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A Note on Tolerance Lattices of Algebras with Restricted Similarity Type

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Abstract

It is shown that for every finite algebra \mathcal{A} of a finite similarity type there exists a finite algebra \mathcal{C} of type $(2,1,1)$ such that the lattices of tolerances (congruences) on \mathcal{A} and \mathcal{C} are isomorphic.

Key words: Finite algebra of a finite similarity type, Lattice of congruences, Lattice of tolerances.

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In [1] I. Chajda proved that for every algebra \mathcal{A} of a finite similarity type there exists an algebra \mathcal{B} of type $(2,1,1)$ such that $\text{Tol } \mathcal{A} \cong \text{Tol } \mathcal{B}$.

Recall that $\mathcal{A} = (A, F)$ is of a finite similarity type if F is a finite set. By $\text{Tol } \mathcal{A}$ we shall mean the lattice of all tolerances on an algebra \mathcal{A} with respect to set inclusion and a tolerance on \mathcal{A} is a binary reflexive and symmetric relation on \mathcal{A} which is a subalgebra of the direct product $\mathcal{A} \times \mathcal{A}$. See [2].

Unfortunately, if an algebra \mathcal{A} is finite and nontrivial, then Chajda's algebra \mathcal{B} of type $(2,1,1)$ satisfying $\text{Tol } \mathcal{A} \cong \text{Tol } \mathcal{B}$ is not finite. In this note we shall show the following:

Theorem *For every finite algebra \mathcal{A} of a finite similarity type there exists a finite algebra \mathcal{C} of type $(2,1,1)$ such that $\text{Tol } \mathcal{A} \cong \text{Tol } \mathcal{C}$.*

Proof Let $\mathcal{A} = (A, F)$ be a finite algebra of finite similarity type. Choose a positive integer $n \geq 2$ such that $\text{card } F \leq n$ and $\text{arity } f \leq n$ for all $f \in F$.

We can write a finite sequence f_1, f_2, \dots, f_n , where $\{f_1, f_2, \dots, f_n\} = F$ and consider every f_i as an n -ary operation on A .

Introduce one binary and two unary operations on $C = A^n$ as follows:

$$\begin{aligned} x \cdot y &= (x_1, y_1, y_2, \dots, y_{n-1}), \\ g(x) &= (f_1(x), f_2(x), \dots, f_n(x)), \\ h(x) &= (x_2, x_3, \dots, x_n, x_1). \end{aligned} \quad (1)$$

Then $\mathcal{C} = (C, \{\cdot, g, h\})$ is a finite algebra of type $(2, 1, 1)$. For $x^{(1)}, x^{(2)}, \dots, x^{(k)} \in C$ ($k \geq 2$) we can put inductively

$$x^{(1)} \cdot x^{(2)} \dots x^{(k)} = x^{(1)} \cdot (x^{(2)} \cdot x^{(3)} \dots x^{(k)}). \quad (2)$$

Define the map $\varphi : \text{Tol } \mathcal{A} \rightarrow \text{Tol } \mathcal{C}$. Let $T \in \text{Tol } \mathcal{A}$. We put

$$\begin{aligned} \langle (x_1, x_2, \dots, x_n), (y_1, y_2, \dots, y_n) \rangle &\in \varphi(T) \\ &\text{if and only if } \langle x_i, y_i \rangle \in T \text{ for } i = 1, 2, \dots, n. \end{aligned} \quad (3)$$

Clearly $\varphi(T)$ is a reflexive and symmetric binary relation on C and from (1) it is easy to show that $\varphi(T) \in \text{Tol } \mathcal{C}$.

Evidently, for $S, T \in \text{Tol } \mathcal{A}$ we have $S \subseteq T$ if and only if $\varphi(S) \subseteq \varphi(T)$ and so φ is an injection. Now, it remains to prove that φ is a surjection. For every $x = (x_1, x_2, \dots, x_n) \in C$ we put $I(x) = x_1$. Suppose that $R \in \text{Tol } \mathcal{C}$. Let $T \subseteq A \times A$ such that

$$\langle u, v \rangle \in T \text{ if and only if there is } \langle x, y \rangle \in R \text{ and } I(x) = u, I(y) = v. \quad (4)$$

Clearly T is a reflexive and symmetric binary relation on A .

First we shall show that

$$\begin{aligned} \langle (x_1, x_2, \dots, x_n), (y_1, y_2, \dots, y_n) \rangle &\in R \\ &\text{whenever } \langle x_i, y_i \rangle \in T \text{ for all } i = 1, 2, \dots, n. \end{aligned} \quad (5)$$

Assume that $\langle x_i, y_i \rangle \in T$. Then there exist $x^{(i)}, y^{(i)} \in C$ such that $\langle x^{(i)}, y^{(i)} \rangle \in R$, $I(x^{(i)}) = x_i$ and $I(y^{(i)}) = y_i$. It follows from (1) and (2) that $x = (x_1, x_2, \dots, x_n) = x^{(1)} \cdot x^{(2)} \dots x^{(n)}$, $y = (y_1, y_2, \dots, y_n) = y^{(1)} \cdot y^{(2)} \dots y^{(n)}$ and so $\langle x, y \rangle \in R$.

Now we shall show that $T \in \text{Tol } \mathcal{A}$. Let $\langle x_i, y_i \rangle \in T$ for $i = 1, 2, \dots, n$. It follows from (5) that $\langle x, y \rangle \in R$, where $x = (x_1, x_2, \dots, x_n)$, $y = (y_1, y_2, \dots, y_n)$. According to (1), we obtain $\langle g(x), g(y) \rangle \in R$ and so $\langle f_1(x), f_1(y) \rangle \in T$. For $k = 1, 2, \dots, n-1$ we have $\langle h^k g(x), h^k g(y) \rangle \in R$ and so $\langle f_i(x), f_i(y) \rangle \in T$ for $i = 2, 3, \dots, n$. Thus $T \in \text{Tol } \mathcal{A}$.

Finally we shall show that $R = \varphi(T)$. Let

$$\langle x, y \rangle = \langle (x_1, x_2, \dots, x_n), (y_1, y_2, \dots, y_n) \rangle \in R,$$

then according to (1), we have $\langle h^k(x), h^k(y) \rangle \in R$ for $k = 1, 2, \dots, n-1$ and so $\langle x_i, y_i \rangle \in T$ for $i = 1, 2, \dots, n$. This means that $\langle x, y \rangle \in \varphi(T)$. Therefore $R \subseteq \varphi(T)$.

Assume that $\langle x, y \rangle \in \varphi(T)$, then by (3) we have $\langle x_i, y_i \rangle \in T$ for $i = 1, 2, \dots, n$, where $x = (x_1, x_2, \dots, x_n)$ and $y = (y_1, y_2, \dots, y_n)$. According to (5), we have $\langle x, y \rangle \in R$. Consequently we obtain $\varphi(T) \subseteq R$ and so $R = \varphi(T)$.

Note It is well known that the set $\text{Con } \mathcal{A}$ of all congruences on an algebra \mathcal{A} is a subset of $\text{Tol } \mathcal{A}$ (the set of all transitive tolerances on \mathcal{A}). From (3) it follows that the map φ has the property:

$$\varphi(\text{Con } \mathcal{A}) \subseteq \text{Con } \mathcal{C}.$$

Now, we shall prove that

$$\varphi(\text{Con } \mathcal{A}) = \text{Con } \mathcal{C}.$$

It suffices to show that T defined by (4) is a congruence on \mathcal{A} , whenever $R \in \text{Con } \mathcal{C}$.

Suppose that $\langle u, v \rangle, \langle v, w \rangle \in T$ and $R \in \text{Con } \mathcal{C}$. According to (4), there exist $x, y^{(1)}, y^{(2)}, z \in C$ such that $\langle x, y^{(1)} \rangle, \langle y^{(2)}, z \rangle \in R$ and $I(x) = u, I(y^{(1)}) = v = I(y^{(2)}), I(z) = w$. It is easy to show that by (1) and (2) we have $x^n = x \cdot x \dots x = (u, u, \dots, u)$. Analogously we can obtain that $(y^{(1)})^n = (v, v, \dots, v) = (y^{(2)})^n$ and $z^n = (w, w, \dots, w)$. We have $\langle x^n, (y^{(1)})^n \rangle, \langle (y^{(2)})^n, z^n \rangle \in R$ and so $\langle x^n, z^n \rangle \in R$ and $I(x^n) = u, I(z^n) = w$. Hence, by (4), we get $\langle u, w \rangle \in T$. Consequently $T \in \text{Con } \mathcal{A}$.

Corollary For every finite algebra \mathcal{A} of a finite similarity type there exists a finite algebra \mathcal{C} of type $(2, 1, 1)$ such that $\text{Con } \mathcal{A} \cong \text{Con } \mathcal{C}$.

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