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Czechoslovak Mathematical Journal, Vol. 20 (1970), No. 1, 109–111

Persistent URL: <http://dml.cz/dmlcz/100949>

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A REMARK ON A PROBLEM OF HARARY

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(Received March 13, 1969)

F. HARARY [1] publicized the following question:

“For any graph G with p points, how can levels $1, 2, \dots, p$ be assigned to the points in order to minimize the maximum of the absolute value of the differences between the levels of all pairs of adjacent points?”

We set

$$\varphi(G) = \min \max |f(u) - f(v)|$$

where the maximum is taken over all edges (lines) uv and the minimum over all $1-1$ mappings (valuations) $f: V(G) \rightarrow \{1, 2, \dots, p\}$; by $V(G)$ we mean a set of all vertices (points) of G .

Otherwise, our notation follows Harary [2]. Particularly, we reserve: a letter p for number of points, q for number of lines, d for diameter, κ for connectivity, β_0 for point independence number, d_i for degrees of points ($d_1 \leq d_2 \leq \dots \leq d_p$).

As in [3], we write P_p^k for the k^{th} power of the path P_p , in which two points u, v of $V(P_p) = \{1, 2, \dots, p\}$ are adjacent if, and only if, $0 < |u - v| \leq k$.

Theorem 1. $\varphi(G)$ is the smallest integer k such that $G \subset P_p^k$.

Proof. Every inclusion $f: G \rightarrow P_p^k$ induces a valuation $f: V(G) \rightarrow \{1, 2, \dots, p\}$ such that $\max |f(u) - f(v)| \leq k$ and vice versa. Hence, $\varphi(G) \leq k$ if and only if $G \subset P_p^k$, q.e.d.

Theorem 1 reduces – from a theoretical viewpoint – the original question to an elementary problem of graph theory: Given a pair of graphs, is one of them a subgraph of the other? In practice, however, an answer to the last question becomes extremely difficult, even with aid of high-speed computers. Therefore, it does not seem likely that one could find an effective algorithm to determine a minimal valuation of G . Nevertheless, Theorem 1 will enable us to determine some relations between φ

and the other fundamental invariants of graphs. For this purpose, we shall need the following lemma; its proof is straightforward and will be omitted.

Lemma. *The invariants of the graph P_p^k satisfy:*

$$q = (p - 1) + (p - 2) + \dots + (p - k) = k \frac{2p - k - 1}{2},$$

$$d = \left\{ \frac{p - 1}{k} \right\}, \quad \varkappa = k, \quad \beta_0 = 1 + \left[\frac{p - 1}{k + 1} \right] = \left\{ \frac{p}{k + 1} \right\},$$

$$d_j = \min \left(p - 1, k + \left[\frac{j - 1}{2} \right], 2k \right).$$

Theorem 2. *For any graph G , we have*

$$\varphi \geq p - \frac{1 + \sqrt{((2p - 1)^2 - 8q)}}{2},$$

$$\varphi \geq \frac{p - 1}{d}, \quad \varphi \geq \varkappa, \quad \varphi \geq \frac{p}{\beta_0} - 1,$$

$$\varphi \geq \max_j \max \left(d_j - \left[\frac{j - 1}{2} \right], \frac{d_j}{2} \right).$$

Proof. By Theorem 1, $G \subset P_k^p$. Now, observe that $p(G_1) = p(G_2)$, $G_1 \subset G_2$ implies $q(G_1) \leq q(G_2)$, $d(G_1) \geq d(G_2)$, $\varkappa(G_1) \leq \varkappa(G_2)$, $\beta_0(G_1) \geq \beta_0(G_2)$ and $d_j(G_1) \leq d_j(G_2)$ for each j . The rest follows by our Lemma.

Theorem 3. *If $m \geq n > 0$ then $\varphi(K_{mn}) = [(m - 1)/2] + n$ and a valuation $f: V(K_{mn}) \rightarrow \{1, 2, \dots, m + n\}$ is minimal whenever*

$$f(M) = \left\{ 1, 2, \dots, \left[\frac{m}{2} \right] \right\} \cup \left\{ \left[\frac{m}{2} \right] + n + 1, \dots, m + n \right\}$$

where M is the independent set of K_{mn} having m points.

Proof. If $1 \in f(M)$, $p \notin f(M)$ ($p = m + n$ here) then $\max |f(u) - f(v)| = p - 1 > \varphi$ and f is not minimal (similarly for $1 \notin f(M)$, $p \in f(N)$). Therefore a minimal valuation f satisfies either

$$(i) \quad 1 \in f(M), \quad p \in f(M)$$

or

$$(ii) \quad 1 \notin f(M), \quad p \notin f(M).$$

In case (i), observe that

$$\max |f(u) - f(v)| = \max (p - \min f(N), \max f(N) - 1).$$

Now, obviously, $f(N) = \{j + 1, j + 2, \dots, j + n\}$ for a minimal f . A simple computation supplies the best values

$$j = \left\lfloor \frac{p - n}{2} \right\rfloor = \left\lfloor \frac{m}{2} \right\rfloor, \quad \max |f(u) - f(v)| = \left\lceil \frac{m + 1}{2} \right\rceil + n - 1.$$

In case (ii), we get similarly

$$\max |f(u) - f(v)| = \left\lceil \frac{n + 1}{2} \right\rceil + m - 1$$

and the assumption $m \geq n$ is in favour of (i), q.e.d.

I thank Professor HARARY for his very valuable comments.

References

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- [3] *F. Harary, R. M. Karp and W. T. Tutte*: A Criterion for Planarity of the Square of a Graph, *J. Combinatorial Theory 2* (1967), 395–405.

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