

Acta Universitatis Palackianae Olomucensis. Facultas Rerum  
Naturalium. Mathematica

---

Ivan Chajda

Class preserving mappings of equivalence systems

*Acta Universitatis Palackianae Olomucensis. Facultas Rerum Naturalium. Mathematica*, Vol. 43 (2004), No. 1, 61--64

Persistent URL: <http://dml.cz/dmlcz/132935>

**Terms of use:**

© Palacký University Olomouc, Faculty of Science, 2004

Institute of Mathematics of the Academy of Sciences of the Czech Republic provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these *Terms of use*.



This paper has been digitized, optimized for electronic delivery and stamped with digital signature within the project *DML-CZ: The Czech Digital Mathematics Library* <http://project.dml.cz>



# Class Preserving Mappings of Equivalence Systems

IVAN CHAJDA

*Department of Algebra and Geometry, Faculty of Science,  
Palacký University, Tomkova 40, 779 00 Olomouc, Czech Republic  
e-mail: chajda@inf.upol.cz*

(Received February 6, 2004)

## Abstract

By an equivalence system is meant a couple  $\mathcal{A} = (A, \theta)$  where  $A$  is a non-void set and  $\theta$  is an equivalence on  $A$ . A mapping  $h$  of an equivalence system  $\mathcal{A}$  into  $\mathcal{B}$  is called a class preserving mapping if  $h([a]_\theta) = [h(a)]_{\theta'}$  for each  $a \in A$ . We will characterize class preserving mappings by means of permutability of  $\theta$  with the equivalence  $\Phi_h$  induced by  $h$ .

**Key words:** Equivalence relation, equivalence system, relational system, homomorphism, strong homomorphism, permuting equivalences.

**2000 Mathematics Subject Classification:** 08A02, 08A35, 03E02

For the basic concepts, the reader is referred to [1],[2],[3]. Let  $R$  and  $S$  be binary relations on a non-void set  $A$ . As usually, their *relational product* will be denoted by  $R \circ S$ , i.e.  $R \circ S = \{\langle a, b \rangle \in A^2; \exists c \in A \text{ with } \langle a, c \rangle \in R \text{ and } \langle c, b \rangle \in S\}$ . We will say that  $R, S$  *permute* (or they are *permutable*) if  $R \circ S = S \circ R$ .

**Lemma 1** *Let  $R, S$  be symmetric relations on  $A$ . Then  $R \circ S \subseteq S \circ R$  is equivalent to  $R \circ S = S \circ R$ .*

**Proof** If  $R \circ S \subseteq S \circ R$  then, due to symmetry,

$$S \circ R = S^{-1} \circ R^{-1} = (R \circ S)^{-1} \subseteq (S \circ R)^{-1} = R^{-1} \circ S^{-1} = R \circ S$$

thus  $S, R$  permute. The converse is trivial. □

By a *relational system* is meant a pair  $\mathcal{A} = (A, R)$ , where  $A \neq \emptyset$  is a set and  $R$  is a binary relation on  $A$ . If  $R$  is an equivalence relation,  $\mathcal{A} = (A, R)$  will be called an *equivalence system*.

We are going to introduce a quotient relational system as follows.

**Definition 1** Let  $\mathcal{A} = (A, R)$  be a relational system and  $\Phi$  be an equivalence on  $A$ . Define a binary relation  $R/\Phi$  on the factor set (i.e. a partition)  $A/\Phi$  as follows:  $\langle [a]_\Phi, [b]_\Phi \rangle \in R/\Phi$  iff there exist  $x \in [a]_\Phi, y \in [b]_\Phi$  with  $\langle x, y \rangle \in R$ . Then  $\mathcal{A}/\Phi = (A/\Phi, R/\Phi)$  will be called a *quotient relational system* of  $\mathcal{A}$  by  $\Phi$ .

**Remark 1** It is evident that if  $R$  is reflexive or symmetric then  $R/\Phi$  has the corresponding property.

**Lemma 2** Let  $\mathcal{A} = (A, R)$  be a relational system and  $R$  be transitive. Let  $\Phi$  be an equivalence on  $A$  and  $\Phi \circ R \subseteq R \circ \Phi$ . Then  $R/\Phi$  is transitive, too.

**Proof** Suppose  $\langle [a]_\Phi, [b]_\Phi \rangle \in R/\Phi$  and  $\langle [b]_\Phi, [c]_\Phi \rangle \in R/\Phi$ . Then there exist  $x, y, y', z \in A$  such that  $x \in [a]_\Phi, y, y' \in [b]_\Phi, z \in [c]_\Phi$  and  $\langle x, y \rangle \in R, \langle y', z \rangle \in R$ . Hence  $\langle x, z \rangle \in R \circ \Phi \circ R \subseteq R \circ R \circ \Phi \subseteq R \circ \Phi$ . Thus there is  $w \in A$  with  $\langle x, w \rangle \in R$  and  $\langle w, z \rangle \in \Phi$ , i.e.  $w \in [z]_\Phi = [c]_\Phi$ . By the Definition,  $\langle [a]_\Phi, [c]_\Phi \rangle \in R/\Phi$  proving transitivity of  $R/\Phi$ .  $\square$

Let  $\mathcal{A} = (A, R), \mathcal{B} = (B, S)$  be relational systems. A mapping  $h : A \rightarrow B$  is called a *homomorphism* of  $\mathcal{A}$  into  $\mathcal{B}$  if  $\langle a, b \rangle \in R$  implies  $\langle h(a), h(b) \rangle \in S$ .

A homomorphism  $h$  of  $\mathcal{A}$  into  $\mathcal{B}$  is called *strong* if for each  $\langle x, y \rangle \in S$  there exist  $a, b \in A$  such that  $\langle a, b \rangle \in R$  and  $h(a) = x, h(b) = y$ . Let  $\mathcal{A} = (A, \theta), \mathcal{B} = (B, \theta')$  be equivalence systems. A mapping  $h : A \rightarrow B$  is called *class preserving* if  $h([a]_\theta) = [h(a)]_{\theta'}$  for each  $a \in A$ .

**Lemma 3** Let  $\mathcal{A} = (A, \theta), \mathcal{B} = (B, \theta')$  be equivalence systems and  $h : A \rightarrow B$  a surjective class preserving mapping. Then  $h$  is a strong homomorphism of  $\mathcal{A}$  onto  $\mathcal{B}$ .

**Proof** It is evident that  $\langle a, b \rangle \in \theta$  implies  $\langle h(a), h(b) \rangle \in \theta'$ , i.e. it is a surjective homomorphism of  $\mathcal{A}$  onto  $\mathcal{B}$ . Suppose  $\langle c, d \rangle \in \theta'$ . Then there is  $a \in A$  with  $h(a) = c$  and  $d \in [c]_{\theta'} = [h(a)]_{\theta'}$ . Hence, there exists  $x \in [a]_\theta$  such that  $h(x) = d$ . Since  $\langle a, x \rangle \in \theta$ ,  $h$  is a strong homomorphism.  $\square$

**Example 1** The converse of Lemma 3 does not hold in general. Consider e.g.  $\mathcal{A} = (A, \theta), \mathcal{B} = (B, \theta')$  where  $A = \{x_1, x_2, y_1, y_2, z_1, z_2\}, B = \{a, b, c\}, \theta' = B \times B$  and  $\theta$  is determined by the partition  $\{x_1, x_2\}, \{y_1, y_2\}, \{z_1, z_2\}$ . Let  $h : A \rightarrow B$  is defined as follows:  $h(x_1) = h(y_1) = a, h(x_2) = h(z_1) = b, h(y_2) = h(z_2) = c$ . Then  $h$  is a surjective strong homomorphism of  $\mathcal{A}$  onto  $\mathcal{B}$  but it is not a class preserving mapping; e.g. for  $x_1$  we have

$$h([x_1]_\theta) = h(\{x_1, x_2\}) = \{a, b\} \neq \{a, b, c\} = [a]_{\theta'} = [h(x_1)]_{\theta'}.$$

**Theorem 1** Let  $\mathcal{A} = (A, \theta)$ ,  $\mathcal{B} = (B, \theta')$  be equivalence systems and  $h : A \rightarrow B$  a surjective mapping. The following are equivalent:

- (a)  $h$  is a class preserving mapping;
- (b)  $h$  is a homomorphism of  $\mathcal{A}$  onto  $\mathcal{B}$  and for each  $x, y \in A$  with  $\langle h(x), h(y) \rangle \in \theta'$  there exists  $z \in A$  such that  $\langle x, z \rangle \in \theta$  and  $h(y) = h(z)$ .

**Proof** (a)  $\Rightarrow$  (b) by Lemma 3 and its proof. Prove (b)  $\Rightarrow$  (a). Since  $h$  is a homomorphism, we easily get  $h([a]_\theta) \subseteq [h(a)]_{\theta'}$ . Suppose  $c \in [h(a)]_{\theta'}$ . Then  $c = h(w)$  for some  $w \in A$ . By (b) there exists  $z \in A$  such that  $\langle a, z \rangle \in \theta$  and  $h(z) = h(w) = c$ . Since  $z \in [a]_\theta$ , we conclude  $h([a]_\theta) = [h(a)]_{\theta'}$ .  $\square$

Let  $h : A \rightarrow B$  be a mapping. Denote by  $\Phi_h$  the so-called  $h$ -induced equivalence on  $A$ , i.e.

$$\langle x, y \rangle \in \Phi_h \quad \text{if and only if} \quad h(x) = h(y).$$

Let  $\Phi$  be an equivalence on  $A$ . Denote by  $h_\Phi$  the so-called *natural mapping*  $h_\Phi : A \rightarrow A/\Phi$  defined by  $h_\Phi(a) = [a]_\Phi$ .

**Theorem 2** Let  $\mathcal{A} = (A, \theta)$  be an equivalence system and  $\Phi$  be an equivalence on  $A$ . Suppose that  $\theta, \Phi$  permute. Then the natural mapping  $h_\Phi$  is a class preserving mapping of  $\mathcal{A}$  onto the quotient equivalence system  $\mathcal{A}/\Phi = (A/\Phi, \theta/\Phi)$ .

**Proof** By Lemma 2 and the previous Remark,  $A/\Phi$  is clearly a quotient equivalence system. Of course,  $h_\Phi$  is a surjective mapping. Suppose  $\langle a, b \rangle \in \theta$ . Then  $\langle [a]_\Phi, [b]_\Phi \rangle \in \theta/\Phi$  thus  $h_\Phi$  is a homomorphism of  $\mathcal{A}$  onto  $\mathcal{A}/\Phi$ . Let  $\langle [x]_\Phi, [y]_\Phi \rangle \in \theta/\Phi$ . Then there exist  $a \in [x]_\Phi, b \in [y]_\Phi$  such that  $\langle a, b \rangle \in \theta$ . Hence  $\langle x, b \rangle \in \Phi \circ \theta = \theta \circ \Phi$ , i.e. there exists  $z \in A$  such that  $\langle x, z \rangle \in \theta$  and  $\langle z, b \rangle \in \Phi$ , i.e.  $h_\Phi(z) = h_\Phi(b)$ . By (b) of Theorem 1,  $h_\Phi$  is a class preserving mapping.  $\square$

**Theorem 3** Let  $\mathcal{A} = (A, \theta)$ ,  $\mathcal{B} = (B, \theta')$  be equivalence systems and  $h : A \rightarrow B$  a surjective strong homomorphism of  $\mathcal{A}$  onto  $\mathcal{B}$ . Then  $h$  is a class preserving mapping if and only if  $\theta$  and the  $h$ -induced equivalence  $\Phi_h$  permute.

**Proof** Let  $h$  be a class preserving mapping and suppose  $\langle x, z \rangle \in \Phi_h \circ \theta$ . Then there exists  $y \in A$  with  $\langle x, y \rangle \in \Phi_h$  and  $\langle y, z \rangle \in \theta$ . Thus  $h(x) = h(y)$  and, as  $h$  is a homomorphism,  $\langle h(x), h(z) \rangle \in \theta'$ . By (b) of Theorem 1, there exists  $u \in A$  with  $\langle x, u \rangle \in \theta$  and  $h(u) = h(z)$ , i.e.  $\langle u, z \rangle \in \Phi_h$ . Hence  $\langle x, z \rangle \in \theta \circ \Phi_h$  showing  $\Phi_h \circ \theta \subseteq \theta \circ \Phi_h$ . By Lemma 1,  $\theta$  and  $\Phi_h$  permute.

Conversely, let  $h$  be a surjective strong homomorphism and suppose  $\theta \circ \Phi_h = \Phi_h \circ \theta$ . Since  $h$  is a homomorphism we have  $h([a]_\theta) \subseteq [h(a)]_{\theta'}$ . Let  $x \in [h(a)]_{\theta'}$ . Then  $\langle x, h(a) \rangle \in \theta'$ . Since  $h$  is a strong homomorphism, there exist  $b, c \in A$  such that  $\langle b, c \rangle \in \theta$  and  $h(b) = x, h(c) = h(a)$ . Thus  $\langle c, a \rangle \in \Phi_h$  and we have  $\langle b, a \rangle \in \theta \circ \Phi_h = \Phi_h \circ \theta$ . Hence, there exists  $z \in A$  with  $\langle b, z \rangle \in \Phi_h, \langle z, a \rangle \in \theta$ . Thus  $z \in [a]_\theta$  and  $h(z) = h(b) = x$ , i.e.  $h$  is a class preserving mapping.  $\square$

## References

- [1] Madarász, R., Crvenković, S.: *Relacione algebre. Matematički Institut, Beograd*, 1992.
- [2] Maltsev, A. I.: *Algebraic systems. Nauka, Moskva*, 1970, (in Russian).
- [3] Riguet, J.: *Relations binaires, fermetures, correspondances de Galois*. Bull. Soc. Math. France **76** (1948), 114–155.