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THE GRAPH OF LABELLINGS OF A GIVEN GRAPH

BOHDAN ZELINKA

At the Czechoslovak conference on Graph Theory and Combinatorics in the Račák Valley in May 1986, P. Tomasta [1] has proposed four problems concerning a certain graph $\Pi(H)$ assigned to a finite undirected graph H . Here we shall not solve these problems, but we shall treat simpler questions concerning $\Pi(H)$.

Let a finite undirected graph H be given, let $V(H)$ be its vertex set, let $n = |V(H)|$. A labelling of H is a bijection $\lambda: V(H) \rightarrow \{1, \dots, n\}$. We shall consider ordered pairs (H, λ) , where λ is a labelling of H . To the given graph H there exist $n!$ such pairs.

Let λ_1, λ_2 be two labellings of H . The pairs $(H, \lambda_1), (H, \lambda_2)$ are called isomorphic if they have the property that for any two numbers i, j from the set $\{1, \dots, n\}$ the vertices $\lambda_1^{-1}(i), \lambda_1^{-1}(j)$ are adjacent in H if and only if the vertices $\lambda_2^{-1}(i), \lambda_2^{-1}(j)$ are adjacent in H .

These concepts may be interpreted in the following way. We consider the set $\{1, \dots, n\}$ as a vertex set. To every pair (H, λ) we assign a graph isomorphic to H and having the vertex set $\{1, \dots, n\}$; two vertices i, j from this set will be adjacent if and only if the vertices $\lambda^{-1}(i), \lambda^{-1}(j)$ are adjacent in H . To the isomorphic pairs $(H, \lambda_1), (H, \lambda_2)$ the same graph corresponds; to non-isomorphic ones two distinct (but obviously isomorphic) graphs correspond. This isomorphism is an equivalence relation on the set of all pairs (H, λ) . Any class of this equivalence will be called a labelled graph H . Such a class will be considered as a certain graph isomorphic to H with the vertex set $\{1, \dots, n\}$.

Let $\Lambda(H)$ be the set of all labelled graphs H . Let $\Pi(H)$ be the graph whose vertex set is $\Lambda(H)$ and in which two vertices H_1, H_2 are adjacent if and only if $|E(H_1) - E(H_2)| = |E(H_2) - E(H_1)| = 1$. (The symbol $E(G)$ for any graph G denotes the edge set of G .)

There are four problems by P. Tomasta concerning $\Pi(H)$. The first two concern the neighbourhoods of vertices. The third asks about the characterization of graphs $\Pi(H)$ with a Hamiltonian circuit and the fourth asks an analogous question concerning a Hamiltonian path. These problems seem to be very difficult; note that the number of vertices of $\Pi(H)$ is usually much greater than that of H . We shall treat much simpler problems: the existence of edges in $\Pi(H)$ and its connectedness. Neither of these problems will be solved completely.

A graph will be called discrete if it has no edges.

Theorem 1. *Let H be a regular graph. Then $\Pi(H)$ is discrete.*

Proof. Let r be the regularity degree of H . If $r = 0$, then H is discrete; therefore $(H, \lambda_1), (H, \lambda_2)$ are isomorphic for any λ_1, λ_2 and $\Pi(H)$ consists of one vertex. Now let $r \geq 1$. If H is complete, then $\Pi(H)$ consists again of one vertex and thus it is discrete; otherwise it has at least two vertices. Suppose that H is not complete and two vertices H_1, H_2 of $\Pi(H)$ are adjacent in $\Pi(H)$. Let $E(H_1) - E(H_2) = \{e_1\}, E(H_2) - E(H_1) = \{e_2\}$. Let u, v be the end vertices of e_1 . By deleting e_1 from H_1 we obtain a graph H_0 in which only the vertices u, v have the degree $r - 1$ and all the others have the degree r . The graph H_2 is to be obtained by adding e_2 to H_0 . If e_2 is adjacent to a vertex w distinct from u and v , then w has the degree $r + 1$ in H_2 , which is not possible, because H_2 is regular of degree r . Thus e_2 can join only u and v and thus $e_2 = e_1$, which is a contradiction. \square

Theorem 2. *Let H be a graph, let $\delta(H)$ (or $\Delta(H)$) be its minimum (or maximum respectively) degree. Let there exist an integer r such that $\delta(H) < r < \Delta(H)$ and no vertex of H has the degree r . Then $\Pi(H)$ is disconnected.*

Proof. Let Λ_1 (or Λ_2) be the subset of $\Lambda(H)$ consisting of labelled graphs with the property that the degree of the vertex 1 is less (or greater respectively) than r . Evidently $\Lambda_1 \neq \emptyset, \Lambda_2 \neq \emptyset, \Lambda_1 \cup \Lambda_2 = \Lambda(H), \Lambda_1 \cap \Lambda_2 = \emptyset$. Let $H_1 \in \Lambda_1, H_2 \in \Lambda_2$. The difference between the degrees of the vertex 1 in H_1 and in H_2 is at least 2 and thus also $|E(H_1) - E(H_2)| \geq 2$ and H_1, H_2 are not adjacent in $\Pi(H)$. As H_1, H_2 were chosen arbitrarily, no vertex of Λ_1 is adjacent to a vertex of Λ_2 and $\Pi(H)$ is disconnected. \square

Theorem 3. *Let H be a disconnected graph in which all connected components have the same number of edges and are not all isomorphic. Then $\Pi(H)$ is disconnected.*

Proof. Let D_1, D_2 be two non-isomorphic connected components of H . Let Λ_1 (or Λ_2) be the subset of $\Lambda(H)$ consisting of all labelled graphs such that the vertex 1 is contained in a connected component of H_1 isomorphic to D_1 (or D_2 respectively). Let $H_1 \in \Lambda_1, H_2 \in \Lambda_2$. Let $E(H_1) - E(H_2) = \{e_1\}, E(H_2) - E(H_1) = \{e_2\}$. Then e_2 joins two vertices of the connected component of H_1 which contains e_1 ; otherwise there would exist in H_2 connected components with different numbers of edges. If e_1 does not belong to the connected component of H_1 containing 1, then 1 in H_2 belongs again to a connected component isomorphic to D_1 , which is a contradiction. If e_1 belongs to the connected component of H_1 containing 1 and this component after deleting e_1 and adding e_2 becomes isomorphic to D_2 , then the number of connected components of H_2 isomorphic to D_1 (or to D_2) is less (or greater respectively) than that of H_1 . Thus H_1 is not isomorphic to H_2 , which is again a contradiction. \square

Theorem 4. *Let H be a disconnected graph consisting of 2-edge-connected*

components. Let there exist an integer r such that no connected component of H has r edges, but there exist those with less and those with more than r edges. Then $\Pi(H)$ is disconnected.

Proof. Let Λ_1 (or Λ_2) be the subset of $\Pi(H)$ consisting of labelled graphs H such that the vertex 1 is in a connected component with less (or more respectively) edges than r . We have $\Lambda_1 \neq \emptyset$, $\Lambda_2 \neq \emptyset$, $\Lambda_1 \cup \Lambda_2 = \Lambda(H)$, $\Lambda_1 \cap \Lambda_2 = \emptyset$. Let $H_1 \in \Lambda_1$, $H_2 \in \Lambda_2$ and suppose that H_1, H_2 are adjacent in $\Pi(H)$. Let $E(H_1) - E(H_2) = \{e_1\}$, $E(H_2) - E(H_1) = \{e_2\}$. As each connected component of H is 2-edge-connected, in the graph obtained from H_1 by deleting e_1 the vertex sets of the connected components are the same as in H_1 . After adding e_2 to it the number of edges of the connected component containing 1 can increase at most by one, hence it cannot become greater than r , which is a contradiction. Hence no vertex of Λ_1 is adjacent to a vertex of Λ_2 and $\Pi(H)$ is disconnected. \square

Theorem 5. Let H be a 2-edge-connected bipartite graph with the bipartition classes A, B , let $|A| \neq |B|$. Then $\Pi(H)$ is disconnected.

Proof. Let Λ_1 (or Λ_2) be the subset of $\Lambda(H)$ consisting of labelled graphs H such that the vertex 1 is in the set corresponding to A (or to B respectively). As H is 2-edge-connected, its bipartition classes remain the same after deleting an arbitrary edge. If $H_1 \in \Lambda_1$, then by deleting an edge e_1 and adding an edge e_2 we obtain again a graph from Λ_1 and thus analogously to the proof of Theorem 4 we may prove that $\Pi(H)$ is disconnected. \square

Theorem 6. Let H be a graph, let \bar{H} be its complement. Then $\Pi(H) \cong \Pi(\bar{H})$.

Proof. Let H_1, H_2 be two graphs from $\Lambda(H)$. The graph H_2 is obtained from H_1 by deleting an edge e_1 and adding an edge e_2 if and only if \bar{H}_2 is obtained from \bar{H}_1 by deleting e_2 and adding e_1 . This implies the assertion. \square

We have shown there exist wide classes of graphs H such that $\Pi(H)$ is discrete and wide classes of graphs H such that $\Pi(H)$ is disconnected. We may ask about two problems.

Problem 1. Characterize the graphs H for which $\Pi(H)$ is discrete, i. e. graphs H with the property that by deleting an arbitrary edge e_1 and adding a new edge $e_2 \neq e_1$ always a graph non-isomorphic to H is obtained.

Problem 2. Characterize the graphs H for which $\Pi(H)$ is disconnected.

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ГРАФ ПОМЕЧЕНИЙ ЗАДАННОГО ГРАФА

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Резюме

Пусть H есть конечный неориентированный граф с n вершинами. П. Томаста ввел граф $P(H)$, вершинами которого являются все графы на множестве вершин $\{1, \dots, n\}$, изоморфные графу H , и в котором две вершины H_1, H_2 смежны тогда и только тогда, когда $|E(H_1) - E(H_2)| = |E(H_2) - E(H_1)| = 1$. Показаны некоторые классы графов H , для которых $P(H)$ является несвязным или дискретным.