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INDEPENDENT BASIS FOR THE IDENTITIES OF ENTROPIC  
GROUPOIDS  
G. POLLÁK. A. SZENDREI

Abstract: The variety  $E$  of entropic groupoids, which is generated by any of the algebras  $\mathcal{U}_{r,s} = \langle \mathbb{R}; \circ \rangle$  where  $\mathbb{R}$  is the set of real numbers,  $r, s \in \mathbb{R}$  are algebraically independent and  $x \circ y = rx + sy$ , is known to be not finitely based [1]. Here we give an independent basis for the identities of  $E$ .

Key words and phrases: variety, identity, equational theory, basis (of identities), independent basis, entropic groupoid.

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Secondary 20L05

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In [1] Ježek and Kepka describe the equational theory of entropic groupoids. In particular it follows that the algebra  $\mathcal{U} = \langle A; \circ \rangle$  defined on the free commutative ring  $A$  with free generators  $a_0, a_1$  by  $x \circ y = a_0 x + a_1 y$ , generates the variety  $E$  of entropic groupoids. They also show that the equational theory of  $E$  (and hence of  $\mathcal{U}$ ) is not finitely based. Here we construct an independent basis for the equational theory of  $E$ . These investigations concern also a question of Fajtlowicz and Mycielski[2] asking whether the

groupoids  $\mathcal{U}_{r,s} = \langle \mathbb{R}; \circ \rangle$  defined on the set  $\mathbb{R}$  of real numbers by  $x \circ y = rx + sy$  have finite bases for their identities. Clearly, if  $r$  and  $s$  are algebraically independent then  $\mathcal{U}_{r,s}$  generates the variety  $\mathcal{E}$ , hence its equational theory is not finitely based.

We use the terminology and notations of [3]. Since all algebras occurring are groupoids, we omit all references to the type. In particular, for any cardinal  $\beta$ ,  $P^{(\beta)}$  stands for the set of polynomial symbols of type  $\langle 2 \rangle$  with variables  $\{x_\gamma: \gamma < \beta\}$ . Clearly,  $\mathcal{R}^{(\beta)} = \langle P^{(\beta)}; \circ \rangle$  is the free groupoid on  $\beta$  generators. For  $p, p' \in P^{(\beta)}$ ,  $p \equiv p'$  means that  $p$  and  $p'$  coincide.

Let  $R$ ,  $A$  and  $M$  denote the free unitary ring, free unitary commutative ring and free monoid with free generators  $a_0, a_1$ , respectively. (We consider  $M$  to be a subset of  $R$ .) The length of a word  $w \in M$  is denoted by  $|w|$ . Define the entropic groupoids  $\mathcal{R} = \langle \mathbb{R}; \circ \rangle$  and  $\mathcal{U} = \langle A; \circ \rangle$  by  $x \circ y = a_0 x + a_1 y$ . Let  $\alpha: R \rightarrow A$  be the natural ring homomorphism with  $a_i \alpha = a_i$  ( $i < 2$ ). Clearly,  $\alpha$  is also a groupoid homomorphism  $\mathcal{R} \rightarrow \mathcal{U}$ . For any  $i < \omega$  let  $\varphi_i: \mathcal{R}^{(\omega)} \rightarrow \mathcal{R}$  be the natural homomorphism with  $x_i \varphi_i = 1$  and  $x_j \varphi_i = 0$  if  $j \neq i$ . Further, set  $\varphi = \sum_{i < \omega} \varphi_i$ . It is not hard to show that for any  $p, q \in P^{(\omega)}$ ,


(\*)  $p \equiv q$  iff for every  $i < \omega$ ,  $p \varphi_i = q \varphi_i$ ;

(\*\*)  $p$  and  $q$  have the same parenthesis structure, i.e.  $p(x_0, \dots, x_\omega) \equiv q(x_0, \dots, x_\omega)$ , iff  $p\varphi = q\varphi$ .

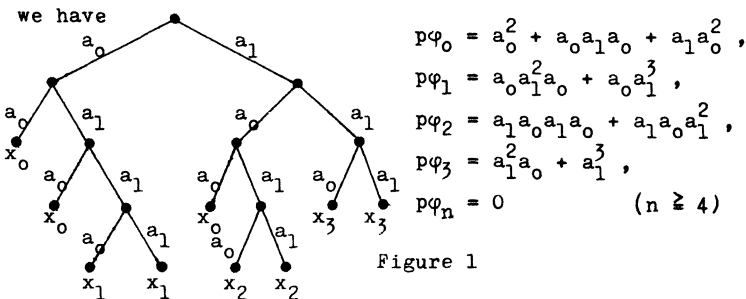
To see this, and also to make it easier to follow the rest of the paper, it is worth noting what the homomorphisms

$\varphi_i$  mean pictorially. There is a natural way to represent a polynomial symbol in  $P^{(\omega)}$  by a binary tree as follows: to  $x_i$  ( $i < \omega$ ) we assign the one-point tree

$$\tilde{x}_i$$

and to any polynomial symbol  $p \cdot q$  we assign the tree arising from  by attaching to its left and right branches the trees corresponding to  $p$  and  $q$ , respectively. Now, consider the tree of a polynomial symbol  $p \in P^{(\omega)}$ , and label all branches going to the left by  $a_0$  and all branches going to the right by  $a_1$ . In this manner, the paths of the tree of  $p$  can be labelled by words from  $M$  and every vertex is uniquely characterized by the word corresponding to the path going downwards to it. This word will be called the weight of the vertex. Since the subterms of  $p$  are in a natural one-to-one correspondence with the vertices of the tree of  $p$ , we can also speak about the weight of a subterm of  $p$ . In particular, the variables are also subterms of  $p$ . Now it is easy to see that for any  $i < \omega$ ,  $p\varphi_i$  is nothing else than the sum of the weights of all occurrences of the variable  $x_i$ .

Example. For  $p = (x_0 \cdot (x_0 \cdot (x_1 \cdot x_1))) \cdot ((x_0 \cdot (x_2 \cdot x_2)) \cdot (x_3 \cdot x_3))$



Clearly, for any  $p \in P^{(\omega)}$  a variable  $x_i$  occurs in  $p$  iff  $p\varphi_i \neq 0$ . Put  $\nu(p) = \{i < \omega : p\varphi_i \neq 0\}$ . For any mapping  $\psi: \nu(p) \rightarrow \{i: i < \omega\}$  we denote by  $p^\psi$  the polynomial symbol arising from  $p$  by substituting  $x_{i\psi}$  for  $x_i$  for all  $i \in \nu(p)$ .

**Proposition 1.** For any  $p, q \in P^{(\omega)}$ , the identity  $p=q$  is in  $\text{Id}(E)$  if and only if  $p\varphi_i^\alpha = q\varphi_i^\alpha$  holds for all  $i < \omega$ .

**Proof:** The statement follows from the fact that for any  $p \in P^{(\omega)}$ ,  $p_\alpha = \sum_{i < \omega} (p\varphi_i^\alpha)x_i$ . The proof is straightforward by induction.

**Example.** Figure 2 shows the tree of a polynomial symbol  $q$  for which  $p=q$  belongs to  $\text{Id}(E)$  ( $p$  is the polynomial symbol in Figure 1).

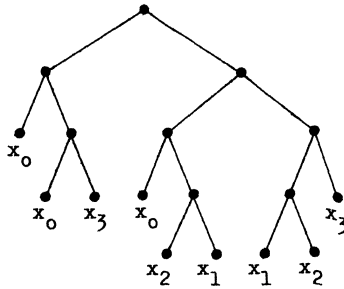


Figure 2

Let  $\bar{P}$  denote the set of all  $p \in P^{(\omega)}$  in which every  $x_i$  ( $i < \omega$ ) occurs at most once; i.e.  $p \in \bar{P}$  iff  $p \in P^{(\omega)}$  and  $p\varphi_i \in M$  for every  $i \in \nu(p)$ . Denote by  $\tilde{P}$  the subset of  $\bar{P}$  consisting of all  $p \in \bar{P}$  such that  $\nu(p) = \{i: i < n\}$  for some  $n < \omega$ , and for every  $i, j \in \nu(p)$ ,  $i > j$  iff either

$|p\varphi_1| < |p\varphi_j|$  or  $|p\varphi_1| = |p\varphi_j|$  and  $p\varphi_j$  precedes  $p\varphi_1$  in the lexicographic order. Pictorially, this means that a polynomial symbol belongs to  $\tilde{\mathcal{F}}$  iff in its tree the variables  $x_0, x_1, x_2, \dots$  are attached to the branches sequentially by levels, starting from the bottom, and within one level from the left to the right (see Figure 3).

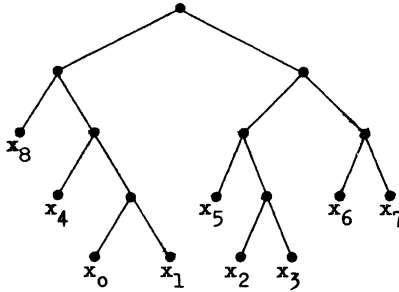


Figure 3

Obviously, for every  $p \in \tilde{\mathcal{F}}$  there is a (unique) one-to-one mapping  $\pi: \nu(p) \rightarrow \{i: i < \omega\}$  such that  $p^\pi \in \tilde{\mathcal{F}}$ . Making use of  $(\ast)$  and  $(\ast\ast)$  it is not hard to see that every polynomial symbol  $p \in \tilde{\mathcal{F}}$  is uniquely determined by  $p\varphi$ .

Proposition 2. If  $p=q$  ( $p, q \in P^{(\omega)}$ ) is in  $\text{Id}(E)$  then there exist  $p' \in \tilde{\mathcal{F}}$  and  $q' \in \tilde{\mathcal{F}}$  such that  $p' = q'$  is also in  $\text{Id}(E)$  and  $p' = q' \vdash p = q$ .

Example. Let  $p$  and  $q$  be the polynomial symbols in Figures 1 and 2, respectively. Then  $p=q$  is in  $\text{Id}(E)$  and the polynomial symbol  $p'$  in Figure 3 is the unique one in  $\tilde{\mathcal{F}}$  such that  $p' \varphi = p\varphi$ . Figure 4 shows two possible choices for  $q'$  satisfying the requirements of Proposition 2.

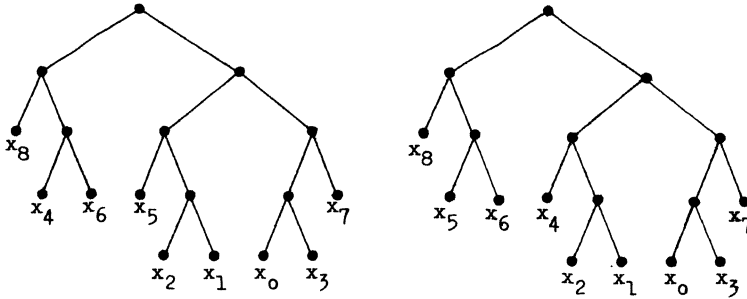


Figure 4

Proof: Let  $p' \in \tilde{\mathbb{F}}$  be the unique polynomial symbol such that  $p'\varphi = p\varphi$  and choose  $q' \in \tilde{\mathbb{F}}$  so that for any  $i < \nu(p')$ , if  $p'\varphi_i$  is an addend in  $p\varphi_j$  then  $q'\varphi_i$  be an addend in  $q\varphi_j$  such that  $q'\varphi_i \alpha = p'\varphi_i \alpha$  (Proposition 1 ensures the existence of such a  $q'$ ). Then, clearly,  $p' = q'$  is in  $\text{Id}(E)$  and  $p = q$  arises from  $p' = q'$  by substituting new (not necessarily distinct) variables.

Let us introduce the following notations: if  $w \in M$ , say  $w = a_{i_0} \dots a_{i_{n-1}}$ , and  $k \leq n$ , put

$$w_k = a_{i_k}, \quad (w)_k = a_{i_0} \dots a_{i_{k-1}}, \quad k(w) = a_{i_k} \dots a_{i_{n-1}},$$

$$\overline{(w)}_k = (w)_{k-1} a_{1-i_{k-1}} \quad \text{and} \quad w^* = w + \sum_{k=1}^n \overline{(w)}_k.$$

It is easy to see that for the polynomial symbols  $s[w] \in P^{(1)}$  ( $w \in M$ ) defined by  $s[1] \equiv x_0$  and for  $n \geq 1$  by  $s[w] \equiv s[{}_1(w)] \cdot x_0$  or  $x_0 \cdot s[{}_1(w)]$  according to whether  $i_0 = 0$  or  $1$ , we have  $s[w]\varphi = w^*$ .

Let  $u, v \in M$ ,  $|u| = n$ ,  $|v| = m$ . Clearly, there exists a polynomial symbol  $q$  such that  $q\varphi = a_0 u^* + a_1 v^*$  (e.g.,

$s[u] \cdot s[v]$  is one). Denote by  $t[a_0 u, a_1 v]$  the unique  $p \in \tilde{\mathcal{F}}$  with  $p\varphi = a_0 u^{\#} + a_1 v^{\#}$ . Observe that these polynomial symbols have exactly 2 subterms of the form  $x_i \cdot x_j$  ( $i, j < \omega$ ).

Example. Figure 5 shows the tree of  $t[a_0 a_1 a_0 a_1, a_1^2 a_0^2]$ .

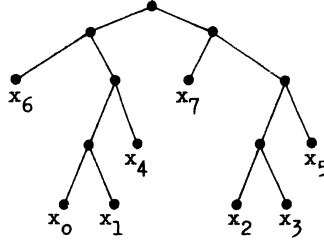


Figure 5

Clearly, if  $u_{n-1} = a_i$  and  $v_{m-1} = a_j$  then

$$t[a_0 u, a_1 v]\varphi_i = a_0 u \quad \text{and} \quad t[a_0 u, a_1 v]\varphi_{j+2} = a_1 v.$$

Denote by  $\sigma(a_0 u, a_1 v)$  the identity  $t[a_0 u, a_1 v] = t^{(i, j+2)}[a_0 u, a_1 v]$  where  $(i, j+2)$  is a transposition. Put

$$\Sigma_0 = \{\sigma(u, v) : u, v \in M, u_0 = a_0, v_0 = a_1, u\alpha = v\alpha\}.$$

Obviously, for  $\sigma(u, v) \in \Sigma_0$  we have  $|u| = |v|$ . This number will be called the depth of  $\sigma(u, v)$ .

Lemma 1. If  $p \in \tilde{\mathcal{F}}$  and  $k, \ell \in \succ(p)$  such that  $p\varphi_k \alpha = p\varphi_\ell \alpha$  then we have  $\sigma(u, v) \vdash_{p=p}^{(k, \ell)}$  for some  $\sigma(u, v) \in \Sigma_0$  of depth  $\leq |p\varphi_k|$ .

Proof: Let  $|p\varphi_k| = |p\varphi_\ell| = n$ ,  $n_{-1}(p\varphi_k) = a_i$  and  $n_{-1}(p\varphi_\ell) = a_j$ . We proceed by induction on the rank of  $p$ . Our claim being trivial if  $p$  is a variable, we can suppose



that it holds for all polynomial symbols of rank smaller than that of  $p$ . We can also assume that  $k \neq l$ , whence  $p\varphi_k \neq p\varphi_l$ . If  $(p\varphi_k)_0 = (p\varphi_l)_0$  then  $x_k$  and  $x_l$  occur in the same subterm of  $p$ , so the lemma follows from the induction hypothesis. Suppose now that they occur in different subterms, say  $(p\varphi_k)_0 = a_0$  and  $(p\varphi_l)_0 = a_1$ . Then it is not hard to show that

$$p \equiv t[p\varphi_k, p\varphi_l](p_0, p_1, \dots) \text{ with } p_i \equiv x_k \text{ and } p_{j+2} \equiv x_l,$$

whence the lemma follows.

Proposition 3.  $\Sigma_0$  is a basis of  $\text{Id}(E)$ .

Proof: By Proposition 2 it suffices to show that for any identity  $p=q$  in  $\text{Id}(E)$  with  $p \in \tilde{\mathcal{P}}$ ,  $q \in \bar{\mathcal{P}}$ , we have  $\Sigma_0 \vdash p=q$ . We proceed by induction. In view of  $(*)$  we shall be done if we prove the following statement: if  $p\varphi_k \neq q\varphi_k$  and for all  $j > k$  we have  $p\varphi_j = q\varphi_j$  then there exists a  $q' \in \bar{\mathcal{P}}$  such that  $\Sigma_0 \vdash q=q'$  and  $p\varphi_i = q'\varphi_i$  for all  $i \geq k$ .

Let  $d = |p\varphi_k|$ . Since  $p \in \tilde{\mathcal{P}}$ , by assumption we have  $p\varphi_j = q\varphi_j$  whenever  $|p\varphi_j| < d$ . Therefore there exists a polynomial symbol  $r \in \mathcal{P}^{(\omega)}$  such that

$$p \equiv r(p_0, p_1, \dots, p_m, x_{k+1}, \dots),$$

$$q \equiv r(q_0, q_1, \dots, q_m, x_{k+1}, \dots)$$

and  $|r\varphi_0| = \dots = |r\varphi_m| = d$ . Since  $p\varphi_k \neq q\varphi_k$ , we may suppose without loss of generality that  $p_0 \equiv x_k$  and  $q_1 \equiv x_k$ . On the other hand,  $p=q$  belongs to  $\text{Id}(E)$ , so that by Proposition 1 we have  $p\varphi_k^\alpha = q\varphi_k^\alpha$ , i.e.  $r\varphi_0^\alpha = r\varphi_1^\alpha$ . Then, by Lemma 1,

$\Sigma_0 \vdash r=r(0,1)$ , whence for

$$q' \equiv r(q_1, q_0, q_2, \dots, q_m, x_{k+1}, \dots)$$

we have  $\Sigma_0 \vdash q = q'$ . Clearly,  $q'$  also satisfies the other requirement.

Lemma 2. Let  $u \in M$ ,  $|u| = n$ , and let  $\sigma(u, v) \in \Sigma_0$  be such that for some  $0 < k < n$  we have  $\overline{(u)_{k+1}} = (v)_{k+1}$ . Then  $\sigma(u, v)$  can be derived from identities of depths  $< n$  in  $\Sigma_0$ .

Proof: Let  $u = a_{i_0} \dots a_{i_{n-1}}$ ,  $v = a_{j_0} \dots a_{j_{n-1}}$  and put  $i = i_{n-1}$ ,  $j = j_{n-1}$ . Since  $u\alpha = v\alpha$ , necessarily  $k < n-1$ . It is not hard to check that

$$\begin{aligned} t[u, v] &\equiv t[(u)_{k+1}, (v)_{k+1}](p_0, \dots, p_3, x_{2n-2k+2}, \dots, x_{2n-1}) \equiv \\ &\equiv t[\overline{(u)_{k+1}}, (v)_{k+1}](p_0, \dots, p_3, x_{2n-2k+2}, \dots, x_{2n-1}) \end{aligned}$$

and the variables  $x_i, x_{j+2}$  occur in  $p_{i_k}, p_{j_{k+2}}$ , respectively. Let  $q$  be the polynomial symbol arising from  $t[u, v]$  by interchanging  $p_{1-i_k}$  and  $p_{j_{k+2}}$ . Clearly,

$$\sigma(\overline{(u)_{k+1}}, (v)_{k+1}) \vdash t[u, v] = q, \quad t^{(i, j+2)}[u, v] = q^{(i, j+2)}.$$

Therefore it remains to show that the identity  $q = q^{(i, j+2)}$  can be derived from an identity of depth  $< n$  in  $\Sigma_0$ . However, this follows from Lemma 1 since by construction

$$q\varphi_i = t[u, v]\varphi_i \quad \text{and} \quad q\varphi_{j+2} = \overline{(u)_{k+1}}(p_{j_{k+2}}\varphi_{j+2}),$$

implying by  $k > 0$  that  $(q\varphi_i)_1 = (q\varphi_{j+2})_1$ . The proof is complete.

Let

$$\Sigma_1 = \{ \sigma(u,v) \in \Sigma_0 : u=u'w, v=v'w, (u)_k^\alpha \neq (v)_k^\alpha \text{ for } 0 < k < |u'|, \text{ and if } (a_i(u)_k)^\alpha = (a_{1-i}(v)_k)^\alpha \text{ for some } i < 2, k < |u'| \text{ then } u_k \neq v_k \}.$$

Proposition 4.  $\Sigma_1$  is a basis of  $\text{Id}(E)$ .

Proof: In virtue of Proposition 3 it suffices to prove that  $\Sigma_1 \vdash \Sigma_0$ . Provisionally, denote by  $\Psi$  the set of all identities in  $\Sigma_0$  that can be derived from  $\Sigma_1$ . Obviously,  $\Sigma_1 \subseteq \Psi \subseteq \Sigma_0$ . Suppose that, contrary to our claim,  $\Psi \neq \Sigma_0$  and choose a  $\sigma(u,v) \in \Sigma_0 - \Psi$  of minimum depth. Let  $m$  be the smallest positive integer such that  $(u)_m^\alpha = (v)_m^\alpha$ , and put  $(u)_m = u', (v)_m = v'$ . Further, let  $u = u' u''$ ,  $v = v' v''$ . Since  $\sigma(u,v) \notin \Sigma_1$ , either  $u'' \neq v''$  or there exist  $k$  and  $i$  ( $k < m, i < 2$ ) such that  $(a_i(u)_k)^\alpha = (a_{1-i}(v)_k)^\alpha$  and  $u_k = v_k$ . We show that in both cases  $\sigma(u,v)$  satisfies the hypotheses of Lemma 2, so that it can be derived from identities of depths  $< |u|$  in  $\Sigma_0$ , which by the minimum property of  $\sigma(u,v)$  implies that  $\Sigma_1 \vdash \sigma(u,v)$ , contradicting our choice.

Indeed, if  $u'' \neq v''$ , say  $u''_\ell \neq v''_\ell$  ( $\ell < |u''|$ ) and  $\ell$  is minimal with respect to this property then

$$\overline{(u)_{n+\ell+1}}^\alpha = ((u)_{n+\ell} v'')^\alpha = (v)_{n+\ell+1}^\alpha.$$

If, in turn,  $(a_i(u)_k)^\alpha = (a_{1-i}(v)_k)^\alpha$  and  $u_k = v_k$  for some  $k < m, i < 2$  then by symmetry we can assume  $u_k = v_k = a_i$ ; so

$$(u)_{k+1}^\alpha = ((v)_k a_{1-i})^\alpha = \overline{(v)_{k+1}}^\alpha,$$

concluding the proof.

Let

$$\Sigma = \{\sigma(u,v) \in \Sigma_1 : \text{at least one of } u, v \text{ ends with } a_0\}.$$

Now we are ready to state our main theorem.

Theorem.  $\Sigma$  is an independent basis of  $\text{Id}(E)$ .

Corollary.  $E$  has no finite basis for its identities.

The crucial part of the proof of the Theorem will be formulated in a separate lemma below. Denote by  $X$  the set of all pairs  $(u,v)$  such that  $\sigma(u,v) \in \Sigma$ . Let  $(u,v) \in X$  and  $u = a_{i_0} \dots a_{i_{n-1}}$ ,  $v = a_{j_0} \dots a_{j_{n-1}}$ . Clearly, by the definition of  $\Sigma$  we have

- (i)  $u\alpha = v\alpha$ ;
- (ii)  $i_0=0, j_0=1$ ;
- (iii) for all  $0 < k < n$ , if  $(u)_k\alpha = (v)_k\alpha$  then  $i_k=j_k$ ;
- (iv) if there exist  $i < 2, 0 < k < n$  such that  $((u)_k a_{i_1})\alpha = ((v)_k a_{j_{1-1}})\alpha$  then  $i_k \neq j_k$ .

Lemma 3. Let  $(u,v) \in X$  and  $p \in P^{(\omega)}$  such that  $p\varphi\alpha = t[u,v]\varphi\alpha$ . Then  $p\varphi = t[u,v]\varphi$ .

Proof: From the definition of  $t[u,v]$  it follows immediately that

$$T = t[u,v]\varphi = u + \sum_{j=2}^n \overline{(u)_j} + v + \sum_{j=2}^n \overline{(v)_j}.$$

Thus, for  $2 \leq j < n$ , the only words of lengths  $j$  entering the sum are  $\overline{(u)_j}$  and  $\overline{(v)_j}$ . Now let  $p\varphi\alpha = T\alpha$ . We have to show that every addend of  $T$  occurs in  $p\varphi$ , too. We proceed by induction on the lengths of the words. From (ii) and

(iv) it follows that either  $\overline{(u)}_2 = a_0 a_1$ ,  $\overline{(v)}_2 = a_1 a_0$  or  $\overline{(u)}_2 = a_0^2$ ,  $\overline{(v)}_2 = a_1^2$ . Since  $p\varphi\alpha = T\alpha$  and the addends of  $p\varphi$  are distinct, in both cases  $\overline{(u)}_2$  and  $\overline{(v)}_2$  must occur in  $p\varphi$ .

Suppose now that  $\overline{(u)}_j$  and  $\overline{(v)}_j$  enter  $p\varphi$  for some  $2 \leq j < n$ . First we show that any addend  $w$  of length  $j+1$  in  $p\varphi$  is of the form  $(u)_j a_i$  or  $(v)_j a_i$  for some  $i < 2$ . Indeed, as  $\overline{(u)}_j$  and  $\overline{(v)}_j$  occur in  $p\varphi$ ,  $p$  must have two subterms with weights  $(u)_j$  and  $(v)_j$ , respectively. If either one of these subterms were the product of two terms of lengths  $\geq 2$ , then  $p$  would have more than two subterms of lengths 2. However, if  $x_k \cdot x_\ell$  is a subterm of  $p$  then  $p\varphi_k\alpha = a_0^{r+1} a_1^s$ ,  $p\varphi_\ell\alpha = a_0^r a_1^{s+1}$ , but  $p\varphi\alpha$  contains only two pairs of members of this kind, namely  $u$ ,  $\overline{(u)}_n$  and  $v$ ,  $\overline{(v)}_n$ . Thus  $p\varphi$  must contain two words of the form  $(u)_j a_i$  and  $(v)_j a_i$ , respectively, if  $j < n-1$  and the four words  $u$ ,  $\overline{(u)}_n$ ,  $v$ ,  $\overline{(v)}_n$  if  $j = n-1$ . Since  $p\varphi\alpha = T\alpha$ ,  $p\varphi$  has no other addend of length  $j+1$ .

Now we are ready to complete the induction step. If  $j = n-1$  then, as we proved in the previous paragraph, every addend of length  $j+1 = n$  of  $T$  must occur in  $p\varphi$ . Suppose now that  $j < n-1$  and  $\overline{(u)}_{j+1}$  doesn't enter  $p\varphi$ . Since  $p\varphi\alpha = T\alpha$ ,  $p\varphi$  has an addend  $w$  such that  $w\alpha = \overline{(u)}_{j+1}\alpha$ . By the above statement  $w$  equals  $(u)_j a_i$  or  $(v)_j a_i$  for some  $i < 2$ . Assume the first. Then, obviously,  $u_j = a_i$  because else we would have  $\overline{(u)}_{j+1} = (u)_j a_i = w$ , contrary to the assumption. Hence  $\overline{(u)}_{j+1}\alpha = ((u)_j a_{1-i})\alpha \neq w\alpha$ , which is not the case. Thus  $w = (v)_j a_i$ . We can assume that  $w \neq \overline{(v)}_{j+1}$  whence  $(v)_j a_i = \overline{(v)}_{j+1}$ ,  $v_j = a_i$ . However, then

$\overline{(u)_{j+1}^\alpha} = (v)_{j+1}^\alpha$ , which contradicts (iii) or (iv) depending on whether  $u_j$  and  $v_j$  (i.e., the last letters of  $(u)_{j+1}$  and  $(v)_{j+1}$ ) are distinct or not. This completes the proof of the lemma.

Let  $p \in \overline{\mathbb{F}}$ ,  $i, j \in \nu(p)$ . We shall say that the variables  $x_i$  and  $x_j$  are linked in  $p$  if  $x_i \circ x_j$  or  $x_j \circ x_i$  is a subterm of  $p$ . Equivalently,  $x_i$  and  $x_j$  are linked iff  $p\varphi_i$  and  $p\varphi_j$  are of the same length and differ in their last letters only. For example, in the polynomial symbol  $t[u, v]$  ( $u, v \in M$ ,  $|u| = |v|$ ),  $x_0, x_1$  and  $x_2, x_3$  are linked with each other, and they are the only variables which are linked with another one.

Proof of the Theorem: To show that  $\Sigma$  is a basis of  $\text{Id}(E)$ , by Proposition 4 it suffices to note that if  $|u| = |v| = n$  and  $u_{n-1} = v_{n-1} = a_1$  then  $\sigma(u, v)$  can be derived from  $\Sigma$  as follows:

$$\sigma((u)_{n-1}a_0, (v)_{n-1}a_0) \vdash t[u, v] = t^{(0,2)}[u, v]$$

and

$$\sigma((u)_{n-1}, (v)_{n-1}) \vdash t^{(0,2)}[u, v] = t^{(1,3)}[u, v].$$

Next we prove that  $\Sigma$  is independent. By way of contradiction suppose that for  $\sigma(u, v) \in \Sigma$ ,  $\Sigma' = \Sigma - \{\sigma(u, v)\}$  we have  $\Sigma' \vdash \sigma(u, v)$ . Choose the permutations  $\pi, \rho$  on  $\{i: i < 2n\}$  so that the shortest derivation of the identity

$$(***) \quad t^\pi[u, v] = t^\rho[u, v]$$

from  $\Sigma'$  be of minimum length among all those of form (\*\*\*)

for which there exists a variable which is linked with different variables on the two sides. Clearly, such an identity is not contained in  $\Sigma'$ . (Observe that when we replaced  $\Sigma_1$  by  $\Sigma$ , we omitted exactly those identities from  $\Sigma_1$  which would have violated this.)

We will arrive at a contradiction by proving that the last step of the shortest derivation of (\*\*\*) cannot be the application of any one of rules (1)-(5) in [3; p. 381]. This is obvious for (1). By the minimality condition it follows immediately for (2) and (3), too, noticing that if for some  $r \in P^{(\omega)}$  we have  $\Sigma' \vdash t^\pi[u, v] = r$  (and hence  $t^\pi[u, v] = r$  belongs to  $\text{Id}(E)$ ) then by Lemma 3 and Proposition 1 there exists a permutation  $\tau$  on  $\{i: i < 2n\}$  such that  $r \equiv t^\tau[u, v]$ .

If  $t^\pi[u, v] \equiv p_0 \circ p_1$ ,  $t^\rho[u, v] \equiv r_0 \circ r_1$  and  $\Sigma' \vdash p_0 = r_0$ ,  $p_1 = r_1$  then clearly  $p_0 = r_0$ ,  $p_1 = r_1$  belong to  $\text{Id}(E)$ , so by the construction of  $t[u, v]$  one easily infers that  $p_0 \equiv r_0$  and  $p_1 \equiv r_1$ . Therefore  $\pi = \rho$ , contradicting our choice. This settles case (4).

Finally, suppose in the last step of the derivation of (\*\*\*) rule (5) is applied and, say, the polynomial symbols  $r_i$  ( $i < m$ ) are substituted for the variables  $x_i$  ( $i < m$ ). By the minimality condition at least one of the  $r_i$ 's is not a variable and hence contains a pair of linked variables, which are linked in  $t^\pi[u, v]$  and  $t^\rho[u, v]$ , too. On the other hand, from the definition of  $t[u, v]$  it follows that  $t^\pi[u, v]$  and  $t^\rho[u, v]$  have exactly two linked pairs of variables. Therefore the relation "linkedness" of the vari-

ables in  $t^\pi[u,v]$  and  $t^0[u,v]$  coincide, contradicting our assumption. The proof of the Theorem is complete.

Remark. Along the same lines one can easily construct an (infinite) independent basis for the identities of algebras  $\langle R;f \rangle$  where  $f$  is an  $n$ -ary ( $n \geq 2$ ) operation

$$f(x_0, \dots, x_{n-1}) = \sum_{i < n} r_i x_i$$

whose coefficients  $r_i$  ( $i < n$ ) are algebraically independent.

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