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## ON A SUM OF OBSERVABLES IN A LOGIC

ANATOLIJ DVUREČENSKIJ

A sum of two observables of a logic defined in a way differing from that of the mean values is studied and some properties are proved.

### Introduction

In the classical probability theory the sum of observables is, doubtless, of great importance. Therefore there are made different attempts to introduce the sum into the theory of logic [2—6], as well as into the quantum measuring theory of noncompatible observables. We shall study the properties of the sum defined by (2.1).

### 1. Logic and observables

Let  $L$  be a  $\sigma$ -lattice with the first and the last elements 0 and 1, respectively, and an orthocomplementation  $\perp : a \mapsto a^\perp$  which satisfies (i)  $(a^\perp)^\perp = a$  for all  $a \in L$ ; (ii) if  $a < b$ , then  $b^\perp < a^\perp$  for  $a, b \in L$ ; (iii)  $a \vee a^\perp = 1$  for all  $a \in L$ . We further assume that if  $a < b$ , then  $b = a \vee (b \wedge a^\perp)$ . A poset  $L$  satisfying the above axioms will be called a logic.

We say that  $a, b$  are (i) orthogonal and we write  $a \perp b$  if  $a < b^\perp$ ; (ii) compatible and we write  $a \leftrightarrow b$  if there are three mutually orthogonal elements  $a_1, b_1, c \in L$  such that  $a = a_1 \vee c, b = b_1 \vee c$ .

An observable is a map  $x$  from  $B(R_1)$  into  $L$  such that (i)  $x(R_1) = 1, x(\emptyset) = 0$ ; (ii)  $x(E) \perp x(F)$  if  $E \cap F = \emptyset, E, F \in B(R_1)$ ; (iii)  $x\left(\bigcup_i E_i\right) = \bigvee_i x(E_i)$  if  $E_i \cap E_j = \emptyset, i \neq j, \{E_i\} \subset B(R_1)$ . If  $f$  is a Borel function on  $R_1$  and  $x$  an observable, then  $f \circ x : E \mapsto x(f^{-1}(E)), E \in B(R_1)$ , is an observable. For an observable  $x$  we denote  $\sigma(x) = \cap \{C \in B(R_1) : x(C) = 1\}$  and we define  $\|x\| = \sup \{|t| : t \in \sigma(x)\}$ . We say that  $x$  is (i) bounded if  $\|x\| < \infty$ ; (ii) bounded above (below) if there is a number  $c \in R_1$  such that  $\sigma(x) \subset (-\infty, c)$  ( $\sigma(x) \subset (c, \infty)$ ). Two observables  $x$  and  $y$  are compatible and we write  $x \leftrightarrow y$  if  $x(E) \leftrightarrow y(F)$  for every  $E, F \in B(R_1)$ .

The conventional measurable space  $(\Omega, \mathcal{S})$  is a logic of compatible observables if we identify  $x(E) = f^{-1}(E)$ ,  $E \in B(R_1)$ , where  $f$  is a  $\mathcal{S}$  — measurable function. The logic  $L(H)$ , that is, the complete lattice of all closed subspaces of a Hilbert space  $H$ , is a very important example of a logic which has noncompatible observables and which is a model for quantum mechanics. In this logic the selfadjoint operators correspond to the observables [8].

Since the notion of observable is an analogy of a measurable function we will now investigate some properties of observables.

**Theorem 1.1.** *Let  $x$  be an observable of a logic  $L$  and  $B_x(t) = x((-\infty, t))$ ,  $t \in R_1$ , then the system  $\{B_x(t) : t \in R_1\}$  has the following properties:*

- (i)  $B_x(s) < B_x(t)$  if  $s < t$ ;
- (ii)  $\bigvee_t B_x(t) = 1$ ,  $\bigwedge_t B_x(t) = 0$ ;
- (iii)  $\bigvee_{t < s} B_x(t) = B_x(s)$ .

*Conversely, if a system  $\{B(t) : t \in R_1\}$  of the elements of a logic  $L$  fulfils (1.1), then there is a unique observable  $x$  such that  $B_x(t) = B(t)$  for every  $t \in R_1$ .*

*Proof.* Let  $x$  be an observable; then (i) is trivial. (ii): let  $B_x(t) < a$  for every  $t \in R_1$ ; then for every integer  $n$  we have  $B_x(n) < a$ . Hence  $a > \bigvee_n B_x(n) = \bigvee_n x((-\infty, n)) = 1$ . Similarly,  $\bigwedge_t B_x(t) = 0$ . (iii): let  $a > B_x(t)$ ,  $t < s$ . If we choose  $t_n \uparrow s$ , then  $a > \bigvee_n B_x(t_n) = B_x(s)$ .

Let now on the logic  $L$  a system  $\{B(t) : t \in R_1\}$  satisfying (i)—(iii) be given. In the first place we show that there is a Boolean sub- $\sigma$ -algebra of  $L$  generating by  $\{B(t) : t \in R_1\}$ .

Let  $r_1, r_2, \dots$  be any distinct enumeration of the rational numbers in  $R_1$ . For every  $n$  let  $\mathcal{A}_n$  be a Boolean subalgebra of  $L$  generated by  $\{B(r_1), \dots, B(r_n)\}$ . This subalgebra surely exists, because if  $(i_1, \dots, i_n)$  is such an enumeration of  $(1, \dots, n)$  that  $r_{i_1} < \dots < r_{i_n}$ , then the set of all finite lattice sums of orthogonal elements  $\{B(r_{i_1}), B(r_{i_2}) \wedge B(r_{i_1})^\perp, \dots, B(r_{i_n}) \wedge B(r_{i_{n-1}})^\perp, B(r_{i_n})^\perp\}$  is a Boolean subalgebra containing all  $B(r_1), \dots, B(r_n)$  and therefore it is  $\mathcal{A}_n$ . Let us put  $\mathcal{A}_0 = \bigcup_n \mathcal{A}_n$ ; then  $\mathcal{A}_0$  is a Boolean subalgebra of  $L$ , too.

By the Zorn lemma it is easy to see that there is a maximal Boolean subalgebra  $\mathcal{M}$  of  $L$  containing  $\mathcal{A}_0$ . The  $\mathcal{M}$  must be a Boolean sub- $\sigma$ -algebra.

Let now  $B(t)$  be an arbitrary element of  $\{B(t) : t \in R_1\}$ . Since there is  $r_n \uparrow t$ , we have  $B(t) = \bigvee_i B(r_{n_i}) \in \mathcal{M}$ . We have shown that there is a Boolean sub- $\sigma$ -algebra of

$L$  generated by  $\{B(t) : t \in R_1\}$  and let it be denoted by  $\mathcal{A}$ .

By the Loomis theorem there is a measurable space  $(\Omega, \mathcal{S})$  and a homomorphism  $h$  from  $\mathcal{S}$  onto  $\mathcal{A}$ . We claim to construct, by induction, the set s  $A_1, A_2, \dots$  from  $\mathcal{A}$  such that

- (a)  $h(A_i) = B(r_i)$ ;
- (b)  $A_i \subset A_j$  if  $r_i < r_j$ ;
- (c)  $\bigcap_{i=1}^{\infty} A_i = \emptyset$ .

We note that if  $A \subset B$ ,  $A, B \in \mathcal{S}$  and if there is  $c \in \mathcal{A}$  such that  $h(A) < c < h(B)$ , then there is  $C \in \mathcal{S}$  such that  $A \subset C \subset B$ ,  $h(C) = c$ . Indeed, since  $h$  maps  $\mathcal{S}$  onto  $\mathcal{A}$ , there is  $C_1 \in \mathcal{S}$  such that  $h(C_1) = c$ . If we define  $C = (C_1 \cap B) \cup A$ , then  $C$  has a given property.

Let  $A_1$  be any set in  $\mathcal{S}$  such that  $h(A_1) = B(r_1)$ . Suppose  $A_1, \dots, A_n \in \mathcal{S}$  have been constructed so that (a) and (b) hold. We shall construct  $A_{n+1}$  as follows. Let  $(i_1, \dots, i_n)$  be the permutation of  $(1, \dots, n)$  such that  $r_{i_1} < \dots < r_{i_n}$ . Then only one condition holds (\*): (i)  $r_{n+1} < r_{i_1}$ ; (ii)  $r_{n+1} > r_{i_n}$ ; (iii) there is a unique  $k = 1, \dots, n$  such that  $r_{i_k} < r_{n+1} < r_{i_{k+1}}$ ; and by the above observation we can select  $A_{n+1}$  such that  $h(A_{n+1}) = B(r_{n+1})$  and (i)  $A_{n+1} \subset A_{i_1}$ ; (ii)  $A_{n+1} \supset A_{i_n}$ ; (iii)  $A_{i_k} \subset A_{n+1} \subset A_{i_{k+1}}$ ; according to (\*). Then the system  $\{A_1, \dots, A_{n+1}\}$  fulfils (a) and (b). Thus, by induction, there follows that there is a sequence  $\{A_j\}$  of sets in  $\mathcal{S}$  with the properties (a) and (b). As

$$h\left(\bigcap_{j=1}^{\infty} A_j\right) = \bigwedge_{j=1}^{\infty} h(A_j) = \bigwedge_{j=1}^{\infty} B(r_j) = 0,$$

we may, replacing  $A_j$  by  $A_j - \bigcap_i A_i$  if necessary, assume that  $\bigcap_i A_i = \emptyset$ .

We define an  $\mathcal{S}$ -measurable function  $f$  as follows:

$$f(\omega) = \begin{cases} 0 & \text{if } \omega \notin \bigcup_{j=1}^{\infty} A_j \\ \inf \{r_j : \omega \in A_j\} & \text{if } \omega \in \bigcup_{j=1}^{\infty} A_j. \end{cases}$$

A function  $f$  is everywhere well defined and it is finite. Moreover

$$f^{-1}((-\infty, r_k)) = \begin{cases} \bigcup_{r_j < r_k} A_j & \text{if } r_k \leq 0 \\ \bigcup_{r_j < r_k} A_j \cup \left(\Omega - \bigcup_i A_i\right) & \text{if } r_k > 0, \end{cases}$$

hence  $f$  is  $\mathcal{S}$ -measurable and  $h(f^{-1}((-\infty, r_k))) = B(r_k)$ . If we define an observable  $x$  by  $x(E) = h(f^{-1}(E))$ ,  $E \in B(R_1)$ , then  $x((-\infty, t)) = B(t)$  for every  $t \in R_1$ . Since  $x_1((-\infty, t)) = x_2((-\infty, t))$  for every  $t \in R_1$  implies  $x_1 = x_2$ , the uniqueness of  $x$  is shown and the proof is finished. Q.E.D.

Remark 1.2. (i) Theorem 1.1 holds if we consider a system  $\{B(t): t \in S\}$  satisfying (1.1), where  $S$  is a countable dense set in  $R_1$ .

(ii) If  $L$  is a non-lattice logic [7], then the assertions of Theorem 1.1 and the first part of Remark 1.2 remain valid, too.

**Theorem 1.3.** *For two observables  $x$  and  $y$  the following conditions are equivalent:*

- (i)  $x \leftrightarrow y$ ;
- (ii)  $B_x(t) \leftrightarrow B_y(s)$  for every  $s, t \in R_1$ ;
- (iii)  $B_x(t) \leftrightarrow B_y(s)$  for every  $s, t \in S$ ,  $S$  is a countable dense set in  $R_1$ .

*Proof.* The implication (i)  $\Rightarrow$  (ii)  $\Rightarrow$  (iii) is trivial. Let now (iii) hold. Let us denote for any  $t \in S$

$$\mathcal{C}_t = \{E \in B(R_1): x(E) \leftrightarrow B_y(t)\}.$$

If we take into account the assertion of Lemma 6.10 [8]: if  $b \leftrightarrow a_n$ ,  $n = 1, 2, \dots$ , then  $b \leftrightarrow a_n^+$ ,  $n = 1, 2, \dots$ ,  $b \leftrightarrow \bigvee_n a_n$ ,  $b \leftrightarrow \bigwedge_n a_n$ ; then  $\mathcal{C}_t = B(R_1)$ . Indeed,  $\mathcal{C}_t$  contains the intervals  $(-\infty, s)$  for every  $s \in S$ . Let  $s \in R_1$ ; then there is  $s_n \uparrow s$ ,  $s_n \in S$ . Hence  $(-\infty, s) \in \mathcal{C}_t$  for every  $s \in R_1$  and, consequently,  $\mathcal{C}_t = B(R_1)$ ,  $t \in S$ . Similarly,  $\mathcal{C}_t = B(R_1)$  for any  $t \in R_1$ . Analogically,  $\mathcal{C} = \{F \in B(R_1): x(E) \leftrightarrow y(F) \text{ for every } E \in B(R_1)\} = B(R_1)$ . Therefore  $x \leftrightarrow y$ . Q.E.D.

## 2. The sum of two observables

If  $x$  and  $y$  are compatible observables, then, by [8, Theorem 6.9], there are an observable  $u$  and two Borel functions  $f, g$  such that  $x = f \circ u$ ,  $y = g \circ u$ . Due to Theorem 6.17 [8] we may define the sum of  $x$  and  $y$  by  $x + y = (f + g) \circ u$  independently of the used  $f, g, u$ . Theorem 1.1 enables us to define the sum for noncompatible observables without using the mean values.

For two observables  $x, y$  we define the following system of the elements of a logic  $L$ :

$$B_{x \oplus y}(t) = \bigvee_{r \in Q} (B_x(r) \wedge B_y(t - r)), \quad t \in R_1, \quad (2.1)$$

where  $Q$  is the set of the rational numbers in  $R_1$ .

**Lemma 2.1.** *If  $x \leftrightarrow y$ , then a system  $\{b_{x \oplus y}(t): t \in R_1\}$  fulfils (1.1) of Theorem 1.1, and then an observable  $x \oplus y$  coincides with the sum of compatible observables.*

*Proof.* There holds

$$B_{x \oplus y}(t) = \bigvee_{r \in Q} (x((-\infty, r)) \wedge y((-\infty, t - r))) =$$

$$\bigvee_{r \in Q} [u(f^{-1}((-\infty, r))) \wedge u(g^{-1}((-\infty, t-r)))] = u((f+g)^{-1}((-\infty, t))) = B_{(f+g) \circ u}(t).$$

Hence  $B_x \oplus_y(t)$  fulfils (1.1) and  $x \oplus y = (f+g) \circ u = x + y$ . Q.E.D.

A logic  $L$  is  $\sigma$ -continuous if for  $a_1 < a_2 < \dots$  and any  $a$

$$a \wedge \left( \bigvee_n a_n \right) = \bigvee_n (a \wedge a_n)$$

holds. A logic  $L$  is said to satisfy the finite chain condition (f.c.c.) if  $\{a_n\} \subset L$  with  $a_1 < a_2 < \dots$  implies that there is  $N$  such that  $a_n = a_N$  for  $n > N$ . It is easy to see that if  $L$  satisfies f.c.c., then it is  $\sigma$ -continuous.

**Lemma 2.2.** *Let  $L$  be a  $\sigma$ -continuous logic and  $S$  a countable dense set in  $R_1$ . Let us denote for the observables  $x, y$   $B_{x \oplus_y}^s(t) = \bigvee_{s \in S} (B_x(s) \wedge B_y(t-s))$ ; then  $B_{x \oplus_y}^s(t) = B_{x \oplus_y}(t)$  for every  $t \in R_1$ .*

*Proof.* We may show that if  $t_n \uparrow t$ , then  $B_{x \oplus_y}^s(t) = \bigvee_n B_{x \oplus_y}^s(t_n)$ . Indeed,

$$\begin{aligned} \bigvee_n B_{x \oplus_y}^s(t_n) &= \bigvee_n \bigvee_{s \in S} (B_x(s) \wedge B_y(t_n - s)) = \\ &= \bigvee_{s \in S} (B_x(s) \wedge \bigvee_n B_y(t_n - s)) = \bigvee_{s \in S} (B_x(s) \wedge B_y(t - s)). \end{aligned}$$

Let now  $n$  be any integer; then for each  $s$  there is  $r = r(s) \in Q$  such that we have  $s < r < s + n^{-1}$ . Therefore  $B_x(s) \wedge B_y(t - n^{-1} - s) < B_x(r) \wedge B_y(t - r)$  and

$$\begin{aligned} B_{x \oplus_y}^s(t - n^{-1}) &< B_{x \oplus_y}(t) \\ B_{x \oplus_y}^s(t) &= \bigvee_n B_{x \oplus_y}^s(t - n^{-1}) < B_{x \oplus_y}(t). \end{aligned}$$

Similarly we show that  $B_{x \oplus_y}(t) < B_{x \oplus_y}^s(t)$ . Q.E.D.

**Theorem 2.3.** *Let  $L$  be a  $\sigma$ -continuous logic and  $x, y$  be observables. Then for  $\{B_{x \oplus_y}(t) : t \in R_1\}$  we have*

- (i)  $B_{x \oplus_y}(s) < B_{x \oplus_y}(t)$ ,  $s < t$  (on any logic, too);
- (ii)  $\bigvee_t B_{x \oplus_y}(t) = 1$  (if  $x, y$  are bounded above, then (ii) holds on any logic);
- (iii)  $\bigwedge_t B_{x \oplus_y}(t) = 0$  (on any logic) if  $x, y$  are bounded below;
- (iv)  $\bigvee_{t \leq s} B_{x \oplus_y}(t) = B_{x \oplus_y}(s)$ ;
- (v)  $B_{x \oplus_y}(t) = B_{y \oplus_x}(t)$  for every  $t \in R_1$ .

*Proof.* Because of (i) of (1.1) the (i) is evident. (ii)

$$\bigvee_t B_{x \oplus_y}(t) = \bigvee_t \bigvee_{r \in Q} (B_x(r) \wedge B_y(t-r)) =$$

$$\begin{aligned}
&= \bigvee_{r \in Q} \bigvee_t (B_x(r) \wedge B_y(t-r)) > \bigvee_{r \in Q} \bigvee_{n=1}^{\infty} (B_x(r) \wedge B_y(n-r)) = \\
&= \bigvee_{r \in Q} (B_x(r) \wedge \bigvee_{n=1}^{\infty} B_y(n-r)) = \bigvee_{r \in Q} (B_x(r) \wedge 1) = 1,
\end{aligned}$$

by the  $\sigma$ -continuity of  $L$ . Similarly for (iv).

If  $x, y$  are bounded above, then there is  $c \in R_1$  such that  $\sigma(x), \sigma(y) \subset (-\infty, c)$ . Then for any  $\varepsilon > 0$   $B_x(c + \varepsilon) = 1 = B_y(c + \varepsilon)$ . Hence  $B_{x \oplus y}(2c + 2\varepsilon) = 1$ .

(iii) There is  $c \in R_1$  such that  $\sigma(x), \sigma(y) \subset \langle c, \infty$ . Then  $B_{x \oplus y}(2c) = 0$ .

(v) Let  $t \in R_1$ ; then the set  $S_t = \{s = t - r : r \in Q\}$  is countable dense in  $R_1$  and, by Lemma 2.2, we have

$$B_{x \oplus y}(t) = \bigvee_{r \in Q} (B_x(r) \wedge B_y(t-r)) = \bigvee_{s \in S_t} (B_y(s) \wedge B_x(t-s)) = B_{y \oplus x}(t).$$

Q.E.D.

**Lemma 2.4.** *Let  $x, y$  be two observables bounded below on a  $\sigma$ -logic  $L$ . Then*

$$\|x \oplus y\| \leq \|x\| + \|y\|. \quad (2.2)$$

*Proof.* If  $x$  or  $y$  is unbounded, then (2.2) holds. Therefore let  $x, y$  be bounded observables. Let us denote

$$\begin{aligned}
a_1 &= \inf \sigma(x), & b_1 &= \sup \sigma(x) \\
a_2 &= \inf \sigma(y), & b_2 &= \sup \sigma(y).
\end{aligned}$$

Then  $B_{x \oplus y}(a_1 + a_2) = 0$  and  $B_{x \oplus y}(b_1 + b_2 + \varepsilon) = 1$  for every  $\varepsilon > 0$ . We prove only  $B_{x \oplus y}(b_1 + b_2 + \varepsilon) = 1$ . If we choose a rational number  $r$  such that  $b_1 + \varepsilon/4 < r < b_1 + \varepsilon/2$ , then  $-r > -b_1 - \varepsilon/2$  and

$$\begin{aligned}
B_x(r) &> B_x(b_1 + \varepsilon/4) = 1, \\
B_y(b_1 + b_2 + \varepsilon - r) &> B_y(b_1 + b_2 + \varepsilon - b_1 - \varepsilon/2) = B_y(b_2 + \varepsilon/2) = 1.
\end{aligned}$$

We have proved that  $\sigma(x \oplus y) \subset \langle a_1 + a_2, b_1 + b_2 \rangle$ . If  $a = \inf \sigma(x \oplus y)$ ,  $b = \sup \sigma(x \oplus y)$ , then  $a_1 + a_2 \leq a \leq b \leq b_1 + b_2$ . We calculate  $\|x \oplus y\| = \max \{|a|, |b|\} \leq \max \{|a_1 + a_2|, |b_1 + b_2|\} \leq \max \{|a_1|, |b_1|\} + \max \{|a_2|, |b_2|\} = \|x\| + \|y\|$ .

Q.E.D.

We denote by  $o$  such an observable that  $o(\{0\}) = 1$ . For  $\alpha \in R_1$  and  $x$  we denote by  $\alpha x$  such an observable that  $(\alpha x)(E) = x(\{t : \alpha(t) \in E\})$ , where  $\alpha(t) \equiv at$ ,  $t \in R_1$  and finally, for  $x, y$  we denote  $x \ominus y = x \oplus (-y)$ .

**Theorem 2.5.** *Let  $O_B(L)$  be the set of all bounded observables on a  $\sigma$ -continuous logic  $L$ . Then  $O_B(L)$  is a normed space with respect to the norm  $\|x\| = \sup \{|t| : t \in \sigma(x)\}$  and the following properties hold*

- (i)  $\|x\| \geq 0$ ,  $x \in O_B(L)$ ,  $\|x\| = 0$  iff  $x = o$ ;
- (ii)  $\|\alpha x\| = |\alpha| \|x\|$ ,  $\alpha \in R_1$ ,  $x \in O_B(L)$ ;
- (iii)  $\|x \oplus y\| \leq \|x\| + \|y\|$ ,  $x, y \in O_B(L)$ ;

- (iv)  $x \oplus y = y \oplus x$ ,  $x, y \in O_B(L)$ ;
- (v)  $x \oplus o = x$ ;  $x \in O_B(L)$ ;
- (vi)  $x \ominus x = o$ ,  $x \in O_B(L)$ ;
- (vii)  $(\alpha + \beta)x = \alpha x \oplus \beta x$ ,  $\alpha, \beta \in R_1$ ,  $x \in O_B(L)$ ;
- (viii)  $\alpha(x \oplus y) = \alpha x \oplus \alpha y$ ,  $\alpha \geq 0$ ,  $x, y \in O_B(L)$ .

Proof. The properties (i)—(ii) follow from [3, Theorem 4.2], (iii) follows from Lemma 2.4, (iv) from Lemma 2.3; (v)—(vii) are the corollaries of the calculus for compatible observables; (viii) follows from the definition of the sum and from Lemma 2.2. Q.E.D.

For a given element  $a \in L$  we define a question observable  $q_a$  by  $q_a(\{0\}) = a^\perp$ ,  $q_a(\{1\}) = a$ , and an observable  $x$  is a question observable iff  $\sigma(x) \subset \{0, 1\}$  [4].

Remark 2.6. The sum defined by (2.1) is not associative in general.

Indeed, let  $a, b, c \in L$ ; then

$$B_{(q_a \oplus q_b) \oplus q_c}(t) = \begin{cases} 0 & t \leq 0 \\ (a \vee b \vee c)^\perp & 0 < t \leq 1 \\ (a \vee b)^\perp \vee ((a \wedge b)^\perp \wedge c^\perp) & 1 < t \leq 2 \\ (a \wedge b \wedge c)^\perp & 2 < t \leq 3 \\ 1 & 3 < t \end{cases}$$

$$B_{q_a \oplus (q_b \oplus q_c)}(t) = \begin{cases} 0 & t \leq 0 \\ (a \vee b \vee c)^\perp & 0 < t \leq 1 \\ (b \vee c)^\perp \vee ((b \wedge c)^\perp \wedge a^\perp) & 1 < t \leq 2 \\ (a \wedge b \wedge c)^\perp & 2 < t \leq 3 \\ 1 & 3 < t. \end{cases}$$

If now  $L = L(R_2)$  and  $a, b, c$  are three mutually distinct noncompatible subspaces, then

$$B_{(q_a \oplus q_b) \oplus q_c}(2) = c^\perp, \quad B_{q_a \oplus (q_b \oplus q_c)}(2) = a^\perp. \quad \text{Q.E.D.}$$

**Lemma 2.7.** *If for  $x_1, \dots, x_n$  we define, by the recurrence formula,  $x_1 \oplus \dots \oplus x_n = (x_1 \oplus \dots \oplus x_{n-1}) \oplus x_n$ ,  $n = 1, 2, \dots$ , then*

- (i)  $q_a \oplus q_b(\{i\}) = \begin{cases} (a \vee b)^\perp & \text{if } i = 0 \\ (a \vee b) \wedge (a \wedge b)^\perp & \text{if } i = 1 \\ a \wedge b & \text{if } i = 2; \end{cases}$
- (ii)  $q_{a_1} \oplus \dots \oplus q_{a_n}(\{0\}) = (a_1 \vee \dots \vee a_n)^\perp$ ;
- (iii)  $q_{a_1} \oplus \dots \oplus q_{a_n}(\{n\}) = a_1 \wedge \dots \wedge a_n$ ;
- (iv)  $\sigma(q_{a_1} \oplus \dots \oplus q_{a_n}) \subset \{0, 1, n\}$ .

Proof. The property (i) follows from the definition of the sum, and (ii)—(iv) may be proved by induction. Q.E.D.



### 3. Comparison with the sum defined by mean values

Gudder in [4] studied the sum of bounded observables defined by mean values. Let  $m$  be a state, that is, a map from  $L$  into  $\langle 0, 1 \rangle$  such that (i)  $m(1) = 1$ ; (ii)  $m\left(\bigvee_i a_i\right) = \sum_i m(a_i)$ , if  $a_i \perp a_j$ ,  $i \neq j$ , then the mean value of an observable in  $m$  is  $m(x) = \int t \, dm_x(t)$  if the integral on the right-hand side exists and is finite, where  $m_x$  is a measure on  $B(R_1)$ :  $m_x(E) = m(x(E))$ ,  $E \in B(R_1)$ . If there is a quite full system  $M$  of states [4] such that for any two bounded observables  $x, y$  there is a unique observable  $z$  such that

$$m(z) = m(x) + m(y), \quad \text{for every } m \in M, \quad (3.1)$$

then  $z$  is called the sum of  $x, y$  and it is written  $z = x + y$ .

It is easy to see that this sum is associative and it coincides with the sum of compatible observables.

Example 3.1. Let  $L = L(R_2)$  and let  $x, y, z$  correspond to

$$\mathbf{M}_1 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \quad \mathbf{M}_2 = \begin{pmatrix} 1/2 & 1/2 \\ 1/2 & 1/2 \end{pmatrix}, \quad \mathbf{M} = \mathbf{M}_1 + \mathbf{M}_2 = \begin{pmatrix} 3/2 & 1/2 \\ 1/2 & 1/2 \end{pmatrix}.$$

Then the sum of  $x, y$  defined (i) by (3.1) is  $z$ ; (ii) by (2.1) is  $q_1$ . The logic  $L(R_2)$  is isomorphic to a logic  $L$  of subsets of the set  $\Omega = \langle 0, \pi/2 \rangle$ , that is, with the logic  $L = \{\emptyset, \Omega, \{\pi/2, \varphi\}, \{\pi/2, \varphi\}^c, 0 \leq \varphi < \pi/2\}$ . Let  $f, g, h$  correspond to  $x, y, z$  in this isomorphism, where

$$f(\omega) = \begin{cases} 0 & \text{if } \omega \in \{\pi/2, 0\} \\ 0 & \text{if } \omega \notin \{\pi/2, 0\}; \end{cases} \quad g(\omega) = \begin{cases} 1 & \text{if } \omega \in \{\pi/2, \pi/4\} \\ 0 & \text{if } \omega \notin \{\pi/2, \pi/4\}; \end{cases}$$

$$h(\omega) = \begin{cases} (2 - \sqrt{2})/2 & \text{if } \omega \notin \{\pi/2, \arctg(1 + \sqrt{2})\} \\ (2 + \sqrt{2})/2 & \text{if } \omega \in \{\pi/2, \arctg(1 + \sqrt{2})\} \end{cases}$$

Now, if we define the sum of measurable functions  $f, g$ :

- (i) by points, that is,  $(f + g)(\omega) = f(\omega) + g(\omega) \Rightarrow f + g$  is no observable;
- (ii) by (3.1), then  $f + g = h$ ;
- (iii) by (2.1), then  $f + g = 1$ .

This example refers to the splitting of the notion of the sum in a transition from a measurable space into logics. Moreover, in [1] it is shown that although  $(f + g)(\omega) = f(\omega) + g(\omega)$  is a measurable function, the additivity of the mean value does not hold in general. ( $f, g$  in [1] are unbounded observables.)

**Lemma 3.2.** *The following propositions are equivalent*

- (i)  $q_a \oplus q_b$  is a question observable;
- (ii)  $q_a \oplus q_b = q_{a \vee b}$ ;
- (iii)  $a \wedge b = 0$ .

S. P. Gudder in [4] showed that  $a \perp b$  iff  $q_a + q_b = q_{a \vee b}$ . This property does not hold for the sum defined by (2.1).

**Corollary 3.2.1.** *If there holds  $a \perp b$  iff  $q_a \oplus q_b = q_{a \vee b}$ , then  $L$  is a Boolean  $\sigma$ -algebra.*

Proof. If  $q_a \oplus q_b = q_{a \vee b}$ , then by (ii) of Lemma 3.2 there follows that  $a \perp b$  iff  $a \wedge b = 0$ . By Zierler [9, Lemma 1.5] there implies that  $L$  is a Boolean  $\sigma$ -algebra.

Q.E.D.

**Lemma 3.3.** *There holds*

$$B_{q_a \ominus q_b}(t) = \begin{cases} 0 & t \leq -1 \\ a^\perp \wedge b & -1 < t \leq 0 \\ a^\perp \cup b & 0 < t \leq 1 \\ 1 & 1 < t \end{cases}$$

Moreover, the following propositions are equivalent

- (i)  $q_a \ominus q_b$  is a question observable;
- (ii)  $q_a \ominus q_b = q_{a \wedge b^\perp}$
- (iii)  $a^\perp \wedge b = 0$ .

We see that the sum of two observables  $x \oplus y$  has not the same properties as the sum defined by (3.1) and the investigation of the sum defined by (2.1) may be made mainly for compatible observables.

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## О СУММЕ НАБЛЮДАЕМЫХ В ЛОГИКЕ

Анатолий Двуреченский

### Резюме

Сумма двух наблюдаемых в логике определяется отличным способом от определения суммы посредством средних значений. Некоторые свойства этой суммы доказаны