the upper bound $nE_s(n)$ ($E_s(n)$ is a nondecreasing function) on the quantity

$$\max\{|v| \mid v \text{ is } r\text{-regular, } \|v\| \leq n, Y(r, s) \not\leq v\}.$$

It suffices because $u \leq v$ for any $v \in Y(\|u\|, |u| - 1)$ except $u = aa \dots a$ (i times) but $f(aa \dots a, n) = n(i - 1)$. We derive for E_s a recurrent relation. Let v be r -regular, $\|v\| \leq n$ and $Y(r, s) \not\leq v$. We split $v = v_1 l_1 v_2 l_2 \dots v_n l_n$ where l_1, \dots, l_n are the last letters of all $x \in S(v)$. Observe that $Y(r, s - 1) \not\leq v_i$ and hence $|v| = \sum_{i=1}^n |v_i| + n \leq (\sum_{i=1}^n \|v_i\|)E_{s-1}(n) + n$. The sum $\sum_{i=1}^n \|v_i\|$ may be estimated by $\Psi_r^s(n, n) + (n - 1)(r - 1)$ (we use the same trick as in Lemma 2.2 — replace v_i by 1-chain of the length $\|v_i\|$). Thus

$$|v| \leq 2nE_{s-1}(n) \cdot (10r)^{\alpha^{s-3}(n) + 4\alpha^{s-4}(n)}$$

by (6). Hence we may choose

$$E_3(n) = 2r \text{ (see Step 1)}$$

$$E_s(n) = 2E_{s-1}(n) \cdot (10r)^{\alpha^{s-3}(n) + 4\alpha^{s-4}(n)}.$$

The solution of this relation is:

$$E_s(n) = 2r \cdot 2^{s-3} \cdot (10r)^{\alpha^{s-3}(n) + \alpha^{s-4}(n) + \dots + \alpha(n) + 4\alpha^{s-4}(n) + \dots + 4}.$$

If replaced r by $\|u\|$ and s by $|u| - 1$ then (7) is obtained. □

CONCLUDING REMARKS

We achieved the exponent $\alpha^{|u|-4}(n)$ in (7) by induction starting with $s = 3$. It is possible that this bound might be improved to (roughly) $\alpha^{\frac{1}{2}|u|}(n)$ but it would require computations far more complex as in [ASS].

More interesting than the best value in (7) is perhaps the fact that $f(u, n)$ is almost linear for any sequence u . Hence a double induction must be used in some form whenever we want to obtain a superlinear lower bound on $f(u, n)$ (cf. [HS], [ASS], [K], [FH] and [WS]). Methods giving such “huge” functions as $n^{\frac{7}{6}}$ or $n \log \log n$ or $n \log^* n$ cannot be successful. It is a remarkable difference in comparison with extremal problems concerning graphs or hypergraphs (Turán theory). Here most common functions are n^β , $\beta > 1$. A certain hybrid occurs in Davenport-Schinzel theory of matrices in [FH] where the maximum number of 1's in a 0-1 matrix (of the size $n \times n$) which does not contain a forbidden subconfiguration is investigated. Here $n\alpha(n)$ figurates as an upper bound as well as $n^{\frac{3}{2}}$ and $n \log n$.

For obtaining a good general upper bound on $f(u, n)$ only basic features of u — such as the length and the number of symbols — were important. It is demonstrated by the fact that we worked instead of u itself with the sets $X(k, l)$ resp. $Y(k, l)$ that are determined by $|u|$ and $\|u\|$. It is probable that this changes if we start to investigate finer properties of the asymptotic growth of $f(u, n)$. But except for the case $u = al_s$ where we know the magnitude of $f(u, n)$ with high precision due the deep result of [ASS] only little about that function is known. One of the basic questions is to determine the set

$$Lin = \{u \mid f(u, n) = O(n)\}$$

— see [AKV] and [Kl] for a partial solution.

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