

Václav Koubek; Jan Reiterman  
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ON THE CATEGORY OF FILTERS

Václav KOUBEK, Jan REITERMAN, Praha

In the present note the category of filters is studied. Denote by  $F_0$  the category, the objects of which are ordered pairs  $[A, \mathcal{F}]$ , where  $A$  is a set and  $\mathcal{F}$  a filter on  $A$ . The morphisms from  $[A, \mathcal{F}]$  to  $[B, \mathcal{G}]$  are all mappings  $\alpha: A \rightarrow B$  with  $\alpha^{-1}(G) \in \mathcal{F}$  for every  $G \in \mathcal{G}$ . Denote by  $F$  the category which we obtained from  $F_0$  by identifications of those mappings  $\alpha, \alpha'$  which are equal on a set  $F \in \mathcal{F}$ . Exact definition c.f. below. The note has four parts. The first contains the basic conventions and exact definition of the category  $F_*$ . The second part contains the characterization of epimorphism and monomorphisms in  $F$ . In the third part the concretizability of the category  $F$  is proved. The fourth part contains some examples of categories the concretizability of which follows immediately from the concretizability of the category  $F$ .

1. Conventions from the set theory

If  $A, B$  are sets,  $f$  a mapping  $f: A \rightarrow B$ , and  $C$  a subset of  $A$  then  $f/C$  denotes the restriction of  $f$  to the domain  $C$ .

If  $A, B$  are sets and  $b_a$  is given for every  $a \in A$ , then the set of all  $b_a, a \in A$  is denoted by  $\{b_a; a \in A\}$ ; the mapping  $a \rightarrow b_a$  is denoted by  $\{b_a | a \in A\}$ .

Conventions from the category theory. If  $K$  is a category, then  $K^\sigma$  denotes the class of all its objects and  $K^m$  the class of all its morphisms. If  $a, b \in K^\sigma$  then  $K(a, b)$  denotes the set of all morphisms from  $a$  into  $b$ . If

$a, b, c \in K^\sigma, f \in K(a, b), g \in K(b, c)$ , then the composition of  $f$  and  $g$  is denoted by  $g \circ f$ .

We recall the following definition: A category  $K$  is said to be concretizable if and only if there exists an isofunctor from  $K$  into  $S$ , where  $S$  is the category of all sets and their mappings. It is well known that a category  $K$  is concretizable if and only if there exists a faithful functor from  $K$  into  $S$ .

Definition of the category  $F$ . Let  $\tilde{C}$  be the class of all ordered pairs  $[A, \mathcal{F}]$ , where  $A$  is a set and  $\mathcal{F}$  is a filter on  $A$ . A triple  $\langle \mathcal{F}, \mathcal{G}, \alpha \rangle$  will be called a morphism from  $[A, \mathcal{F}]$  into  $[B, \mathcal{G}]$  if and only if  $\alpha$  is a mapping,  $\alpha: A \rightarrow B$  such that

$$G \in \mathcal{G} \Rightarrow \alpha^{-1}(G) \in \mathcal{F}.$$

We define composition of two morphisms as follows:

$$\langle \mathcal{G}, \mathcal{H}, \beta \rangle \circ \langle \mathcal{F}, \mathcal{G}, \alpha \rangle = \langle \mathcal{F}, \mathcal{H}, \beta \circ \alpha \rangle.$$

Denote by  $F_0$  the category such that  $F_0^\sigma = \tilde{C}$  and  $F_0^m$  is the class of all morphisms described above with the composition defined above. We define an equivalence on  $F$  as follows:

$$\langle \mathcal{F}_1, \mathcal{G}_1, \alpha_1 \rangle \sim \langle \mathcal{F}_2, \mathcal{G}_2, \alpha_2 \rangle \equiv (\mathcal{F}_1 = \mathcal{F}_2) \& \\ \& (\mathcal{G}_1 = \mathcal{G}_2) \& (\exists F \in \mathcal{F}_1) (\alpha_1 / F = \alpha_2 / F) .$$

It is easy to see that  $\sim$  is a congruence on  $F_0^m$  and consequently it defines a factorcategory  $F$ , morphisms of which are equivalence-classes of morphisms of  $F_0$  with respect to  $\sim$ . We shall denote the morphisms of the category  $F$  by  $f, g, h \dots$ .

We shall write  $\alpha \in f$ , whenever  $\langle \mathcal{F}, \mathcal{G}, \alpha \rangle \in f$  and we shall say that the mapping  $\alpha$  designates the morphism  $f$ .

2.

Lemma 1: A morphism  $f \in F([A, \mathcal{F}], [B, \mathcal{G}])$  is an epimorphism if and only if the following holds:

$$(1) \quad (\forall \alpha \in f) (\forall F \in \mathcal{F}) (\alpha(F) \in \mathcal{G}) .$$

Remark: The condition (1) is equivalent to the condition (1')

$$(1') \quad (\exists \alpha \in f) (\forall F \in \mathcal{F}) (\alpha(F) \in \mathcal{G}) .$$

Proof of the remark is evident.

Proof of Lemma 1: Let us assume that the condition holds and  $f$  is not an epimorphism, i.e.

$$(\exists [C, \mathcal{H}] \in F^\sigma) (\exists g, h \in F([B, \mathcal{G}], [C, \mathcal{H}])) (g + h, g \circ f = h \circ f).$$

The last equality implies

$$(\forall \alpha \in \mathfrak{f})(\forall \beta \in \mathfrak{g})(\forall \gamma \in \mathfrak{h})(\exists F \in \mathfrak{F})(\beta \circ \alpha / F = \gamma \circ \alpha / F).$$

It means that  $\beta / \alpha (F) = \gamma / \alpha (F)$ , consequently  $\mathfrak{h} = \mathfrak{g}$  which is a contradiction.

Let us assume that the condition (1) does not hold. Then there exists  $F \in \mathfrak{F}$  such that  $\alpha (F) \notin \mathfrak{g}$ . On the other hand the set  $B - \alpha (F)$  is not a member of  $\mathfrak{G}$  because  $(\alpha^{-1}(B - \alpha (F))) \cap F = \emptyset$ .

Denote:

$$\mathfrak{G}_1 = \{G \cap \alpha (F); G \in \mathfrak{G}\},$$

$$\mathfrak{G}_2 = \{G \cap (B - \alpha (F)); G \in \mathfrak{G}\}.$$

It is easy to see that  $\mathfrak{G}_1$  (or  $\mathfrak{G}_2$ ) is a filter on a set  $\alpha (F)$  (or  $B - \alpha (F)$  respectively). Let  $C = C_1 \cup C_2 \cup C_3$ , where  $C_i$  are disjoint sets such that

$$\text{card } C_1 = \text{card } \alpha (F),$$

$$\text{card } C_2 = \text{card } C_3 = \text{card } (B - \alpha (F)).$$

Let  $\omega: \alpha (F) \rightarrow C_1, \pi_1: (B - \alpha (F)) \rightarrow C_2, \pi_2: (B - \alpha (F)) \rightarrow C_3$  be arbitrary bijective mappings.

Define the filter  $\mathfrak{H}$  on the set  $C$  as follows:

$$(X \in \mathfrak{H}) \equiv (\omega^{-1}(X \cap C_1) \in \mathfrak{G}_1 \& \pi_1^{-1}(X \cap C_2) \in \mathfrak{G}_2 \& \pi_2^{-1}(X \cap C_3) \in \mathfrak{G}_2).$$

The mappings  $\varepsilon, \mu: B \rightarrow C$  defined by

$$\varepsilon / \alpha (F) = \mu / \alpha (F) = \omega, \varepsilon / (B - \alpha (F)) = \pi_1, \mu / (B - \alpha (F)) = \pi_2$$

designate the morphisms  $g, h$  such that  $g \neq h$ ,

$$g \circ f = h \circ f.$$

Consequently,  $f$  is not an epimorphism.

Lemma 2: A morphism  $f \in \mathbb{F}([A, \mathcal{F}], [B, \mathcal{G}])$  is a monomorphism if and only if the following holds:

$$(2) (\forall \alpha \in f) (\exists F \in \mathcal{F}) (\forall x, y \in F) (x \neq y \Rightarrow \alpha(x) \neq \alpha(y)).$$

Remark: The condition (2) is equivalent to the condition (2')

$$(2') (\exists \alpha \in f) (\exists F \in \mathcal{F}) (\forall x, y \in F) (x \neq y \Rightarrow \alpha(x) \neq \alpha(y)).$$

Proof of the remark is evident.

Proof of Lemma 2: Clearly, if (2) is satisfied then  $f$  is a monomorphism. Let us assume that the condition (2) does not hold, i.e.

$$(\exists \alpha \in f) (\forall F \in \mathcal{F}) (\exists a_F, b_F \in F) (a_F \neq b_F \& \alpha(a_F) = \alpha(b_F)).$$

Put  $C = \{[a_F, b_F]; F \in \mathcal{F}\}$ . Let  $\mathcal{H}$  be a filter on the set  $C$  a base of which is the set of all

$$\{[a_F, b_F]; F \subset G\}, \text{ where } G \in \mathcal{F}.$$

The mappings  $\varepsilon, \mu: C \rightarrow A$  defined by

$$\varepsilon([a_F, b_F]) = a_F, \quad \mu([a_F, b_F]) = b_F$$

designate the morphisms  $g, h$  of  $[C, \mathcal{H}]$  into  $[A, \mathcal{F}]$  such that

$$g \neq h, \quad f \circ g = f \circ h.$$

Consequently, the morphism  $f$  is not a monomorphism.

3.

Definition: Denote by  $\mathcal{U}$  the full subcategory of  $\mathbb{F}$  the objects of which are all  $[A, \mathcal{F}]$  where  $\mathcal{F}$  is an ultrafilter.

Convention: Let  $\mathcal{T}$  be the class of all cardinal numbers. For every  $t \in \mathcal{T}$  choose a set  $X_t$  with  $\text{card } X_t = t$ . The sets  $X_t$  will be fixed in the sequel.

Definition: For every object  $[A, \mathcal{F}] \in \mathcal{F}^\sigma$  put  $\min_{F \in \mathcal{F}} \text{card } F = \|[A, \mathcal{F}]\|$ . The number  $\|[A, \mathcal{F}]\|$  will be called essential cardinality of the filter  $\mathcal{F}$ .

Lemma 3: There exists a skeleton  $\mathcal{U}_1$  of  $\mathcal{U}$  with the following property: if  $[A, \mathcal{F}] \in \mathcal{U}_1^\sigma$  then

$$A = X_{\|[A, \mathcal{F}]\|}$$

Proof is evident.

Lemma 4: The category  $\mathcal{U}$  is concretizable.

Proof: It is sufficient to prove that  $\mathcal{U}_1$  is concretizable.

1) First we prove that:

$$[X_t, \mathcal{F}], [X_u, \mathcal{G}] \in \mathcal{U}_1^\sigma; t < u \Rightarrow \mathcal{U}_1([X_t, \mathcal{F}], [X_u, \mathcal{G}]) = \emptyset.$$

Assume that there exist  $f \in \mathcal{U}_1([X_t, \mathcal{F}], [X_u, \mathcal{G}])$ .

If  $\alpha \in f$ ,  $F \in \mathcal{F}$ , then  $\alpha(F) \in \mathcal{G}$ . For,  $\mathcal{G}$  is an ultrafilter and  $\alpha^{-1}(X_u - \alpha(F)) \cap F = \emptyset$ . Thus,  $\text{card } \alpha(F) = u$  while  $\text{card } F = t < u$ . That is a contradiction.

2) Consequently,

$$\bigcup_{b \in \mathcal{U}_1^\sigma} \mathcal{U}_1(a, b) = \bigcup_{b \in \mathcal{U}_1^\sigma, \|b\| \leq \|a\|} \mathcal{U}_1(a, b).$$

The right side hand is evidently a set, which implies that  $\mathcal{U}_1$  is concretizable because we can use the

Mac-Lane's representation for the category  $\mathcal{U}_1^*$  dual to  $\mathcal{U}_1$ .

Definition: Let  $\mathbb{K}$  be arbitrary category. Define the category  $H^{\mathbb{K}}$  as follows. The object of the category  $H^{\mathbb{K}}$  are all sets of objects of the category  $\mathbb{K}$ . Let  $a, b$  be the objects of the category  $H^{\mathbb{K}}$ . Morphisms from  $a$  to  $b$  are exactly all collections  $\{f_m \mid m \in a\}$  where  $f_m \in \mathbb{K}(m, N_m)$ ,  $N_m \in b$ . We define the composition:

$$\{g_m \mid m \in b\} \circ \{f_m \mid m \in a\} = \{g_{N_m} \circ f_m \mid m \in a\}.$$

Remark: It is evident that  $H^{\mathbb{K}}$  is a category.

Lemma 5: If the category  $\mathbb{K}$  is concretizable then the category  $H^{\mathbb{K}}$  is concretizable.

Proof is evident.

Theorem: The category  $\mathbb{F}$  is concretizable.

Proof: 1) The category  $H^{\mathcal{U}}$  is concretizable.

2) Now we shall construct a functor  $\Psi: \mathbb{F} \rightarrow H^{\mathcal{U}}$ . For every  $[A, \mathcal{F}] \in \mathbb{F}^{\sigma}$  define  $\Psi[A, \mathcal{F}]$  as the set of all  $[A, \mathcal{H}]$ , where  $\mathcal{H}$  is an ultrafilter on  $A$  and  $\mathcal{F} \subset \mathcal{H}$  (i.e.  $F \in \mathcal{F} \Rightarrow F \in \mathcal{H}$ ). If

$$f \in \mathbb{F}([A, \mathcal{F}], [B, \mathcal{G}]), \alpha \in f, [A, \mathcal{H}] \in \Psi[A, \mathcal{F}],$$

then the set  $\{\alpha(H); H \in \mathcal{H}\}$  is a base of an ultrafilter on  $B$  which will be called  $f(\mathcal{H})$ . (The ultrafilter  $f(\mathcal{H})$  does not depend on a choice of  $\alpha \in f$ .) Define:

$$\Psi(f) = \{f_{[A, \mathcal{H}]} \mid [A, \mathcal{H}] \in \Psi[A, \mathcal{F}]\},$$



where  $f_{[A, \mathcal{A}]} \in \mathcal{U}([A, \mathcal{A}], [B, f(\mathcal{A})])$  such that  $\alpha \in f_{[A, \mathcal{A}]}$  whenever  $\alpha$  is a mapping  $\alpha: A \rightarrow B$  with  $\alpha \in f$ .

3) Now we prove that  $\Psi$  is an isofunctor from  $\mathbb{F}$  into  $H^{\mathcal{U}}$ . The mapping  $\Psi/F^{\sigma}$  is one-to-one because

$$\mathcal{F} = \bigcap_{[A, \mathcal{A}] \in \Psi[A, \mathcal{F}]} \mathcal{A}$$

for each filter  $\mathcal{F}$  on  $A$ . We shall prove that for each  $a, b \in F^{\sigma}$ ,  $a = [A, \mathcal{F}], b = [B, \mathcal{G}], \Psi_{/F}(a, b)$  is one-to-one. Let  $f, g$  be two morphisms from  $a$  to  $b$ ,  $f \neq g$ . Choose  $\alpha \in f, \beta \in g$  and set

$$C = \{x \in A; \alpha(x) \neq \beta(x)\}.$$

Since  $f \neq g$ ,  $C \cap F \neq \emptyset$  holds for each  $F \in \mathcal{F}$ . Consequently,  $\{C \cap F; F \in \mathcal{F}\}$  is a base of a filter  $\mathcal{O}$  on  $A$ . Let  $\mathcal{H}$  be an ultrafilter on  $A$  with  $\mathcal{H} \supset \mathcal{O}$ . Since  $\mathcal{O} \supset \mathcal{F}$ ,  $\mathcal{H} \in \Psi[A, \mathcal{F}]$ , it is easy to see that  $H \cap C \neq \emptyset$  for every  $H \in \mathcal{H}$ . Therefore  $f_{[A, \mathcal{H}]} \neq g_{[A, \mathcal{H}]}$ , consequently  $\Psi(f) \neq \Psi(g)$ .

4) The assertion of the theorem follows now immediately from 3) and 1).

#### 4. Some examples

1) We recall: a directed set is an ordered pair  $[A, \kappa]$ , where  $A$  is a set and  $\kappa$  a partial order on  $A$  such that

$$(\forall a \in A)(\forall b \in A)(\exists c \in A)(a \kappa c \& b \kappa c).$$

Let  $[A_1, \kappa_1], [A_2, \kappa_2]$  be two directed sets. A triple  $\langle \kappa_1, \kappa_2, \alpha \rangle$  will be called a morphism from  $[A_1, \kappa_1]$  into  $[A_2, \kappa_2]$  if and only if  $\alpha$  is a  $\kappa_1 - \kappa_2$  compatible mapping,  $\alpha: A_1 \rightarrow A_2$ , i.e.  $\alpha$  is a mapping from  $A_1$  into  $A_2$  such that

$$a, b \in A_1, a \kappa_1 b \Rightarrow \alpha(a) \kappa_2 \alpha(b).$$

We define the composition of two morphisms as follows:

$$\langle \kappa_2, \kappa_3, \beta \rangle \circ \langle \kappa_1, \kappa_2, \alpha \rangle = \langle \kappa_1, \kappa_3, \beta \circ \alpha \rangle.$$

It is clear that directed sets as objects with morphisms just described form a category. Denote this category by  $\mathbb{R}_0$ . Denote by  $\mathbb{R}$  the factorcategory of  $\mathbb{R}_0$  with respect to the congruence  $\sim$  where  $\sim$  is defined as follows:

$$\begin{aligned} &\langle \kappa_1, \kappa_2, \alpha_i \rangle \in \mathbb{R}_0 ([A_1, \kappa_1], [A_2, \kappa_2]), \quad i = 1, 2, \\ &\langle \kappa_1, \kappa_2, \alpha_1 \rangle \sim \langle \kappa_1, \kappa_2, \alpha_2 \rangle \equiv \\ &\equiv (\exists x \in A_1)(\forall y \in A_1)(x \kappa_1 y \Rightarrow \alpha_1(y) = \alpha_2(y)). \end{aligned}$$

2) Denote by  $\mathbb{P}$  the class of all triples  $[t, T, \mathcal{T}]$  where  $[T, \mathcal{T}]$  is a topological space and  $t \in T$ . A continuous mapping  $f$  from  $[T, \mathcal{T}]$  into  $[S, \mathcal{S}]$  will be called a morphism from  $[t, T, \mathcal{T}]$  into  $[s, S, \mathcal{S}]$  if and only if  $f(t) = s$ . The composition of morphisms

is the usual composition of mappings. Clearly, elements of  $\mathcal{P}$  as objects and morphisms just described form a category. Denote by  $\mathcal{T}_0$  this category. Denote by  $\mathcal{T}$  the factorcategory of  $\mathcal{T}_0$  with respect to the congruence  $\sim$ , where  $\sim$  is defined as follows:

$$\alpha, \beta \in \mathcal{T}_0 ([t, \mathcal{T}, \mathcal{I}], [s, \mathcal{S}, \mathcal{Y}]) ,$$

$$\alpha \sim \beta \equiv (\exists U \in \mathcal{U}_t^{\mathcal{T}})(\alpha/U = \beta/U) .$$

(  $\mathcal{U}_t^{\mathcal{T}}$  denote the system of all neighborhoods of the point  $t$  in the topology  $\mathcal{T}$  . )

3) Let  $\mathcal{Q}$  be the class of all ordered pairs  $[M, \mu]$ , where  $M$  is a set and  $\mu$  a non-trivial measure on  $M$ . If  $[M, \mu] \in \mathcal{Q}$ , let us denote by  $\mathcal{D}\mu$  (or  $\mathcal{D}_0\mu$ ) the system of all  $\mu$ -measurable sets (or the system of all  $N \subset M$  such that  $\mu(N) = 0$ , respectively). A mapping  $\alpha: M_1 \rightarrow M_2$  will be called a morphism from  $[M_1, \mu_1]$  into  $[M_2, \mu_2]$  if and only if

$$(N \in \mathcal{D}\mu_2 \Rightarrow \alpha^{-1}(N) \in \mathcal{D}\mu_1) \& (N \in \mathcal{D}_0\mu_2 \Rightarrow \alpha^{-1}(N) \in \mathcal{D}_0\mu_1) .$$

The composition of morphisms is the usual composition of mappings. It is easy to see that elements of  $\mathcal{Q}$  and morphisms just described form a category. Denote this category by  $|\mathcal{M}|_0$ . Denote by  $|\mathcal{M}|$  the functorcategory of  $|\mathcal{M}|_0$  with respect to congruence  $\sim$ , where  $\sim$  is defined as follows:

$$[M, \mu], [N, \nu] \in \mathcal{Q}, \alpha, \beta \in |\mathcal{M}|_0([M, \mu], [N, \nu]),$$

$(\alpha \sim \beta) \equiv (\alpha = \beta \text{ } \mu\text{-almost everywhere}) .$

Proposition: The categories  $\mathcal{R}$ ,  $\mathcal{T}$ ,  $\mathcal{M}$  are concretizable. It follows almost immediately from the fact that the category  $\mathcal{F}$  is concretizable. The categories  $\mathcal{R}$ ,  $\mathcal{T}$ ,  $\mathcal{M}$  can be represented as subcategories of the category  $\mathcal{F}$  .

Matematicko-fyzikální fakulta  
Karlovy university  
Praha 8 Karlín  
Sokolovská 83, Czechoslovakia

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