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VALUATIONS OF STRUCTURES
J. MLČEK

Abstract: This paper is a contribution to the development of the alternative set theory. A typical special result among those presented is the following: Let $\mathcal{A} = \langle \mathfrak{a}, f \rangle$ be a set-semigroup and let $\mathcal{A}/Q = \langle Q, f/Q^2 \rangle$ where $Q \subseteq \mathfrak{a}$ is a π -class be a substructure of \mathcal{A} . Then there exists a set-mapping $h: \mathfrak{a} \rightarrow \mathbb{R}N(\geq 0)$ ($\mathbb{R}N(\geq 0)$ is the class of non-negative rationals) such that $h(f(x, y)) \leq h(x) + h(y)$ and $h(x) \doteq 0 \Leftrightarrow x \in Q$ holds for each $x, y \in \mathfrak{a}$. (As usual, we write $z \doteq 0$ if $|z| < n$ for all finite natural numbers n .)

We present more general results; namely, they concern some richer structures than that of a semigroup, deal also with proper classes, and the universe Q of the substructure \mathcal{A}/Q is a σ - or π -class.

As a consequence of our results we obtain a metrization theorem.

Key words: Structure, valuation, σ -class, π -class, metrization.

Classification: 02K10, 02K99, 08A05, 54J05

§ 0. Introduction. Great numbers of important structures are constructed in the alternative set theory by using π -classes. For example, real numbers are constructed as factor-classes of the π -equivalence \doteq on the class $\mathbb{R}N$ of rational numbers. (See [V].) The topological structure is comprehended as a π -equivalence on a set-theoretically definable class. In this paper we study structures which are described by using σ -classes and π -classes only. Let us explain

our problems in more detail on the structure $\langle a^2, \sim \rangle$, where a is a set and \sim is a π -equivalence on a . Using some ideas of the proof of the classic metrization lemma, we can prove that there is a set-mapping $h: a^2 \rightarrow \mathbb{R}N(\geq 0)$ ($\mathbb{R}N(\geq 0)$ denotes the class of non-negative rationals) such that $h(x,z) \leq h(x,y) + h(y,z)$, $h(x,y) = h(y,x)$, $h(x,y) \stackrel{\circ}{=} 0 \equiv x \sim y$, $h(x,y) = 0 \equiv x = y$ hold. (h is called metric of \sim on a .) We can say that h is a valuation of a^2 in $\mathbb{R}N(\geq 0)$ such that h respect (in the sense mentioned above) the following couples of operations: the operation \circ (the composition of pairs) and $+$; the operation Cn of converse and the identity mapping Id . Moreover, the values of all elements of \sim are exactly in $[\geq 0] = \{x \in \mathbb{R}N(\geq 0); x \stackrel{\circ}{=} 0\}$. We shall describe a class of structures of the type $\langle A, F, E \rangle$, where F is a binary function and E is a unary function, such that the following statement holds: if \mathcal{A} is a set-structure of this class and \mathcal{A}/Q is a substructure of \mathcal{A} with the universe Q , which is a π -class, then the pair $\langle \mathcal{A}, \mathcal{A}/Q \rangle$ is valued in $\langle \langle \mathbb{R}N(\geq 0), +, Id \rangle, \langle [\geq 0], +, Id \rangle \rangle$ by a set-mapping similarly as a set-metric of \sim on a values $\langle \langle a^2, \circ, Cn \rangle, \langle \sim, \circ, Cn \rangle \rangle$ in $\langle \langle \mathbb{R}N(\geq 0), +, Id \rangle, \langle [\geq 0], +, Id \rangle \rangle$.

Note that we do not work with set-structures only but the structure \mathcal{A} mentioned can be generally a structure from a standard system \mathcal{M} and the universe Q of the substructure \mathcal{A}/Q can be a $\pi^{\mathcal{M}}$ -or a $\sigma^{\mathcal{M}}$ -class. Then we construct a valuation of the pair $\langle \mathcal{A}, \mathcal{A}/Q \rangle$ as a class of \mathcal{M} .

(For the notions of the standard systems and $\pi^{\mathcal{M}}$ -and $\sigma^{\mathcal{M}}$ -class see [M1].)

Let us mention one consequence of our general results. Recall that $x \stackrel{\circ}{=} y$ iff for each set-formula $\varphi(z)$ in FL we have

$\varphi(x) \equiv \varphi(y)$. The following statement holds: there is a metric of \equiv on V which is an element of a revelation Sd_V^* of the codable class Sd_V of all set-theoretically definable classes (i.e., roughly speaking, there is a "formally set-theoretically definable" metric of \equiv on V . (For the notion of the revelations see [S-V 1].))

Further results concerning the problems of valuations will be presented in another paper.

§ 1. Preliminaries

1.0.0. We use usual definitions and notions of the alternative set theory and definitions, notions and symbols introduced in [ML]. We shall use results obtained in [ML].

1.0.1. Throughout this paper let \mathcal{M} denote a standard system.

§ 2. e-structures. Valuations

2.0.0. By a structure we mean a $m+n+1$ -tuple $\mathcal{A} = \langle A, F_i, R_j \rangle_{i \in m, j \in n}$, $m, n \in FN$, where, for each $i \in m$, F_i is a $a(i)$ -ary function, $\text{dom}(F_i) = A^{a(i)}$, $F_i^* A^{a(i)} \subseteq A$, $a(i) \in FN$ and, for each $j \in n$, $R_j \subseteq A^{b(j)}$, $b(j) \in FN$.

We say that a class $B \subseteq A$ is a universe in \mathcal{A} iff, for each $i \in m$, $F_i^* B^{a(i)} \subseteq B$ holds. A substructure of the structure \mathcal{A} is a structure $\langle B, F_i \upharpoonright B^{a(i)}, R_j \cap B^{b(j)} \rangle_{i \in m, j \in n}$ where B is a universe in \mathcal{A} . We denote the substructure presented by \mathcal{A}/B . If there is no danger of confusion, we write $\langle B, F_i, R_j \rangle$ instead of $\langle B, F_i \upharpoonright B^{a(i)}, R_j \cap B^{b(j)} \rangle_{i \in m, j \in n}$.

2.0.1. A covariant (contravariant resp.) e-structure is a structure $\langle A, F, E \rangle$ where F is a binary function, E is a

unary function and the following holds: (1) F is associative on A ,

$$(2) E \circ E = \text{Id}$$

$$(3) F(E(x), E(y)) = E(F(x, y))$$

$$(F(E(x), E(y)) = E(F(y, x)) \text{ resp.})$$

holds for each $x, y \in A$.

An e-structure is a covariant or a contravariant e-structure. An e-structure $\mathcal{A} = \langle A, F, E \rangle$ is a commutative e-structure iff F is commutative on A .

Then \mathcal{A} is covariant and contravariant simultaneously. An e-structure $\langle A, F, \text{Id} \rangle$ is covariant. It is contravariant iff it is commutative. Let $\mathcal{A} = \langle A, F, E \rangle$ be an e-structure. We define the binary relation on A as follows:

$$x \triangleleft_{\mathcal{A}} y \equiv (\exists z \in A)(F(x, z) = y).$$

If there is no danger of confusion, we shall write simply \triangleleft instead of $\triangleleft_{\mathcal{A}}$.

Proposition. The relation $\triangleleft_{\mathcal{A}}$ is transitive on A .

2.0.2. Examples. (1) A structure $\langle A, F \rangle$ is a semigroup iff $\langle A, F, \text{Id} \rangle$ is a covariant e-structure.

(2) $\langle \mathbb{N}, +, \text{Id} \rangle$ is a commutative e-structure.

(3) Let $\mathbb{RN}(\geq 0) = \{x \in \mathbb{RN}; x \geq 0\}$, $\mathbb{RN}(> 0) = \{x \in \mathbb{RN}; x > 0\}$. $\langle \mathbb{RN}(\geq 0), +, \text{Id} \rangle$ and $\langle \mathbb{RN}(> 0), \cdot, {}^{-1} \rangle$ are commutative e-structures.

(4) We put, for $X \subseteq \mathbb{N}$, $X_2 = \{2^{\alpha}; \alpha \in X\}$. $\langle \mathbb{N}_2, \cdot, \text{Id} \rangle$ is a commutative e-structure.

(5) Let a be a set, $a \neq \emptyset$. Then $\langle P(a), \cup, \text{Id} \rangle$, $\langle P(a), \cap, \text{Id} \rangle$ are commutative e-structures.

(6) We define the mapping $F^0: (V^2 \cup \{0\})^2 \rightarrow V^2 \cup \{0\}$ as follows: $F^0(\langle x, y \rangle, \langle u, v \rangle) = \langle x, v \rangle$ (0 resp.) iff $y = u$ ($y \neq u$

resp.) and $F^0(w,0) = F^0(0,w) = 0$ for each $w \in V^2 \cup \{0\}$.

F^0 is an associative function on $V^2 \cup \{0\}$ and, consequently, $\langle V^2 \cup \{0\}, F^0, Id \rangle$ is an e-structure, which is not commutative. Let R be a transitive relation. Then $\langle R \cup \{0\}, F^0, Id \rangle$ is an e-structure and the following holds:

$(\forall u \in R \cup \{0\})(u \triangleleft 0) \& (\forall u \in R \cup \{0\})(0 \triangleleft u \equiv u = 0)$.

2.0.3. Lemma. Let $\langle A, F, E \rangle$ be an e-structure. Let A_0, A_1 be classes such that $A_0 \subseteq A_1 \subseteq A$ and $\llbracket F, E \rrbracket(A_0, A_1)$ hold. Let $Q_i = E^*A_i \cap A_i$ for $i = 0, 1$.

Then $Q_0 \subseteq A_0 \subseteq Q_1 \subseteq A_1$ and, for $i = 0, 1$, $F^*Q_0^2 \subseteq Q_1$, $E^*Q_i \subseteq Q_i$.

Proof. The relation $Q_i \subseteq A_i$, $i = 0, 1$, is obvious. 1) We prove that $A_0 \subseteq Q_1$. Let $x \in A_0$. We have $E(x) \in A_1$, $x \in A_1$ and $x = E(E(x))$. Thus $x \in A_1 \cap E^*A_1$. 2) We prove that $F^*Q_0^2 \subseteq Q_1$. Let $x, y \in Q_0$. Thus $x, y \in A_0$ and $x = E(u)$, $y = E(v)$ hold with some $u, v \in A_0$. We have $F(x, y) \in A_1$, $F(u, v) \in A_1$ and $F(v, u) \in A_1$. Thus $F(x, y) = F(E(u), E(v)) \in E^*A_1$ holds. We deduce from this that $F(x, y) \in A_1 \cap E^*A_1$. 3) Let us prove that $E^*Q_i \subseteq Q_i$ holds for $i = 0, 1$. Let $x \in Q_i$. Then $x \in A_i$ and there is a $y \in A_i$ such that $x = E(y)$. Consequently, $E(x) \in A_i \cap E^*A_i$ holds.

2.0.4. Let \mathcal{A} be an e-structure. Let Q, B be universes in \mathcal{A} . The triple $\langle \mathcal{A}, \mathcal{A}/Q, \mathcal{A}/B \rangle$ is called a triad over \mathcal{A} . Let $\mathcal{A}(Q, B)$ denote this triad. A triad of the type $\mathcal{C}^{\mathcal{M}}$ (or a $\mathcal{C}^{\mathcal{M}}$ -triad) is a triad $\mathcal{A}(Q, B)$ such that $\mathcal{A} \in \mathcal{M}$, $B \in \mathcal{M}$ and Q is a $\mathcal{C}^{\mathcal{M}}$ -class. We define a triad of the type $\mathcal{M}^{\mathcal{M}}$ (or a $\mathcal{M}^{\mathcal{M}}$ -triad) analogously.

Examples. (1) $\langle N, +, Id \rangle (FN, \{0\})$, $\langle N_2, \cdot, Id \rangle (FN_2, \{1\})$ are \mathcal{C}^0 -triads.

(2) Let a be a set, $a \neq 0$ and let Q be an ideal on $P(a)$.

Then $\langle P(a), \cup, Id \rangle (Q, \{0\})$ is a triad. Suppose, moreover, that Q is a σ (σ' resp.)-class. Then the triad presented is a σ - triad (σ' -triad resp.).

(3) The equivalence \cong on RN is defined as follows:

$$(\forall x, y \in RN)(x \cong y \equiv (\forall n)(|x-y| < \frac{1}{n} \vee (x > n \& y > n) \vee (x < -n \& y < -n))).$$

We put $[\geq 0] = \{y \in RN(\geq 0); y \cong 0\}$. Then

$\langle RN(\geq 0), +, Id \rangle ([\geq 0], \{0\})$ is a σ^0 -triad.

2.1.0. Let $a = \langle A, F, E \rangle$, $\tilde{a} = \langle \tilde{A}, \tilde{F}, \tilde{E} \rangle$ be e-structures. A mapping $H: A \rightarrow \tilde{A}$ is called valuation of a in \tilde{a} iff for each $x, y \in A$ holds:

$$H(F(x, y)) \triangleleft_{\tilde{a}} F(H(x), H(y))$$

$$H(E(x)) = E(H(x)).$$

Let $a(Q, B)$, $\tilde{a}(\tilde{Q}, \tilde{B})$ be triads. A mapping $H: A \rightarrow \tilde{A}$ is called valuation of the triad $a(Q, B)$ in the triad $\tilde{a}(\tilde{Q}, \tilde{B})$ iff H is a valuation of a in \tilde{a} and we have for each $x \in A$:

$$x \in Q \equiv H(x) \in \tilde{Q}, \quad x \in B \equiv H(x) \in \tilde{B}.$$

Example. The mapping $H: N \rightarrow N_2$ sending ∞ to 2^∞ is a valuation of $\langle N, +, Id \rangle (FN, \{0\})$ in $\langle N_2, \cdot, Id \rangle (FN_2, \{1\})$.

Proposition. Let a be an e-structure and let \triangleleft_a be reflexive on A . Let $a(Q, B)$ be a triad over a and let $A' \subseteq A$ be an universe in a .

(1) $a/A'(Q \cap A', B \cap A')$ is a triad over a/A' .

(2) Identity mapping Id is a valuation of $a/A'(Q \cap A', B \cap A')$ in $a(Q, B)$.

Proof. (1) follows from the fact that $Q \cap A'$ and $B \cap A'$ are universes in a/A' . (2) Identity mapping is a valuation of a/A' in a (by using of the reflexivity of \triangleleft_a).

Proposition. Let $\tilde{a} = \langle \tilde{A}, \tilde{F}, \tilde{E} \rangle$ be a commutative e-structure and let $\tilde{a}(\tilde{Q}, \tilde{B})$ be a triad. Suppose that there exist

points $a, q, b \in \tilde{A}$ such that $b \triangleleft q \triangleleft a$ and $b \in \tilde{B}, q \in \tilde{Q}-\tilde{B}, a \in \tilde{A}-\tilde{Q}$.

Then, for each triad \mathcal{T} , there is a valuation of \mathcal{T} in $\tilde{a}(\tilde{Q}, \tilde{B})$.

Proof. Let H be a mapping, defined as follows:

$H(x) = b \equiv x \in B, H(x) = q \equiv q \in Q-B, H(x) = a \equiv x \in A-Q$, where $\langle A, F, E \rangle (Q, B) = \mathcal{T}$. The H is the required valuation.

§ 3. Valuation lemmas

3.0.0. We shall prove two lemmas which have the important role for the construction of valuations of $\sigma^{\mathcal{M}}$ -triads and $\pi^{\mathcal{M}}$ -triads. At first, we introduce the following definition: let $\mathcal{A} = \langle A, F, E \rangle$ be an e-structure and let B be an universe in \mathcal{A} . A σ -string (π -string resp.) R is called σ (π resp.)-string in \mathcal{A} over B iff $B = R(0), A = R(\text{dom}(R)-1)$ and $\llbracket F, F_3 \rrbracket (R(\alpha), R(\alpha+1)), E^*R(\alpha) \subseteq R(\alpha)$ holds for each $\alpha \in \text{dom}(R)-1$ ($A = R(0), B = R(\text{dom}(R)-1)$ and $\llbracket F, F_3 \rrbracket (R(\alpha+1), R(\alpha)), E^*R(\alpha) \subseteq R(\alpha)$ holds for each $\alpha \in \text{dom}(R)-1$ resp.), where $F_3: A^3 \rightarrow A$ is the function satisfying $F_3(x, y, z) = F(F(x, y), z)$.

3.0.1. σ -valuation lemma. The following holds in the sense of \mathcal{M} : Let \mathcal{A} be an e-structure and let B be an universe in \mathcal{A} . Let Q be a σ -string in \mathcal{A} over B and let $\xi+1 = \text{dom}(Q)$.

Then there is a valuation H of the triad $\mathcal{A}(B, B)$ in $\langle N, +, \text{Id} \rangle (\{0\}, \{0\})$ such that $Q(\alpha) \subseteq \{x \in A; H(x) \triangleleft 2^\alpha\} \subseteq \subseteq Q(\alpha+1)$ holds for each $\alpha \in \xi$.

π -valuation lemma. The following holds in the sense of \mathcal{M} : Let \mathcal{A} be an e-structure and let B be an universe in \mathcal{A} .

Let Q be a π -string in \mathcal{A} over B and let $\xi+1 = \text{dom}(Q)$.

Then there is a valuation H of the triad $\mathcal{A}(B, B)$ in $\langle \mathbb{N}(\geq 0), +, \text{Id} \rangle (\{0\}, \{0\})$ such that $Q(\alpha+1) \subseteq \{x \in A; H(x) \leq 2^{-(\alpha+1)}\} \subseteq Q(\alpha)$ holds for each $\alpha \in \xi$.

The π -valuation lemma follows from the σ -valuation lemma. Really, let G be a valuation of $\mathcal{A}(B, B)$ in $\langle \mathbb{N}, +, \text{Id} \rangle (\{0\}, \{0\})$ such that $Q(\xi-\alpha) \subseteq \{x \in A; G(x) \leq 2^{-\alpha}\} \subseteq Q(\xi-(\alpha+1))$ holds for each $\alpha \in \xi$. We put $\beta = \xi - \alpha$. Thus, $Q(\beta) \subseteq \{x \in A; G(x) \leq 2^{\beta-\beta}\} \subseteq Q(\beta-1)$ holds for each $1 \leq \beta \in \xi$. The required valuation is the mapping $H = 2^{-\xi} \cdot G$.

3.0.2. The proof of the σ -valuation lemma.

I. A path in A is a function t such that $\text{dom}(t) \subseteq \mathbb{N}$ and $\text{rng}(t) \subseteq A$. We construct the function $[F]$ with domain

$\cup \{t\} \times \{ \langle \alpha, \beta \rangle; \alpha \leq \beta \ \& \ \beta \in \text{dom}(t) \}$; t is a path in A by induction over \mathbb{N} :

$$[F](t, \langle \alpha, \alpha \rangle) = t(\alpha)$$

$$[F](t, \langle \alpha, \beta+1 \rangle) = F([F](t, \langle \alpha, \beta \rangle), t(\beta+1)).$$

We shall write more simply $[F](t, \alpha, \beta)$ instead of $[F](t, \langle \alpha, \beta \rangle)$.

Lemma 1. Let t be a path in A , $\alpha \leq \gamma+1 \leq \beta \in \text{dom}(t)$.

Then

$$[F](t, \alpha, \beta) = F([F](t, \alpha, \gamma), [F](t, \gamma+1, \beta))$$

holds.

This follows by induction on $\beta - \alpha$.

Let t be a path in A , $\text{dom}(t) = \vartheta+1$. We define the path \bar{t} with $\text{dom}(\bar{t}) = \vartheta+1$ as follows: $\bar{t}(\alpha) = t(\vartheta-\alpha)$.

$\tilde{F}: A^2 \rightarrow A$ is the function so that $\tilde{F}(x, y) = F(y, x)$ holds for each $x, y \in A$. $[\tilde{F}]$ is defined similarly as $[F]$.

The following lemma can be proved by induction on $\beta - \alpha$.

Lemma 2. Let t be a path in A , $\text{dom}(t) = \nu^{\delta} + 1$. Then

$$[FJ](t, \alpha, \beta) = [\tilde{F}](t, \nu^{\delta} - \beta, \nu^{\delta} - \alpha)$$

holds for each $\alpha \leq \beta \leq \nu^{\delta}$.

II. We put for each $x \in A$: $G_Q(x) = \min\{\alpha \leq \xi; x \in Q(\alpha)\}$. Thus, G_Q is a function, $G_Q: A \rightarrow N$, and we have $G_Q(x) \leq \alpha \equiv x \in Q(\alpha)$, $\alpha < G_Q(x) \equiv x \notin Q(\alpha)$ for each $\alpha \leq \xi$. We shall write more simply G instead of G_Q . The index Q denotes only that G_Q is constructed from Q and this notion will be used in 3.0.3.

We define the function G^* , $G^*: A \rightarrow N$, as follows:

$$G^*(x) = 0 \text{ iff } x \in B, \quad G^*(x) = 2^{G(x)} \text{ iff } x \in A - B.$$

Let t be a path in A . We put

$$\mathcal{V}_Q(t) = \sum\{G^*(x); x \in \text{rng}(t)\}.$$

We shall write more simply \mathcal{V} instead of \mathcal{V}_Q . \mathcal{V} is a function, $\text{rng}(\mathcal{V}) \subseteq N$.

We deduce from the definition of \mathcal{V} that $\mathcal{V}(t) = 0 \equiv \text{rng}(t) \subseteq B$ and $\mathcal{V}(t) = 0 \rightarrow (\forall \alpha, \beta \in \text{dom}(t))(\alpha \leq \beta \rightarrow [FJ](t, \alpha, \beta) \in B)$.

Let t be a path in A , $\text{dom}(t) = \sigma^{\delta} + 1$. Writing $[FJ](t)$ ($[\tilde{F}J](t)$ resp.) we mean $[FJ](t, 0, \sigma^{\delta})$ ($[\tilde{F}J](t, 0, \sigma^{\delta})$ resp.). Note that whenever $[FJ](t, \alpha, \beta)$ appears, then we assume that $\langle t, \langle \alpha, \beta \rangle \rangle$ is an element of $\text{dom}([FJ])$. We use the similar convention for the terms $[FJ](t)$, $[\tilde{F}J](t, \alpha, \beta)$, $[\tilde{F}J](t)$.

Lemma 3. Let $z \in A$ and suppose that $[FJ](t) = z$. Then

$$(*) \quad \mathcal{V}(t) \neq 0 \rightarrow 2^{G(z)} \leq 2 \cdot \mathcal{V}(t)$$

holds.

Proof. By induction on $\text{dom}(t)$.

(i) Suppose that $\text{dom}(t) = 2$. Assume, for example that $G(t(0)) \leq G(t(1))$. Thus $G(z) \leq G(t(1) + 1)$ holds and we have $2^{G(z)} \leq 2 \cdot 2^{G(t(1))}$. If $t(1) \in B$ then $G(t(1)) = 0$ and, consequently, $G(t(0)) = 0$. We deduce from this that $t(0) \in B$, which

is a contradiction. Thus, $t(1) \notin B$ holds and we have $2 \cdot 2^{G(t(1))} \leq 2 \cdot (G^*(t(0)) + 2^{G(t(1))}) = 2 \cdot \mathcal{V}(t)$.

(ii) Suppose that the statement $(*)$ holds whenever $\text{dom}(t) \leq \beta + 1$ and $\beta + 1 \geq 3$ is fixed. Let t be a path in A and let $\text{dom}(t) = \beta + 2$. Let $[F](t) = z$ and assume that $\mathcal{V}(t) \neq 0$. We shall prove that $2^{G(z)} \leq 2 \cdot \mathcal{V}(t)$ holds.

We put $c = \mathcal{V}(t)$. Let σ be the maximal natural number such that $2^\sigma \leq c$. If $\sigma \geq \xi - 1$ then $2^{G(z)} \leq 2 \xi \leq 2^{\sigma+1} \leq 2 \cdot 2^\sigma \leq 2 \cdot c$ and, consequently, the statement in question is proved. Assume $\sigma < \xi - 1$.

(α) Suppose that $G^*(t(0)) \leq \frac{c}{2}$. Let $\gamma \in \mathbb{N}$ be a maximal number such that

$$\mathcal{V}(t \wedge \gamma + 1) = \sum_{\alpha=0}^{\gamma} G^*(t(\alpha)) \leq \frac{c}{2}.$$

Obviously, $0 \leq \gamma \leq \beta$. Moreover, $0 \neq G^*(t(\gamma + 1)) \leq c$ and

$\sum_{\alpha=\gamma+2}^{\beta+1} G^*(t(\alpha)) \leq \frac{c}{2}$. We put $z_1 = [F](t, 0, \gamma)$, $z_3 = [F](t, \gamma + 2, \beta + 1)$.

Suppose that $\sum_{\alpha=0}^{\gamma} G^*(t(\alpha)) \neq 0$. We deduce from the induction hypothesis that $2^{G(z)} \leq 2 \cdot \frac{c}{2} = c$. Thus, the following relation holds:

- (*) $G(z_1) \leq \sigma$. It is easy that
- (**) $G(t(\gamma + 1)) \leq \sigma$. We deduce as above that
- (***) $G(z_3) \leq \sigma$

follows from $\sum_{\alpha=\gamma+2}^{\beta+1} G^*(t(\alpha)) \neq 0$.

The relations (*), (**), (***) hold too in the case if

$\sum_{\alpha=0}^{\gamma} G^*(t(\alpha)) = 0$ or $\sum_{\alpha=\gamma+2}^{\beta+1} G^*(t(\alpha)) = 0$. We have $z = [F](t) = F(F(z_1, t(\gamma + 1)), z_3) = F_3(z_1, t(\gamma + 1), z_3)$ and $F_3^3 Q^3(\sigma) \subseteq Q(\sigma + 1)$.

We deduce from this that $z \in Q(\sigma + 1)$. Consequently, $G(z) \leq \sigma + 1$

holds, and

$$2^{G(z)} \leq 2^{\alpha+1} = 2 \cdot 2^\alpha \leq 2c = 2 \cdot \mathcal{V}(t)$$

follows immediately.

(β) Suppose that $G^*(t(0)) > \frac{c}{2}$. Then $G^*(t(\beta+1)) \leq \frac{c}{2}$. Thus, $G^*(\bar{t}(0)) = G^*(t(\beta+1)) \leq \frac{c}{2}$ holds. We have $[\tilde{F}](t) = z = [F](t)$ (by using the lemma 2). We deduce similarly as in the case (∞) that $2^{G(z)} \leq 2c$ holds.

III. The following definition of the function $H: A \rightarrow N$ is justified:

$$H(x) = \min \{ \mathcal{V}(t); [F](t) = x \}.$$

We shall prove that H is the valuation in question.

(a) $H(x) = 0 \equiv x \in B$. Suppose that $H(x) = 0$. Then there exists a path t in A such that $H(x) = \mathcal{V}(t)$ and $[F](t) = x$. Thus, $x \in B$ holds. Suppose that $x \in B$. We have $G^*(x) = 0$ and $H(x) = 0$ follows from the relation $H(x) \leq \mathcal{V}(\{ \langle x, 0 \rangle \}) = G^*(x) = 0$.

(b) $Q(\infty) \subseteq \{ x \in A; H(x) \leq 2^\alpha \} \subseteq Q(\alpha+1)$ holds for each $\alpha \in \mathbb{F}$. At first, we prove that

$$(\times \times) \quad x \in A-B \rightarrow 2^{-1} \cdot 2^{G(x)} \leq H(x) \leq 2^{G(x)} \text{ holds.}$$

Proof. Let t be a path in A such that $[F](t) = x$ and $\mathcal{V}(t) = H(x)$. We have $\mathcal{V}(t) \neq 0$ and, consequently, $2^{-1} \cdot 2^{G(x)} \leq \mathcal{V}(t) \leq H(x)$. The statement ($\times \times$) follows from this and from the relation $H(x) \leq \mathcal{V}(\{ \langle x, 0 \rangle \}) = G^*(x) = 2^{G(x)}$. We are proving (b). Let $x \in A$ be such that $H(x) \leq 2^\alpha$ and $x \in B$. We have $2^{G(x)-1} \leq H(x) \leq 2^\alpha$ and, consequently $x \in Q(\alpha+1)$ holds. Conversely, let $x \in Q(\alpha)-B$. We have $G(x) \leq \alpha$. We deduce from this that $H(x) \leq 2^{G(x)} \leq 2^\alpha$.

(c) $H(F(x,y)) \leq H(x) + H(y)$ holds for each $x, y \in A$. This follows immediately from the construction of H .

(d) $H(B(x)) = H(x)$ holds for each $x \in A$.

We shall prove (d) by using the following lemma.

- Lemma 5. Let t be a path in A , $\text{dom}(t) = \nu^{\beta} + 1$, and let $\alpha \leq \beta \leq \nu^{\beta}$. (1) $\mathcal{V}(E \circ t) \leq \mathcal{V}(t)$.
 (2) If \mathcal{Q} is covariant then $[FJ](E \circ t, \alpha, \beta) = E([FJ](t, \alpha, \beta))$.
 (3) If \mathcal{Q} is contravariant then $[FJ](E \circ \bar{t}, \alpha, \beta) = E([FJ](t, \mathcal{V} - \beta, \mathcal{V} - \alpha))$.

The proof of this lemma is straightforward and we omit it.

We prove that

$$(\square) \quad H(y) \leq H(E(y))$$

holds for each $y \in A$. Suppose that $E(y) = x$. Let t be a path in A such that $[FJ](t) = x$ and $\mathcal{V}(t) = H(x)$. Assume covariant \mathcal{Q} . Then $[FJ](E \circ t) = E([FJ](t)) = E(x) = y$. Assume contravariant \mathcal{Q} . Then $[FJ](E \circ \bar{t}) = E([FJ](t)) = E(x) = y$. We have $\mathcal{V}(E \circ \bar{t}) \leq \mathcal{V}(E \circ t) \leq \mathcal{V}(t) = H(x)$ and, consequently, (\square) is proved. We deduce from (\square) that

$$H(y) \leq H(E(y)) \leq H(E(E(y))) = H(y).$$

Thus, the statement (d) is proved. The proof of the \mathcal{G} -valuation lemma is finished.

3.0.3. Remark. (1) The valuation H from the previous proof is defined as follows: $\langle x, y \rangle \in H \equiv y \in A \ \& \ x = \min \{ \mathcal{V}_{\mathcal{Q}}(t); [FJ](t) = x \}$. Thus, there is a normal formula $\Phi'(x, y, X, Y)$ of the language FL such that

$$\langle x, y \rangle \in H \equiv \Phi'(x, y, \mathcal{Q}, \mathcal{V}_{\mathcal{Q}}).$$

The function $\mathcal{V}_{\mathcal{Q}}$ is constructed by a normal formula again.

We deduce from this that there exists a normal formula

$\Phi(x, y, X, Y)$ of the language FL, satisfying

$$\langle x, y \rangle \in H \equiv \Phi(x, y, \mathcal{Q}, \mathcal{Q}).$$

(2) Let \mathcal{Q}, \mathcal{R} be \mathcal{G} -artings in \mathcal{Q} over B , where B is an universe in an e -structure $\mathcal{A} = \langle A, F, E \rangle$. Let $\text{dom}(\mathcal{Q}) = \text{dom}(\mathcal{R})$

and suppose that $Q(\alpha) \subseteq R(\alpha)$ holds for each $\alpha \in \text{dom}(Q)$. We put

$$H^Q = \{ \langle x, y \rangle; \Phi(x, y, A, Q) \}, \quad H^R = \{ \langle x, y \rangle; \Phi(x, y, A, R) \}.$$

Then $H^R(x) \subseteq H^Q(x)$ holds for each $x \in A$.

Proof. Let x be an element of A . Then $G_R(x) \subseteq G_Q(x)$. (For G_Q see the previous proof.) We deduce from this that $\mathcal{V}_R(t) \subseteq \mathcal{V}_Q(t)$ for each path t in A . The required proposition follows from this immediately.

§ 4. Scales for $\sigma^{\mathcal{M}}$ -triads and $\pi^{\mathcal{M}}$ -triads

4.0.0. A triad \mathcal{T} is called scale for the type $\sigma^{\mathcal{M}}$ ($\pi^{\mathcal{M}}$ resp.) iff \mathcal{T} is a σ^0 (π^0 resp.)-triad and, for each triad $\tilde{\mathcal{T}}$ of the type $\sigma^{\mathcal{M}}$ ($\pi^{\mathcal{M}}$ resp.), there exists a valuation H of $\tilde{\mathcal{T}}$ in \mathcal{T} such that $H \in \mathcal{M}$.

4.0.1. Theorem

(1) The triad $\langle N, +, \text{Id} \rangle$ ($\text{FN}, \{0\}$) is a scale for the type $\sigma^{\mathcal{M}}$.

(2) The triad $\langle \text{RN}(\geq 0), +, \text{Id} \rangle$ ($\{ \geq 0 \}, \{0\}$) is a scale for the type $\pi^{\mathcal{M}}$.

Proof. Let $\mathcal{A} = \langle A, F, E \rangle$ be an e-structure and let $\mathcal{A}(Q, B)$ be a $\sigma^{\mathcal{M}}$ -triad over \mathcal{A} . We have $\llbracket F, E \rrbracket(Q, Q)$. Thus, there is a σ -string S of Q , $S \in \mathcal{M}$, and $B \subseteq S(0) \subseteq S(\alpha) \subseteq A$, $\llbracket F, E \rrbracket(S(\alpha), S(\alpha+1))$ holds for each $\alpha+1 \in \text{dom}(S)$. (This follows from [M1] 2.1.0). Put, for each $\alpha \in \text{dom}(S)$,

$$\langle x, \alpha \rangle \in P \equiv x \in S(\alpha) \cap E^*S(\alpha)$$

We deduce from 2.0.3 that P is a σ -string of Q and $B \subseteq P(0) \subseteq P(\alpha) \subseteq A$, $F^*P^2(\alpha) \subseteq P(\alpha+1)$, $E^*P(\alpha) \subseteq P(\alpha)$ hold for each $\alpha+1 \in \text{dom}(P)$. Evidently, P is an element of \mathcal{M} . Let $\tilde{\sigma} \in N\text{-FN}$ be such that $2\tilde{\sigma} \subset \text{dom}(P)$. Let R be a relation, satis-

fying: $\text{dom}(R) = \sigma + 1$, $R^{\{0\}} = B$, $R^{\{\sigma\}} = A$, $1 \leq \alpha < \sigma \rightarrow$
 $\rightarrow R^{\{\alpha\}} = P(2\alpha)$. It is easy that $R \in \mathcal{M}$ and R is a σ -
 string of Q . Moreover, R is a σ -string in \mathcal{A} over B . We de-
 duce from the σ -valuation lemma that there is a valuation
 $H \in \mathcal{M}$ of $\mathcal{A}(B, B)$ in $\langle N, +, \text{Id} \rangle (\{0\}, \{0\})$ and $x \in Q \equiv (\exists n)$
 $(H(x) \leq 2^n)$ holds. Consequently, H is a valuation of $\mathcal{A}(Q, B)$
 in $\langle N, +, \text{Id} \rangle (FN, \{0\})$ and the part (1) of the theorem is pro-
 ved. The part (2) can be proved quite analogously as the part
 (1).

4.0.2. Remark. Let $\mathcal{A}(Q, B)$ be a triad and suppose that
 $\mathcal{A} \in \text{Sd}_V$, $B \in \text{Sd}_V$. Assume that Q is a σ -class which is not a
 σ^0 -class. Then there exists a valuation H of $\mathcal{A}(Q, B)$ in
 $\langle N, +, \text{Id} \rangle (FN, \{0\})$ and $H \in \text{Sd}_V^*$. But no valuation of $\mathcal{A}(Q, B)$ in
 $\langle N, +, \text{Id} \rangle (FN, \{0\})$ is an element of Sd_V .

Proof. The existence of a valuation, which is a Sd_V^* -
 class, follows from the previous theorem (because $\mathcal{A}(Q, B)$ is
 a $\sigma^{\text{Sd}_V^*}$ -triad).

Suppose that there is a valuation of $\mathcal{A}(Q, B)$ in
 $\langle N, +, \text{Id} \rangle (FN, \{0\})$ and let $H \in \text{Sd}_V$. Let $\xi \in N - FN$. Then $R =$
 $= \{ \langle x, \alpha \rangle ; H(x) < \alpha \ \& \ \alpha \in \xi \}$ is a σ -string of Q and $R \in \text{Sd}_V$.
 Thus Q is a σ^0 -class, which is a contradiction.

4.1.0. Let Q be an equivalence on a class A . The map-
 ping $H: A^2 \rightarrow RN(\geq 0)$ is called metric of Q on A iff the fol-
 lowing holds for each $x, y, z \in A$:
 $H(x, z) \leq H(x, y) + H(y, z)$, $H(x, y) = H(y, x)$, $H(x, y) \neq 0 \equiv \langle x, y \rangle \in Q$,
 $H(x, y) = 0 \equiv x = y$.

Metrization theorem. Let Q be an equivalence on A ,
 $A \in \mathcal{M}$, and let Q be a $\sigma^{\mathcal{M}}$ -class. Then there exists a met-
 ric H of Q on A , $H \in \mathcal{M}$.

Proof. Let $E^0: V^2 \cup \{0\} \rightarrow V^2 \cup \{0\}$ be the mapping defined as follows: $E^0(\langle x, y \rangle) = \langle y, x \rangle$, $E^0(0) = 0$. Then $\mathcal{A} = \langle A^2 \cup \{0\}, F^0, E^0 \rangle$ is a contravariant e-structure and $\mathcal{T} = \langle \mathcal{A}(\mathbb{Q} \cup \{0\}, \{\langle x, x \rangle; x \in A\} \cup \{0\}) \rangle$ is a \mathcal{A} -triad. Let $G \in \mathcal{M}$ be a valuation of \mathcal{T} in $\langle \mathbb{R}N(\geq 0), +, Id \rangle (\mathbb{I} \geq 0, \{0\})$. A metric in question is the mapping $G \wedge A^2$.

Corollary. (1) There exists a metric H of $\underline{\cong}$ on V , so that $H \in Sd_V^*$.

(2) There is no metric of $\underline{\cong}$ on V which is an element of Sd_V .

Proof. (1) follows from the metrization theorem. (2) follows from [M1], 1.0.7 and from 4.0.2. (For the equivalence $\underline{\cong}$ see also § 0.)

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