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*Archivum Mathematicum*, Vol. 33 (1997), No. 4, 279--299

Persistent URL: <http://dml.cz/dmlcz/107617>

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## A MULTIPLICATION OF E-VARIETIES OF REGULAR $\mathcal{E}$ -SOLID SEMIGROUPS BY INVERSE SEMIGROUP VARIETIES

MARTIN KUŘIL

ABSTRACT. A multiplication of e-varieties of regular  $E$ -solid semigroups by inverse semigroup varieties is described both semantically and syntactically. The associativity of the multiplication is also proved.

### 1. INTRODUCTION

We investigate here an operator on the lattice of all e-varieties of regular semigroups. In [7] we defined semantically a partial multiplication on this lattice:  $\mathcal{U} \square \mathcal{V}$  is defined if  $\mathcal{U}$  is an e-variety of regular semigroups and  $\mathcal{V}$  is an e-variety of inverse semigroups. The definition is based on a certain semidirect product of regular semigroups by inverse semigroups. In the case that  $\mathcal{U}$  is an e-variety of orthodox semigroups we also described our multiplication syntactically in terms of biinvariant congruences for orthodox semigroups introduced in [5] by Kadourek and Szendrei.

In this paper we present a syntactical description of our multiplication in the case that the first factor is an e-variety of regular  $\mathcal{E}$ -solid semigroups. The description is essentially based on the notion of biinvariant congruences for regular  $\mathcal{E}$ -solid semigroups given in [6] by Kadourek and Szendrei. Moreover, we prove the associativity:  $\mathcal{U} \square (\mathcal{V} \square \mathcal{W}) = (\mathcal{U} \square \mathcal{V}) \square \mathcal{W}$  for any e-variety  $\mathcal{U}$  of regular  $\mathcal{E}$ -solid semigroups and any inverse semigroup varieties  $\mathcal{V}, \mathcal{W}$ .

For basic notions in the theory of semigroups the reader is referred to [4].

### 2. SEMANTICS

Let  $S = (S, \cdot)$  be a semigroup. The set of all endomorphisms of  $S$  is denoted by  $\text{End}(S)$ . Let  $E(S)$  stand for the set of all idempotents of  $S$ . Denote by  $\langle E(S) \rangle$  the subsemigroup of  $S$  generated by  $E(S)$  provided that  $E(S) \neq \emptyset$ . Clearly, for

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1991 *Mathematics Subject Classification*: 20M07, 20M17.

*Key words and phrases*: regular semigroup, inverse semigroup, e-variety, biinvariant congruence, bifree object.

Received January 19, 1996.

The author acknowledges the support of the Grant no. 201/93/2121 of the Grant Agency of the Czech Republic.

any element  $a \in S$ , there is at most one element  $b \in S$  satisfying  $ab = ba = a$ . If such an element  $b$  really exists then we denote it by  $a^{-1}$ .

In [7] we used the following non-standard semidirect product of semigroups:

Let  $(S, \cdot)$  be an inverse semigroup. For  $E \subseteq S$ , the unique inverse of  $a \in E$  is denoted by  $a'$ . Let  $\alpha : (S, \cdot) \rightarrow (\text{End}(S) \circ)$ , where  $\circ$  is the composition  $(\alpha \circ \beta)(a) = \alpha(\beta(a))$  ( $\alpha \in \text{End}(S)$ ,  $\beta \in \text{End}(S)$ ), be a homomorphism.

Put  $S \times_\varphi E = \{(a, e) \in S \times E \mid (a')(e) = a\}$  and define

$$(a, e) \cdot (b, f) = ( (a')(b) \cdot (e)(f) \cdot a, e \cdot f )$$

for  $(a, e), (b, f) \in S \times_\varphi E$ .

**2.1 Result.** ([7], 2.1 Lemma, 2.2 Lemma)

- (i)  $(S \times_\varphi E, \cdot)$  is a semigroup
- (ii) If  $S$  is regular, then  $S \times_\varphi E$  is also regular.

Notice that this non-standard semidirect product of semigroups is in essence the so called  $\alpha$ -semidirect product of inverse semigroups introduced by Billhardt in [1].

**2.2 Result.** ([7], 2.3 Lemma) *Let  $(a, e) \in S \times_\varphi E$ . Then  $(a, e)$  is an idempotent in  $S \times_\varphi E$  if and only if  $a \in E$  and  $e \in E$ .*

**2.3 Lemma.** *Let  $(a, e) \in S \times_\varphi E$ . Then  $(a, e) \in (S \times_\varphi E)^*$  if and only if  $a \in E$  and  $e \in E$ .*

**Proof.**

1. Let  $(a, e) \in (S \times_\varphi E)^*$ . Then  $(a, e) = (a_1, e_1) \cdot (a_k, e_k)$  for some  $(a_1, e_1), (a_k, e_k) \in (S \times_\varphi E)$ . We know that  $a_1, a_k \in E$  and  $e_1, e_k \in E$  (see 2.2). Put  $(a_i, e_i) = (a_1, e_1) \cdot (a_i, e_i) \cdot (a_1, e_1)$ . We will show that  $a_i \in E$  and  $e_i \in E$  ( $a_1, e_1 \in E$ ). Clearly,  $a_1 \in E$ ,  $e_1 \in E$ . Let  $1 \leq i \leq k$  and  $a_{i-1} \in E$ ,  $e_{i-1} \in E$ . We have  $(a_{i-1}, e_{i-1}) \cdot (a_i, e_i) \in (S \times_\varphi E)$ , since  $(a_{i-1}, e_{i-1}) \in \text{End}(S)$ . Further,  $(a_{i-1}, e_{i-1}) \cdot (a_i, e_i) \in (S \times_\varphi E)$ . We see that  $a_i = (a_{i-1}, e_{i-1}) \cdot (a_i, e_i) \in E$ . Finally,  $e_i = e_{i-1} \cdot e_i \in E$ .
2. Let  $(a, e) \in (S \times_\varphi E)$  and  $(a, e) \in (S \times_\varphi E)^*$ . Then  $(a, e) = (a_1, e_1) \cdot (a_k, e_k)$  for some  $(a_1, e_1), (a_k, e_k) \in (S \times_\varphi E)$ . Put  $(a_i, e_i) = (a')(e_i) \cdot (a_1, e_1)$ . Clearly,  $a_i \in E$  ( $a_1, e_1 \in E$ ) and  $e_i = a_1, e_k$ . Further,  $(a_i, e_i) \in S \times_\varphi E$ , since  $(a')(e_i) = a_i \cdot (a_1, e_1)$ . Using 2.2 we obtain  $(a_i, e_i) \in (S \times_\varphi E)^*$  ( $a_1, e_1 \in E$ ). We will prove that  $(a_1, e_1) \cdot (a_i, e_i) = (a_1, e_1) \cdot (a_i, e_i)$ . Let  $1 \leq i \leq k$ . Then
 
$$\begin{aligned} (a_1, e_1) \cdot (a_{i-1}, e_{i-1}) \cdot (a_i, e_i) &= (a_1, e_{i-1}) \cdot (a_i, e_i) \\ &= ( (a_1)'(e_{i-1}) \cdot (a_i, e_i) ) \cdot (a_i, e_i)^2 \\ &= ( (a_1)'(e_{i-1}) \cdot (a_i, e_i) ) \cdot (a_i, e_i) \\ &= (a_1, e_i) \end{aligned}$$

For  $i = k$  we get  $(a, e) = (a_1, e_k) = (a_1, e_1) \cdot (a_k, e_k) \in (S \times_\varphi E)^*$ . □

A semigroup  $(S, \cdot)$  is called regular  $\alpha$ -solid if it is regular and  $(S, \cdot)$  is completely regular.

**2.4 Lemma.** *If  $\mathcal{V}$  is regular  $\mathcal{R}$ -solid, then  $\mathcal{V} \times_{\varphi}$  is also regular  $\mathcal{R}$ -solid.*

**Proof.** We know that  $\mathcal{V} \times_{\varphi}$  is regular (see 2.1(ii)). We have to show that  $(\mathcal{V} \times_{\varphi})$  is completely regular. Let  $(a, b) \in (\mathcal{V} \times_{\varphi})$ . Then  $a \in \mathcal{V}$ ,  $b \in \mathcal{R}$ , by 2.3. Since  $\mathcal{V}$  is regular  $\mathcal{R}$ -solid, there exists  $c \in \mathcal{V}$  such that  $ac = ca = a$ . Put  $d = (c')(b)$ . Then  $d \in \mathcal{R}$  and  $ad = da = a$ . Clearly,  $(a, b) \in \mathcal{V} \times_{\varphi} \mathcal{R}$ . Using 2.3 we obtain  $(a, b) \in (\mathcal{V} \times_{\varphi})$ . Further,  

$$\begin{aligned} (a, b)(c, d) &= ((c')(b))(c, d) \cdot ((c')(b))^2 = (c, d)(c, d) \\ &= ((c')(b))(c, d) \cdot ((c')(b))^2 = (c, d) = (c, d) \end{aligned}$$

Similarly,  $(a, b)(c, d) = (a, b)$  and  $(c, d)(a, b) = (c, d)$ . □

For any class  $\mathcal{V}$  of regular semigroups, we will denote by  $\mathcal{H}(\mathcal{V})$ ,  $\mathcal{R}(\mathcal{V})$  and  $\mathcal{D}(\mathcal{V})$ , respectively, the classes of all homomorphic images, regular subsemigroups and direct products of semigroups in  $\mathcal{V}$ .

We adopt the following notations for classes of regular semigroups:

- R** — the class of all regular semigroups;
- ES** — the class of all regular  $\mathcal{R}$ -solid semigroups;
- I** — the class of all inverse semigroups.

A class  $\mathcal{V} \subseteq \mathbf{R}$  satisfying  $\mathcal{H}(\mathcal{V}) \subseteq \mathcal{V}$ ,  $\mathcal{R}(\mathcal{V}) \subseteq \mathcal{V}$  and  $\mathcal{D}(\mathcal{V}) \subseteq \mathcal{V}$  is called an e-variety. The classes **R**, **ES**, **I** are examples of e-varieties. The concept of e-variety was introduced by Hall in [3]. Simultaneously and independently Kađourek and Szendrei in [5] have considered e-varieties of orthodox semigroups, which they called bivarieties of orthodox semigroups.

Denote by  $\langle \mathcal{V} \rangle$  the least e-variety of regular semigroups containing the class  $\mathcal{V} \subseteq \mathbf{R}$ .

Let  $\mathcal{U} \subseteq \mathbf{R}$  and  $\mathcal{V} \subseteq \mathbf{I}$  be e-varieties. In [7] we defined a multiplication  $\square$  in the following way:

$$\mathcal{U} \square \mathcal{V} = \{ \langle \mathcal{V} \times_{\varphi} \mathcal{U} \mid \mathcal{U} \in \mathcal{U}, \mathcal{V} \in \mathcal{V} : (\cdot) \rightarrow (\text{End}(\cdot) \circ) \text{ is a homomorphism} \rangle \}$$

**2.5 Result.** ([7], 2.5 Lemma) *Let  $\mathcal{V} \neq \emptyset$ . Let  $\mathcal{V}_i$  be a semigroup for  $i \in I$ . Let  $\mathcal{V}_i$  be an inverse semigroup for  $i \in I$ . Finally, let  $\mathcal{V}_i : (\mathcal{V}_i \cdot) \rightarrow (\text{End}(\mathcal{V}_i) \circ)$  be a homomorphism for  $i \in I$ . Then*

$$\prod_{i \in I} (\mathcal{V}_i \times_{\varphi_i} \mathcal{V}_i) \cong \prod_{i \in I} \mathcal{V}_i \times_{\varphi} \prod_{i \in I} \mathcal{V}_i$$

where the homomorphism

$$: \left( \prod_{i \in I} \mathcal{V}_i \cdot \right) \rightarrow \left( \text{End} \left( \prod_{i \in I} \mathcal{V}_i \right) \circ \right)$$

is given by

$$((\mathcal{V}_i)_{i \in I})((\mathcal{V}_i)_{i \in I}) = (\mathcal{V}_i(\mathcal{V}_i)(\mathcal{V}_i))_{i \in I}$$

The isomorphism is given by

$$((\mathcal{V}_i \mathcal{V}_i))_{i \in I} \mapsto ((\mathcal{V}_i)_{i \in I} (\mathcal{V}_i)_{i \in I})$$

**2.6 Result.** ([6], Proposition 2.3) *Let  $\mathcal{V} \subseteq \mathbf{ES}$ . Then*

- (i)  $\text{r}(\mathcal{V}) \subseteq \text{r}(\mathcal{V})$
- (ii)  $\langle \mathcal{V} \rangle = \text{r}(\mathcal{V})$

**2.7 Lemma.** *Let  $\mathcal{U} \subseteq \mathbf{ES}$  and  $\mathcal{V} \subseteq \mathbf{I}$  be e-varieties. Then*

$$\mathcal{U} \square \mathcal{V} = \text{r}(\{ \times_{\varphi} \mid \in \mathcal{U} \in \mathcal{V} : (\cdot) \rightarrow (\text{End}(\cdot) \circ) \text{ is a homomorphism} \})$$

**Proof.** Put  $\mathcal{W} = \{ \times_{\varphi} \mid \in \mathcal{U} \in \mathcal{V} : (\cdot) \rightarrow (\text{End}(\cdot) \circ) \text{ is a homomorphism} \}$ . It follows from 2.4 that  $\mathcal{W} \subseteq \mathbf{ES}$ . Then  $\mathcal{U} \square \mathcal{V} = \text{r}(\mathcal{W})$  by 2.6(ii). It is clear that  $\text{r}(\mathcal{W}) \subseteq \text{r}(\mathcal{W})$ . Further,  $(\mathcal{W}) \subseteq (\mathcal{W})$ , by 2.5. Then  $\text{r}(\mathcal{W}) \subseteq \text{r}(\mathcal{W})$ . This together with 2.6(i) gives  $\text{r}(\mathcal{W}) \subseteq \text{r}(\mathcal{W})$ . □

### 3. SYNTAX

Recall the notions of biidentities and biinvariant congruences in the class of regular  $\text{-solid}$  semigroups introduced by Kađourek and Szendrei in [6].

A unary semigroup is an algebra  $= (\cdot)'$  with an associative multiplication and with a unary operation  $'$ .

Let  $\text{ be a non-empty set. We add new symbols } (\text{ and } )'$  to the set  $\text{ and obtain a set } \text{ }_0 = \cup \{ (\cdot)'\}$ . Let us denote the free semigroup on the alphabet by  $\text{ }^+$ . Let  $(\cdot)$  be the smallest one among the subsets  $\text{ in } \text{ }_0^+$  which satisfy

- (i)  $\subseteq$
- (ii)  $\in$  implies  $\in$
- (iii)  $\in$  implies  $(\cdot)' \in$ .

The set  $(\cdot)$  will be often considered as a unary semigroup with a binary operation given by the concatenation of words and with a unary operation  $' : (\cdot) \rightarrow (\cdot)'$  given by  $\mapsto (\cdot)'$ . The unary semigroup  $(\cdot)$  is the free unary semigroup on the set  $\text{.$

In order to reduce the number of brackets in formulas, we will omit them if it causes no confusion. For example, we will often write  $'$  instead of  $(\cdot)'$ .

Consider a set  $'$  disjoint from  $\text{ and a bijection } ' : \rightarrow ' \mapsto '.$  The union  $\cup '$  will be denoted by  $\overline{\text{}}$ . For any  $\in \text{, we will identify } (\cdot)'$  with  $'$ , and so  $\overline{\text{}}$  becomes a subset in  $(\cdot)$ .

If  $\text{ is an inverse semigroup and } \in \text{ then the unique inverse of } \text{ is denoted by } '.$  In this way a unary operation  $'$  on  $\text{ is given and the inverse semigroup } = (\cdot)$  can be considered as a unary semigroup  $= (\cdot)'$ . Moreover, this unary semigroup satisfies the identities

- (id 1)  $(\cdot)' =$
- (id 2)  $(\cdot)' = ' '$
- (id 3)  $' =$
- (id 4)  $' ' = ' '.$



the smallest subsemigroup in  $\overline{S}$  containing the set  $\overline{S^{-1}}$  and closed under the partial operation assigning the word  $(\ )^{-1}$  to any word  $\alpha$  with  $\alpha(\ ) = 1$  (see [6], Section 2, for the definition of  $(\ )$ ). If we consider the unary homomorphism  $\overline{\alpha} : \overline{S} \rightarrow \overline{S}$  extending the mapping  $\alpha \mapsto \overline{\alpha} \mapsto \alpha' = (\ )' (\ \in \overline{S})$  then, for any  $\alpha \in \overline{S}$ , the condition  $\alpha(\ ) = 1$  is equivalent to  $\alpha(\ ) (\ ) = 1(\ )$  and the restriction of  $\overline{\alpha}$  to  $\overline{S^\infty}(\ )$  is an isomorphism between  $\overline{S^\infty}(\ )$  and  $r(\ )$ .

If  $(\ )$  is a regular semigroup, then a mapping  $\overline{\alpha} : \overline{S} \rightarrow \overline{S}$  is called matched if  $(\ ) \cdot (\ )' \cdot (\ ) = (\ )$  and  $(\ )' \cdot (\ ) \cdot (\ )' = (\ )'$  for all  $\alpha \in \overline{S}$ .

To any matched mapping  $\overline{\alpha} : \overline{S} \rightarrow \overline{S}$ , where  $\overline{S}$  is a regular  $\overline{S}$ -solid semigroup, we now define a homomorphism  $\alpha : r(\ ) \rightarrow r(\ )$  as follows. We proceed by induction with respect to the complexity of words from  $r(\ )$ , and we put

- (i)  $\alpha(\ ) = (\ ) (\ \in \overline{S})$
- (ii)  $\alpha(\ ) = (\ ) (\ ) (\ \in r(\ ))$
- (iii)  $\alpha(\ )' = (\ (\ ))^{-1} (\ \in r(\ ) \ (\ ) = 1(\ ))$ ,

where  $(\ (\ ))^{-1}$  denotes the group inverse of  $(\ )$  in the maximal subgroup of  $\overline{S}$  containing  $(\ )$ . Of course, we must show that this  $\alpha(\ )$  really lies in a subgroup of  $r(\ )$ . This will be the content of the next result. We will then see that  $\alpha$  is well defined and we will call the homomorphism  $\alpha$  the extension of the matched mapping  $\overline{\alpha} : \overline{S} \rightarrow \overline{S}$  to  $r(\ )$ .

**3.2 Result.** ([6], Lemma 2.1) *The above definition is correct, that is, for any  $\alpha \in r(\ )$  with  $\alpha(\ ) = 1(\ )$ , the element  $\alpha(\ )$  lies in a subgroup of  $r(\ )$ , provided that  $\overline{S}$  is a regular  $\overline{S}$ -solid semigroup.*

By a biidentity over  $\overline{S}$  we will mean any pair  $\alpha = \beta$  of words  $\alpha, \beta \in r(\ )$ . We will say that a biidentity  $\alpha = \beta$  is satisfied in a regular  $\overline{S}$ -solid semigroup  $\overline{S}$  if, for any matched mapping  $\overline{\alpha} : \overline{S} \rightarrow \overline{S}$ , we have  $\alpha(\ ) = \beta(\ )$  where  $\alpha : r(\ ) \rightarrow r(\ )$  is the extension of  $\overline{\alpha}$  to  $r(\ )$ . As usual, we will say that a biidentity is satisfied in a class  $\mathcal{V}$  of regular  $\overline{S}$ -solid semigroups if it is satisfied in each member of  $\mathcal{V}$ .

Given a class  $\mathcal{V}$  of regular  $\overline{S}$ -solid semigroups, put

$$(\mathcal{V}) = \{(\alpha = \beta) \in r(\ ) \times r(\ ) \mid \text{the biidentity } \alpha = \beta \text{ is satisfied in } \mathcal{V}\}$$

For any set  $\Sigma \subseteq r(\ ) \times r(\ )$  of biidentities, put

$$[\Sigma] = \{ \alpha \in \mathbf{ES} \mid \alpha \text{ satisfies all biidentities in } \Sigma \}$$

We will write  $(\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n)$  to indicate that only elements  $\alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n \in \overline{S}$  may occur in  $\alpha \in r(\ )$ . If  $\alpha = (\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n) \in r(\ )$  and  $\beta = (\ \beta_1 \ \beta'_1 \ \dots \ \beta_n \ \beta'_n) \in r(\ )$  then  $(\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n)$  is obtained by substituting  $\beta_1 \ \beta'_1 \ \dots \ \beta_n \ \beta'_n$  into  $\alpha$  for  $\alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n$  respectively. It is clear that  $(\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n) \in (\ )$ . It is easy to see that if  $(\ \alpha_i \ \alpha'_i) (\ ) = 1(\ ) (1 \leq i \leq n)$  then  $(\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n) \in r(\ )$ .

A congruence  $\sim$  on  $r(\ )$  will be called biinvariant if  $(\mathbf{ES}) \subseteq \sim$  and it has the following property: whenever  $\alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n \in r(\ )$  such that

$$(\ \alpha_1 \ \alpha'_1 \ \dots \ \alpha_n \ \alpha'_n) \sim (\ \beta_1 \ \beta'_1 \ \dots \ \beta_n \ \beta'_n)$$

and

$$i i i (\mathbf{ES} ) i i i i (\mathbf{ES} ) i \text{ for } = 1 2$$

then also

$$( 1 1 n n ) ( 1 1 n n )$$

Observe that the second assumption implies  $( i i ) ( ) = 1( )$  for  $= 1$ , as the class of all groups is contained in **ES**, so that this definition is correct.

The set of all fully invariant congruences on the unary semigroup  $( )$  will be denoted by  $( )$  and the set of all biinvariant congruences on the semigroup  $r( )$  will be denoted by  $r( )$ .

Now, we can present the syntax of our multiplication.

Let  $= \{ 1 2 \}$  be a set of variables. Given  $\in ( )$ , define a new alphabet  $\rho = ( ) \times$ .

Define a left action  $*$  of  $( )$  on  $( \rho )$  by

$$\begin{aligned} * ( ) &= ( ) \\ * &= ( * )( * ) \\ * ' &= ( * )' \end{aligned}$$

for  $\in ( ) \in ( \rho )$ .

We will frequently use the following lemma without references.

**3.3 Lemma.** *Let  $\in ( ) \in ( )$ . Then*

- (i) *implies  $* = *$  for all  $\in ( \rho )$*
- (ii)  *$*( * ) = ( )*$  for all  $\in ( \rho )$ .*

**Proof.** The assertions are clear. □

**3.4 Lemma.** *Let  $\in ( ) \in ( )$ . If  $\in r( \rho )$ , then  $* \in r( \rho )$ .*

**Proof.** By induction with respect to  $.$  Let  $\in ( ) \in$ . Then  $*( ) = ( ) \in r( \rho )$   $*( )' = ( )' \in r( \rho )$ .

Let  $\in r( \rho )$   $* \in r( \rho )$ . Then  $* = ( * )( * ) \in r( \rho )$ .

Let  $\in r( \rho )$   $( \rho ) = 1( \rho )$   $* \in r( \rho )$ . Then  $*' = ( * )' \in r( \rho )$ , since  $( * ) ( \rho ) = 1( \rho ) ( ( \rho ) \in ( \rho )$  and  $\mapsto *$  is an endomorphism on  $( \rho )$ ). □

Now, let  $\in ( ) \supseteq ( )$ . Define

$$\rho : ( ) \rightarrow ( \rho )$$

by

$$\begin{aligned} \rho ( ) &= ( ' ) \\ \rho ( ) &= ( ' ' * \rho ( ) )( * \rho ( ) ) \\ \rho ( ' ) &= ' * ( \rho ( ) )' \end{aligned}$$

where  $\in ( ) \in ( )$ .



**3.5 Remark.** The mapping  $\rho$  is defined unambiguously:

Let  $\alpha \in (\mathcal{A})$ . Suppose that the values  $\rho(\alpha)$ ,  $\rho(\alpha')$ ,  $\rho(\alpha'')$ ,  $\rho(\alpha''')$  and  $\rho(\alpha''''')$  are determined unambiguously. We show that  $\rho((\alpha)) = \rho(\alpha)$ :

$$\begin{aligned} \rho((\alpha)) &= ((\alpha) \alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') \\ \rho(\alpha) &= (\alpha) (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha''''') (\alpha'''' \alpha'''''). \end{aligned}$$

The following lemma will be also often used without references.

**3.6 Lemma.** Let  $\alpha \in (\mathcal{A}) \supseteq (\mathcal{B}) \in (\mathcal{A})$ . Then  $\alpha' * \rho(\alpha) = \rho(\alpha)$ .

**Proof.** By induction with respect to  $\alpha$ . Let  $\alpha \in \mathcal{A}$ . Then

$$\begin{aligned} \alpha' * \rho(\alpha) &= \alpha' * (\alpha') \\ &= (\alpha' \alpha') \\ &= (\alpha') \\ &= \rho(\alpha). \end{aligned}$$

Let  $\alpha \in (\mathcal{A})$ ,  $\alpha' * \rho(\alpha) = \rho(\alpha)$ . Then

$$\begin{aligned} (\alpha') * \rho(\alpha) &= \alpha' \alpha'' \alpha''' \alpha'''' \alpha'''''' (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha'''' \alpha''''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha'''' \alpha''''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha'''' \alpha''''''') (\alpha'''' \alpha''''') \\ &= (\alpha' \alpha'' \alpha''' \alpha'''' \alpha''''''') (\alpha'''' \alpha''''') \\ &= \rho(\alpha). \end{aligned}$$

Let  $\alpha \in (\mathcal{A})$ . Then

$$\begin{aligned} (\alpha')' * \rho(\alpha') &= \alpha' * (\alpha' * (\rho(\alpha)))' \\ &= \alpha' * (\alpha' * (\rho(\alpha)))' \\ &= \alpha' * (\rho(\alpha))' \\ &= \rho(\alpha'). \end{aligned}$$

□

**3.7 Lemma.** Let  $\alpha \in (\mathcal{A}) \supseteq (\mathcal{B}) \in (\mathcal{A})$ . If  $(\alpha)$ , then  $\rho(\alpha) (\rho) \rho(\alpha)$ .

**Proof.** Having in mind that the variety **I** is the class of all unary semigroups satisfying the identities (id 1) — (id 4) we prove the lemma in the following seven steps.

1.  $\rho((\alpha')) (\rho) \rho(\alpha) (\alpha \in (\mathcal{A}))$   

$$\begin{aligned} \rho((\alpha')) &= (\alpha') * (\rho(\alpha'))' \\ &= \alpha' * (\alpha' * (\rho(\alpha)))' \\ &= ((\alpha' * \rho(\alpha)))' \\ &= ((\rho(\alpha)))' \\ &= (\rho) \rho(\alpha). \end{aligned}$$

$$\begin{aligned}
 2. \quad & \rho((\quad)') (\quad) \rho(\quad) (\quad) (\quad) \in (\quad) \\
 & \rho((\quad)') = (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' * ((\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad))' \\
 & = ((\quad)' * (\quad) (\quad)) (\quad)' * (\quad) (\quad)' \\
 & \quad (\quad) (\quad)' * (\quad) (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad)' * (\quad) (\quad)' (\quad)' * (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = \rho(\quad)'.
 \end{aligned}$$

$$\begin{aligned}
 3. \quad & \rho(\quad)' (\quad) \rho(\quad) (\quad) (\quad) \in (\quad) \\
 & \rho(\quad)' = ((\quad) (\quad) (\quad))' * (\quad) (\quad) (\quad) * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad) (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' \\
 & = \rho(\quad) (\quad)' (\quad)' * (\quad) (\quad)' \rho(\quad) \\
 & = \rho(\quad) (\quad) (\quad)' \rho(\quad) \\
 & \quad (\quad) \rho(\quad)'.
 \end{aligned}$$

$$\begin{aligned}
 4. \quad & \rho(\quad)' (\quad)' (\quad) \rho(\quad) \rho(\quad)' (\quad)' (\quad) \in (\quad) \\
 & \text{Let } (\quad) \in (\quad). \text{ Then} \\
 & \rho(\quad)' = (\quad)' (\quad)' (\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad) (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = \rho(\quad) (\quad) (\quad)'.
 \end{aligned}$$

Further,

$$\begin{aligned}
 \rho(\quad)' (\quad)' & = (\quad)' (\quad)' (\quad)' (\quad)' * (\quad) (\quad) (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' (\quad)' * (\quad) (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' \rho(\quad) (\quad)' (\quad)' \\
 & = (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' \rho(\quad) (\quad)' \\
 & \quad (\quad) (\quad)' (\quad)' * (\quad) (\quad)' (\quad)' \rho(\quad) (\quad)' \\
 & = \rho(\quad)' (\quad)'.
 \end{aligned}$$

$$\begin{aligned}
 5. \quad & (\quad) \text{ and } \rho(\quad) (\quad) \rho(\quad) \text{ implies } \rho(\quad) (\quad) \rho(\quad) (\quad) \in (\quad) \\
 & \rho(\quad) = (\quad)' (\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad) \\
 & \rho(\quad) = (\quad)' (\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad) \\
 & \text{Since } (\quad), \text{ we have } (\quad)' (\quad)' * (\quad) (\quad) = (\quad)' (\quad)' * (\quad) (\quad). \\
 & \text{Since } \rho(\quad) (\quad) \rho(\quad), \text{ we have } (\quad)' (\quad)' * (\quad) (\quad) (\quad) = (\quad)' (\quad)' * (\quad) (\quad). \\
 & \text{So, } \rho(\quad) (\quad) \rho(\quad).
 \end{aligned}$$

$$\begin{aligned}
 6. \quad & (\quad) \text{ and } \rho(\quad) (\quad) \rho(\quad) \text{ implies } \rho(\quad) (\quad) \rho(\quad) (\quad) \in (\quad) \\
 & \rho(\quad) = (\quad)' (\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad) \\
 & \rho(\quad) = (\quad)' (\quad)' * (\quad) (\quad) (\quad) * (\quad) (\quad) \\
 & \text{Since } \rho(\quad) (\quad) \rho(\quad), \text{ we have } (\quad)' (\quad)' * (\quad) (\quad) (\quad) = (\quad)' (\quad)' * (\quad) (\quad). \text{ Fur-} \\
 & \text{ther, from } (\quad) \text{ we get } (\quad)' (\quad)' * (\quad) (\quad) = (\quad)' (\quad)' * (\quad) (\quad). \text{ So, } (\quad)' (\quad)' * \\
 & \rho(\quad) (\quad) \rho(\quad) = (\quad)' (\quad)' * (\quad) (\quad). \text{ Finally, } (\quad)' (\quad)' * (\quad) (\quad) = (\quad)' (\quad)' * (\quad) (\quad), \text{ since } (\quad). \text{ Now,} \\
 & \text{we see that } \rho(\quad) (\quad) \rho(\quad).
 \end{aligned}$$

7.  $(\rho)$  and  $\rho(\rho) \in \rho(\rho)$  implies  $\rho(\rho) \in \rho(\rho)$   
 $\rho(\rho) = \rho * (\rho(\rho))'$   
 $\rho(\rho) = \rho * (\rho(\rho))'$

Since  $(\rho)$ , we get  $\rho(\rho) = \rho$  and then  $\rho * (\rho(\rho))' = \rho * (\rho(\rho))'$ . Since  $\rho(\rho) \in \rho(\rho)$ , we get  $(\rho(\rho))' \in \rho(\rho)$  and  $\rho * (\rho(\rho))' \in \rho * (\rho(\rho))'$ . So,  $\rho(\rho) \in \rho(\rho)$ .  $\square$

**3.8 Corollary.** Let  $\rho \in \rho(\rho) \supseteq \rho(\rho) \in \rho(\rho)$ . If  $\rho(\rho) = 1(\rho)$ , then  $\rho(\rho) \in \rho(\rho) = 1(\rho)$ .

**Proof.** Let  $\rho \in \rho(\rho) \supseteq \rho(\rho) \in \rho(\rho)$ . We know that  $\rho^2(\rho) \in \rho(\rho)$  (see 3.1). From 3.7 we get  $\rho(\rho^2) \in \rho(\rho)$ . Further,

$$\begin{aligned} \rho(\rho^2) &= (\rho * \rho(\rho))(\rho * \rho(\rho)) \\ &= (\rho * \rho(\rho))(\rho * \rho(\rho)) \\ &= (\rho * \rho(\rho))(\rho * \rho(\rho)) \\ &= (\rho(\rho))^2. \end{aligned}$$

Thus,  $(\rho(\rho))^2 \in \rho(\rho)$ .  $\square$

**3.9 Corollary.** Let  $\rho \in \rho(\rho) \supseteq \rho(\rho) \in \rho(\rho)$ . If  $\rho \in r(\rho)$ , then  $\rho(\rho) \in r(\rho)$ .

**Proof.** By induction with respect to  $\rho$ . Let  $\rho \in r(\rho)$ . Then  $\rho(\rho) = (\rho * \rho) \in r(\rho)$ ,  
 $\rho(\rho) = \rho * (\rho * \rho) = (\rho * \rho) = (\rho * \rho) \in r(\rho)$ .

Let  $\rho \in r(\rho)$   $\rho(\rho) \in r(\rho)$ .

$$\rho(\rho) = (\rho * \rho(\rho))(\rho * \rho(\rho))$$

We know that  $\rho * \rho(\rho) \in r(\rho)$   $\rho * \rho(\rho) \in r(\rho)$  (see 3.4). Thus,  $\rho(\rho) \in r(\rho)$ .

Let  $\rho \in r(\rho)$   $(\rho) = 1(\rho)$   $\rho(\rho) \in r(\rho)$ .

$$\rho(\rho) = \rho * (\rho(\rho))'$$

We know that  $\rho(\rho) \in \rho(\rho) = 1(\rho)$  (see 3.8). Then  $(\rho(\rho))' \in r(\rho)$ . Using 3.4 we obtain  $\rho(\rho) \in r(\rho)$ .  $\square$

Now, let  $\rho \in \rho(\rho) \supseteq \rho(\rho) \in r(\rho)$ . Define

$$\rho \subseteq r(\rho) \times r(\rho)$$

by

$$(\rho) \iff \rho \text{ and } \rho(\rho)$$

$(\rho \in r(\rho))$ .

The correctness of the definition is based on 3.9.

**3.10 Remark.** If  $\rho \in \rho(\rho) \supseteq \rho(\rho) \in r(\rho)$ , then  $\rho \in r(\rho)$ . We will prove it in 4.10.

4. RELATIONSHIPS BETWEEN SYNTAX AND SEMANTICS

**4.1 Result.** ([6], Corollary 2.11) *For any infinite set  $X$ , the rules*

$$\mathcal{V} \mapsto (\mathcal{V} \text{ on } X) \text{ and } \rho \mapsto [\rho]$$

*define mutually inverse order-reversing bijections between all e-varieties of regular  $X$ -solid semigroups and all biinvariant congruences on  $r(X)$ .*

We will denote the one-to-one correspondence from 4.1 simply by the symbol  $\leftrightarrow$ . Since it causes no confusion, we will use the symbol  $\leftrightarrow$  also for the well-known one-to-one correspondence between all varieties of unary semigroups and all fully invariant congruences on the free unary semigroup  $r(X)$ .

Now, we recall the notion of a bifree object. Let  $\mathcal{V}$  be a class of regular semigroups. By a bifree object in  $\mathcal{V}$  on a non-empty set  $X$ , we mean a pair  $(S, \rho)$  where  $S \in \mathcal{V}$  and  $\rho : \overline{X} \rightarrow S$  is a matched mapping such that the following universal property is satisfied: for any semigroup  $T \in \mathcal{V}$  and any matched mapping  $\sigma : \overline{X} \rightarrow T$ , there exists a unique homomorphism  $\tau : S \rightarrow T$  such that  $\tau \circ \rho = \sigma$ . In cases when the mapping  $\rho$  is obvious, we omit it and we term simply  $S$  to be a bifree object in  $\mathcal{V}$  on  $X$ . Note that in any class of regular semigroups, there exists, up to isomorphism, at most one bifree object on any non-empty set.

**4.2 Result.** ([6], Theorem 2.5) *If  $X$  is an infinite set and  $\mathcal{V}$  is a class of regular  $X$ -solid semigroups closed under taking regular subsemigroups and direct products then  $r(X)$  ( $\mathcal{V}$  on  $X$ ) is a bifree object in  $\mathcal{V}$  on  $X$ .*

**4.3 Lemma.** *Let  $(S, \rho) \in \mathcal{V}$  ( $\mathcal{V}$  on  $X$ ) and  $(T, \sigma) \in \mathcal{V}$  ( $\mathcal{V}$  on  $X$ ). Then the mapping*

$$\tau : (S, \rho) \rightarrow (\text{End}(T, \sigma), \sigma)$$

*given by*

$$\tau(s) = (s * t) \quad (s \in S, t \in r(X))$$

*is a correctly defined homomorphism.*

**Proof.**

1. correctness of the definition:

It follows from 3.4 that  $(S, \rho) \in \mathcal{V}$  and  $(T, \sigma) \in \mathcal{V}$  implies  $(S * T, \sigma) \in \mathcal{V}$ .

Now, let  $(S, \rho) \in \mathcal{V}$  and  $(T, \sigma) \in \mathcal{V}$ . We will show that  $(S * T, \sigma) \in \mathcal{V}$ .

We have  $\sigma * \rho = \sigma$ . Since  $(S, \rho) \in \mathcal{V}$ , we get  $(S * T, \sigma) \in \mathcal{V}$ . So,  $(S * T, \sigma) \in \mathcal{V}$ .

2.  $\tau : (S, \rho) \rightarrow (\text{End}(T, \sigma), \sigma)$  is an endomorphism (for any  $(S, \rho) \in \mathcal{V}$ ):

Let  $(S, \rho) \in \mathcal{V}$ . Then

$$\begin{aligned} \tau(s) \tau(t) &= (s * t) * (u * v) = (s * (t * u)) * v = (s * t) * (u * v) \\ &= ((s * t) * u) * v = (s * (t * u)) * v = (s * t) * (u * v). \end{aligned}$$

3.  $\tau$  is a homomorphism:

Let  $(S, \rho) \in \mathcal{V}$  and  $(T, \sigma) \in \mathcal{V}$ . Then

$$\begin{aligned} \tau(s) \tau(t) &= (s * t) * (u * v) = (s * (t * u)) * v = ((s * t) * u) * v \\ &= (s * t) * (u * v) = ((s * t) * u) * v. \end{aligned}$$

□

**4.4 Lemma.** Let  $\rho \in \text{ES}$ ,  $\rho' \in \text{ES}$ ,  $\rho \supseteq \rho'$ . Further, let  $\alpha: \text{End}(\rho) \rightarrow \text{End}(\rho')$  be the homomorphism from 4.3. Finally, let  $\beta: \text{End}(\rho) \rightarrow \text{End}(\rho) \times_{\varphi} \text{End}(\rho')$  be given by

$$\beta \mapsto (\rho(\beta)) (\beta \in \overline{\text{End}(\rho)})$$

Then

- (i)  $\beta$  is a matched mapping
- (ii)  $\beta(\rho) = (\rho(\beta))$  for all  $\beta \in \text{End}(\rho)$  (where  $\beta$  is the extension of the matched mapping  $\beta$ ).

**Proof.** Note that  $\rho \in \text{ES}$ ,  $\rho' \in \text{ES}$ ,  $\rho \supseteq \rho'$ .  $\beta \in \text{ES}$  (see 4.1, 4.2 and 2.4).

- (i) Choose  $\beta \in \text{End}(\rho)$ . Then  $\beta(\rho) (\beta') (\beta) =$   
 $= ((\beta' \beta)) ((\beta' \beta)') ((\beta' \beta))$   
 $= ((\beta' \beta)') ((\beta' \beta)) (\beta') ((\beta' \beta)') ((\beta' \beta))$   
 $= ((\beta' * (\beta' \beta))) ((\beta' * (\beta' \beta))') ((\beta' \beta))$   
 $= ((\beta' \beta)) ((\beta' \beta)') ((\beta' \beta))$   
 $= ((\beta' \beta)') ((\beta' \beta)) ((\beta' \beta)') ((\beta' \beta))$   
 $= ((\beta' * (\beta' \beta)) ((\beta' \beta)')) ((\beta' * (\beta' \beta)))$   
 $= ((\beta' \beta)) ((\beta' \beta)') ((\beta' \beta))$   
 $= ((\beta' \beta))$   
 $= (\beta)$ .

Similarly,  $\beta'(\rho') (\beta) (\beta') = (\beta')$ .

- (ii) We proceed by induction with respect to  $\rho$ . Let  $\beta \in \text{End}(\rho)$ .  $\beta(\rho) = (\rho(\beta))$   
 $= (\rho(\beta)) (\rho(\beta))$   
 $= ((\beta' \beta)') (\rho(\beta)) (\beta') (\rho(\beta))$   
 $= ((\beta' \beta) * \rho(\beta)) ((\beta' \beta) * \rho(\beta))$   
 $= (\rho(\beta))$ .

Let  $\beta \in \text{End}(\rho)$ .  $\beta(\rho) = 1(\beta) (\beta) = (\rho(\beta))$ . Note that  $\rho(\beta) \in \text{ES}$  by 3.9 and  $\rho(\beta) (\rho) = 1(\rho)$  by 3.8. We want to prove  $\beta'(\rho') = (\rho(\beta'))$ . In view of  $\beta'(\rho') = (\beta'(\rho'))^{-1}$  we have to show that

$$\begin{aligned} (\rho(\beta)) (\rho(\beta')) (\rho(\beta)) &= (\rho(\beta)) (\rho(\beta')) \\ (\rho(\beta')) (\rho(\beta)) (\rho(\beta')) &= (\rho(\beta')) (\rho(\beta')) \\ (\rho(\beta)) (\rho(\beta')) &= (\rho(\beta')) (\rho(\beta)). \end{aligned}$$

$$\begin{aligned} \text{We see that } (\rho(\beta)) (\rho(\beta')) (\rho(\beta)) &= \\ &= ((\beta' \beta)') (\rho(\beta)) (\beta') (\rho(\beta')) (\rho(\beta)) \\ &= ((\beta' * \rho(\beta)) (\beta' * \rho(\beta))) (\beta') (\rho(\beta)) \\ &= (\rho(\beta)) (\rho(\beta))' (\beta') (\rho(\beta)) \\ &= ((\beta' \beta)') (\rho(\beta)) (\rho(\beta))' (\beta') (\rho(\beta)) (\beta') \\ &= ((\beta' * \rho(\beta)) (\beta' * \rho(\beta))) (\beta' * \rho(\beta)) \\ &= (\rho(\beta)) (\rho(\beta))' \rho(\beta) \\ &= (\rho(\beta)) (\rho(\beta))', \text{ since } \rho \supseteq \text{ES } \rho \text{ and } \rho(\beta) (\rho(\beta))' \rho(\beta) \in \text{ES } \rho \text{ } \rho(\beta). \end{aligned}$$

Similarly,  $(\rho(\sigma) \sigma')(\rho(\sigma))(\rho(\sigma) \sigma') = (\rho(\sigma) \sigma')$ . Further,  
 $(\rho(\sigma))(\rho(\sigma) \sigma') =$   
 $= (\sigma' \sigma')(\rho(\sigma))(\sigma)(\rho(\sigma) \sigma')$   
 $= ((\sigma' * \rho(\sigma))(\sigma' * (\sigma' * (\rho(\sigma)))) \sigma')$   
 $= (\rho(\sigma))(\rho(\sigma))' \sigma')$   
 $(\rho(\sigma) \sigma')(\rho(\sigma)) =$   
 $= (\sigma' \sigma')(\rho(\sigma) \sigma')(\sigma')(\rho(\sigma)) \sigma'$ .  
 We know that  $\sigma^2(\sigma)$  (see 3.1). So,  $(\sigma) \sigma' \sigma'$ .  
 Then  $(\rho(\sigma) \sigma')(\rho(\sigma)) =$   
 $= ((\sigma' * (\sigma' * (\rho(\sigma))))(\sigma' * (\sigma' * \rho(\sigma)))) \sigma'$   
 $= ((\sigma^2)' * \rho(\sigma))'(\sigma^2)' * \rho(\sigma) \sigma'$   
 $= ((\rho(\sigma))' \rho(\sigma) \sigma')$   
 $= (\rho(\sigma))(\rho(\sigma))' \sigma'$ ,  
 since  $\supseteq (\mathbf{ES} \rho)$  and  $\rho(\sigma)(\rho(\sigma))' (\mathbf{ES} \rho)(\rho(\sigma))' \rho(\sigma)$ .  $\square$

**4.5 Corollary.** Let  $\sigma \in (\sigma) \supseteq (\sigma) \in \sigma_r(\rho)$ . Let  $\mathcal{U} \subseteq \mathbf{ES}$  be an e-variety and  $\mathcal{V} \subseteq \mathbf{I}$  be a variety such that  $\mathcal{U} \leftrightarrow \mathcal{V} \leftrightarrow \sigma$ . Then

$$\square \supseteq (\mathcal{U} \square \mathcal{V} \sigma)$$

**Proof.** Let  $\sigma \in \sigma_r(\sigma) (\mathcal{U} \square \mathcal{V} \sigma)$ . We will show that  $(\square \sigma)$ , i.e. and  $\rho(\sigma) \rho(\sigma)$ . Note that  $\sigma_r(\rho) \in \mathcal{U} (\sigma) \in \mathcal{V}$  (see 4.1 and 4.2). We use the homomorphism  $\sigma : (\sigma) \rightarrow (\text{End}(\sigma_r(\rho)) \circ)$  from 4.3 and the matched mapping  $\sigma : \sigma \rightarrow \sigma_r(\rho) \times_{\varphi} (\sigma)$  from 4.4. Now,  $\sigma_r(\rho) \times_{\varphi} (\sigma) \in \mathcal{U} \square \mathcal{V}$ . Thus the biidentity  $\sigma =$  is satisfied in  $\sigma_r(\rho) \times_{\varphi} (\sigma)$ , and therefore  $(\sigma) = (\sigma)$  (where  $\sigma$  is the extension of the matched mapping  $\sigma$ ). Hence, by 4.4(ii),  $(\rho(\sigma) \sigma') = (\rho(\sigma) \sigma')$ .  $\square$

**4.6 Lemma.** Let  $\sigma \in \mathbf{ES} \in \mathbf{I} : (\sigma) \rightarrow (\text{End}(\sigma) \circ)$  be a homomorphism,  $\sigma : \sigma \rightarrow \times_{\varphi} \sigma$  be a matched mapping such that

$$\begin{aligned} (\sigma_i) &= (\sigma_i \sigma_i) \\ (\sigma'_i) &= (\sigma_i \sigma_i) \quad (\text{for } i = 1, 2, \dots). \end{aligned}$$

Let  $\sigma \in (\sigma) \supseteq (\sigma)$ . Suppose that all identities from  $\sigma$  are satisfied in  $\sigma$ . Let  $\sigma_2 : \sigma \rightarrow \sigma$  be given by

$$i \mapsto \sigma_i$$

Let  $\sigma_1 : \sigma \rightarrow \sigma$  be given by

$$\begin{aligned} (\sigma_i) &\mapsto (\sigma_2(\sigma))(\sigma_i) \\ (\sigma'_i) &\mapsto (\sigma_2(\sigma_i))(\sigma_i) \quad (\text{for } (\sigma_i) \in \rho), \end{aligned}$$

where  $\sigma_2 : (\sigma) \rightarrow \sigma$  is the unary homomorphism extending  $\sigma_2$ . Finally, let  $\sigma : \sigma_r(\sigma) \rightarrow \times_{\varphi} \sigma$  be the extension of the matched mapping  $\sigma$ . Then

(i) the mapping  $\sigma_1$  is matched



$$\begin{aligned}
 &= (1(\rho(\ )) \ 2(\ ))(1(\rho(\ )) \ 2(\ )) \\
 &= (2(\ ' \ '))(1(\rho(\ ))) (2(\ ))(1(\rho(\ ))) \ 2(\ ) \ 2(\ )).
 \end{aligned}$$

It was proved in section (ii) that  $(2(\ ' \ '))(1(\rho(\ ))) = 1(\ ' \ ' * \rho(\ ))$ ,  $(2(\ ))(1(\rho(\ ))) = 1(* \rho(\ ))$ .

Now,

$$\begin{aligned}
 (\ ) &= (1((\ ' \ ' * \rho(\ ))(* \rho(\ ))) \ 2(\ )) \\
 &= (1(\rho(\ )) \ 2(\ )).
 \end{aligned}$$

Finally, let  $\in r(\ ) \ (\ ) = 1(\ ) \ (\ ) = (1(\rho(\ )) \ 2(\ ))$ . We have to show that

$$\begin{aligned}
 &(1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ )) = (1(\rho(\ )) \ 2(\ )), \\
 &(1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ ')) = (1(\rho(\ ')) \ 2(\ ')), \\
 &(1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ ')) = (1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ )).
 \end{aligned}$$

$$\begin{aligned}
 &\text{We see that } (1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ )) = \\
 &= ((2(\ ' \ '))(1(\rho(\ ))) (2(\ ))(1(\rho(\ '))) \ 2(\ ')) \\
 &\quad (1(\rho(\ )) \ 2(\ )) \\
 &= (1(\rho(\ )) \ 1((\rho(\ ))') \ 2(\ '))(1(\rho(\ )) \ 2(\ )) \\
 &= ((2(\ ' \ ' \ ' \ '))(1(\rho(\ )) \ 1((\rho(\ ))')) (2(\ '))(1(\rho(\ ))) \\
 &\quad 2(\ ' \ ')) \\
 &= (1(\rho(\ )) \ 1((\rho(\ ))') \ 1(\rho(\ )) \ 2(\ ' \ ')) \\
 &= (1(\rho(\ )) \ 2(\ )).
 \end{aligned}$$

$$\begin{aligned}
 &\text{Similarly, } (1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ ')) = \\
 &= (1(\rho(\ ')) \ 2(\ ')).
 \end{aligned}$$

$$\begin{aligned}
 &\text{Further, } (1(\rho(\ )) \ 2(\ ))(1(\rho(\ ')) \ 2(\ ')) = \\
 &= (1(\rho(\ )) \ 1((\rho(\ ))') \ 2(\ ')), \\
 &\quad (1(\rho(\ ')) \ 2(\ '))(1(\rho(\ )) \ 2(\ )) = \\
 &= ((2(\ ' \ ' \ ' \ '))(1(\rho(\ '))) (2(\ '))(1(\rho(\ ))) \ 2(\ ' \ ')) \\
 &= (1(\ ' * (\rho(\ ))') \ 1(\ ' * \rho(\ )) \ 2(\ ' \ ')) \\
 &= (1(\ ' * (\rho(\ ))') \ 1(\ ' * \rho(\ )) \ 2(\ ' \ ')) \\
 &= (1((\rho(\ ))') \ 1(\rho(\ )) \ 2(\ ' \ ')) \\
 &= (1(\rho(\ )) \ 1((\rho(\ ))') \ 2(\ ' \ ')).
 \end{aligned}$$

We used the following facts:

$$2(\ ) \ (\text{see 3.1}), \quad 2(\ ' \ 2(\ )) \in (\ ) \ 2(\ ' \ ) \in (\ ). \quad \square$$

**4.7 Corollary.** Let  $\in (\ ) \supseteq (\ ) \in r(\rho)$ . Let  $\in \mathbf{ES} \in \mathbf{I}$ . Suppose that all identities from are satisfied in and all bidentities from are satisfied in . Finally, let  $:(\cdot) \rightarrow (\text{End}(\ ) \circ)$  be a homomorphism. Then  $\square \subseteq (\{ \times_{\varphi} \})$ .

**Proof.** Let  $\in r(\ ) \ (\square) : \overline{\ } \rightarrow \times_{\varphi}$  be a matched mapping. We have to show that  $(\ ) = (\ )$ , where  $:(\cdot) \rightarrow \times_{\varphi}$  is the extension of . We know that  $\rho(\ ) = \rho(\ )$ . Consider the mappings  $_1$  and  $_2$  from 4.6. The mapping  $_1$  is matched by 4.6(i). Let  $_1 : r(\rho) \rightarrow$  be the extension of  $_1$  and  $_2 : (\ ) \rightarrow$  be the unary homomorphism extending  $_2$ . Then  $_1(\rho(\ )) = _1(\rho(\ ))$  and  $_2(\ ) = _2(\ )$ . Thus  $(\ ) = (\ )$  (by 4.6(iii)).  $\square$



**4.8 Result.** ([2], Lemma 1) Let  $\alpha : S \rightarrow T$  be a surjective homomorphism of regular semigroups and let  $\beta \in \text{Con } S = \text{Con } T$ . Then there exist  $\gamma \in \text{Con } S$  such that  $\beta = \gamma \circ \alpha$  and  $\alpha(\gamma) = \alpha(\beta) = \beta$ .

**4.9 Corollary.** Let  $\beta \in \text{Con } (U \square V) \supseteq \alpha \in \text{Con } r(\rho)$ . Let  $\mathcal{U} \subseteq \mathbf{ES}$  be an e-variety,  $\mathcal{V} \subseteq \mathbf{I}$  be a variety such that  $\mathcal{U} \leftrightarrow \mathcal{V} \leftrightarrow \cdot$ . Then

$$\square \subseteq (\mathcal{U} \square \mathcal{V} \quad \cdot)$$

**Proof.** Let  $\beta \in r(\alpha) \quad (\square)$   $\in \mathcal{U} \square \mathcal{V}$  and let  $\gamma : \bar{\quad} \rightarrow \quad$  be a matched mapping. We will show that  $\beta = \gamma$ , where  $\gamma : r(\alpha) \rightarrow \quad$  is the extension of  $\alpha$ . It follows from 2.7 that there exist  $\beta \in \mathcal{U} \in \mathcal{V}$ , a homomorphism  $\delta : (\cdot) \rightarrow (\text{End}(\cdot) \circ)$ , a regular subsemigroup  $\beta$  in  $\times_{\varphi}$  and a surjective homomorphism  $\delta : \rightarrow \quad$ . By 4.8, there is a matched mapping  $\bar{\delta} : \bar{\quad} \rightarrow \quad$  such that  $\bar{\delta}(\beta) = \beta$  for all  $\beta \in \bar{\quad}$ . Then  $\bar{\delta}(\beta) = \beta$  for all  $\beta \in r(\alpha)$  ( $\bar{\delta} : r(\alpha) \rightarrow \quad$  is the extension of  $\bar{\delta}$ ). Now, we use 4.7. We have  $\bar{\delta}(\beta) = \beta$ . Thus  $\bar{\delta}(\beta) = \bar{\delta}(\beta) \quad \beta = \beta$ . □

**4.10 Theorem.** Let  $\beta \in \text{Con } (U \square V) \supseteq \alpha \in \text{Con } r(\rho)$ . Let  $\mathcal{U} \subseteq \mathbf{ES}$  be an e-variety,  $\mathcal{V} \subseteq \mathbf{I}$  be a variety such that  $\mathcal{U} \leftrightarrow \mathcal{V} \leftrightarrow \cdot$ . Then

- (i)  $\square \in \text{Con } r(\alpha)$
- (ii)  $\mathcal{U} \square \mathcal{V} \leftrightarrow \square$
- (iii) The mapping  $\gamma : r(\alpha) \square \rightarrow r(\rho) \times_{\varphi} (\cdot)$  defined by

$$\gamma((\beta \square \gamma)) = (\rho(\beta) \quad \gamma)$$

where  $\gamma$  is the homomorphism from 4.3, is an embedding.

**Proof.**

(i) and (ii) Note that  $\mathcal{U} \square \mathcal{V} \subseteq \mathbf{ES}$  (see 2.4). By 4.5 and 4.9 we have  $\square = (\mathcal{U} \square \mathcal{V} \quad \cdot)$ . Thus  $\square \in \text{Con } r(\alpha)$  and  $\mathcal{U} \square \mathcal{V} \leftrightarrow \square$  by 4.1.

(iii) It follows immediately from the definition of  $\square$  that  $\gamma$  is a correctly defined injective mapping.

$\gamma$  is a homomorphism:

Let  $\beta \in r(\alpha)$ . Then

$$\begin{aligned} ((\beta \square \gamma))((\beta \square \gamma)) &= (\beta \square \gamma) \\ &= (\rho(\beta) \quad \gamma) = ((\beta \quad \gamma) * \rho(\beta))(\beta \quad \gamma) \\ &= ((\beta \quad \gamma) * \rho(\beta))(\beta \quad \gamma) \\ &= (\rho(\beta) \quad \gamma)(\beta \quad \gamma) \\ &= ((\beta \square \gamma))((\beta \square \gamma)). \end{aligned} \quad \square$$

**4.11 Remark.** Theorem 4.10 together with Result 4.2 show that bifree objects in  $\mathcal{U} \square \mathcal{V}$  are isomorphic to some subsemigroups in suitable semidirect products of bifree objects in  $\mathcal{U}$  by free objects in  $\mathcal{V}$ , for any e-variety  $\mathcal{U} \subseteq \mathbf{ES}$  and any variety  $\mathcal{V} \subseteq \mathbf{I}$ .

This section is concluded with a corollary of Theorem 4.10. First, the following result ensures that if  $\mathcal{U}$  and  $\mathcal{V}$  are varieties of inverse semigroups then  $\mathcal{U} \square \mathcal{V}$  is also a variety of inverse semigroups.

**4.12 Result.** ([1], Proposition 1) *Let  $(S, \cdot)$  be inverse semigroups,  $\theta : (S, \cdot) \rightarrow (\text{End}(S), \circ)$  be a homomorphism. Then  $(S, \cdot) \times_{\varphi} (S, \cdot)$  is also an inverse semigroup.*

**4.13 Corollary.** *Let  $(S, \cdot) \supseteq (T, \cdot) \in \mathcal{E}(\rho) \supseteq (T, \rho)$ . Let  $\mathcal{U}, \mathcal{V} \subseteq \mathbf{I}$  be varieties such that  $\mathcal{U} \leftrightarrow (S, \cdot) \leftrightarrow \mathcal{V} \leftrightarrow (T, \cdot)$ . Denote by  $\square$  the fully invariant congruence on  $(S, \cdot)$  corresponding to the variety  $\mathcal{U} \square \mathcal{V}$ . Then*

$$((S, \cdot) \square) \iff (S, \cdot) \text{ and } \rho((S, \cdot) \square) = \rho(S, \cdot) \text{ (for all } (S, \cdot) \in \mathcal{E}(\rho)).$$

**Proof.** Let  $\theta$  be the biinvariant congruence on  $(S, \cdot)$  corresponding to the e-variety  $\mathcal{U}$ . It follows from 4.10(ii) that  $\theta \square$  is the biinvariant congruence on  $(S, \cdot)$  corresponding to the e-variety  $\mathcal{U} \square \mathcal{V}$ . Clearly,  $\theta \square = \theta \cap ((S, \cdot) \times_{\varphi} (S, \cdot)) \square \theta \square = ((S, \cdot) \square) \cap ((S, \cdot) \times_{\varphi} (S, \cdot))$ . Let  $(s, t) \in \theta \square$ . There are  $(s_0, t_0) \in (S, \cdot)$  such that  $(s, t) \theta (s_0, t_0)$ . Then  $(s_0, t_0) \theta (s_0, t_0) \square (s_0, t_0)$ . Further,  $\rho((s_0, t_0) \theta) = \rho(s_0, t_0) = \rho((s_0, t_0) \square)$  (by 3.7). Thus  $\rho((s, t) \theta \square) = \rho(s, t) = \rho((s, t) \square)$ . Now,

$$\begin{aligned} ((S, \cdot) \square) &\iff \theta \square (S, \cdot) \square \theta \\ &\iff \theta \square (\theta \square) \square \theta \\ &\iff \theta \square \theta \text{ and } \rho(\theta \square) \square \rho(\theta) \\ &\iff \theta \square \text{ and } \rho(\theta \square) = \rho(\theta) \\ &\iff \theta \square \text{ and } \rho(S, \cdot) = \rho(S, \cdot). \end{aligned}$$

We used also the facts that  $\rho((s_0, t_0) \theta) = \rho(s_0, t_0) \in \mathcal{E}(\rho)$  (by 3.9). □

### 5. ASSOCIATIVITY

We specify our notation in this section. Let  $S$  be a countable set,  $\theta \in \mathcal{E}(\rho)$ ,  $\theta \supseteq (S, \cdot)$ . Put

$$\rho = (S, \cdot) \times_{\varphi} (S, \cdot)$$

and define

$$\theta : (S, \cdot) \rightarrow (S, \rho)$$

in the same way as the mapping  $\rho$  in the section 3 (of course, we replace the set  $S = \{s_1, s_2, \dots\}$  by an arbitrary countable set  $S$ ).

Throughout this section, let  $\mathcal{U} \subseteq \mathbf{ES}$  be an e-variety and  $\mathcal{V} \subseteq \mathbf{I}$  be varieties. We will prove syntactically that

$$((\mathcal{U} \square \mathcal{V}) \square \mathcal{W}) = \mathcal{U} \square (\mathcal{V} \square \mathcal{W})$$

Note that  $\mathcal{V} \square \mathcal{W}$  is a variety of inverse semigroups by 4.12 and so the right side of the equation mentioned above is meaningful.

Let

$$\begin{aligned} \mathcal{U} &\in \mathcal{E}(\rho) \supseteq (S, \cdot) \leftrightarrow \mathcal{W}, \\ \mathcal{V} &\in \mathcal{E}(\rho) \supseteq (S, \rho) \leftrightarrow \mathcal{V}, \end{aligned}$$

$$\begin{aligned} \in & \quad r(\sigma \square \rho) \leftrightarrow \mathcal{U}, \\ ' \in & \quad r((\rho)\sigma) \quad ' \leftrightarrow \mathcal{U}. \end{aligned}$$

In view of 4.10 we have to prove that

$$(\ ' \square ) \square = \square ( \square )$$

Choose  $\in \in r(\ )$ . Then (by the definition of  $\square$  and by 4.13)

$$\begin{aligned} (\ \square ( \square ) ) & \iff ( \square ) \\ & \iff ( \square ) ( ) \quad ( \square ) ( ) \\ & \iff ( ) ( ) \quad ( ) ( ) \\ & \iff ( \square ) ( ) \quad ( \square ) ( ) \end{aligned}$$

and

$$\begin{aligned} (( \ ' \square ) \square ) & \iff ( ) ( ) ( \ ' \square ) \quad ( ) ( ) \\ & \iff ( ) ( ) \quad ( ) ( ) \\ & \iff ( \rho ) ( ( ) ( ) ) \quad ' \quad ( \rho ) ( ( ) ( ) ) \end{aligned}$$

Clearly, it suffices to prove that

$$( \square ) ( ) \quad ( \square ) ( )$$

is equivalent to

$$( \rho ) ( ( ) ( ) ) \quad ' \quad ( \rho ) ( ( ) ( ) )$$

Define  $\sigma \square \rho \rightarrow (\rho)\sigma$  by

$$( ( \square ) ) \mapsto ( ( ) ( ) ( ) )$$

$$( \in ( ) \in ).$$

**5.1 Lemma.** *is a correctly defined injective mapping.*

**Proof.**

1. Let  $\in ( ) ( \square )$ . We want to show:  $( ) ( ) ( ) ( )$ .  
It follows immediately from 4.13.
2. Let  $\in ( ) \in ( ( ) ( ) ( ) ) = ( ( ) ( ) ( ) )$ . We want to show that  $( ( \square ) ) = ( ( \square ) )$ .  
 $( ) = ( )$  implies  $=$ .  
 $( ) ( ) = ( ) ( )$  together with  $\implies ( \square )$  (see 4.13)  $\square$

Now, we extend the mapping

$$\sigma \square \rho \rightarrow (\rho)\sigma$$

to the unary homomorphism

$$: ( \sigma \square \rho ) \rightarrow ((\rho)\sigma)$$

**5.2 Lemma.**  $( * ( \square ) ( ) ) = ( ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) )$   
 for any  $\in ( )$ .

**Proof.** By induction with respect to :

1. Let  $\in$  . Then  $( * ( \square ) ( ) ) = ( * ( ' ( \square ) ) )$   
 $= ( ( ' ( \square ) ) ) = ( ( ) ( ' ) ( ' ) )$   
 $= ( ( ' ' ' * ( ) ( ) ) ( * ( ) ( ' ) ) ( ' ) )$   
 $= ( ' ' * ( ) ( ) ) * ( ( * ( ) ( ' ) ) ( ' ) )$ .

Now,  $* ( ) ( ' ) =$   
 $= * ( ' ' * ( ' ) ) ( * ( ' ) )$   
 $= ( ' ) ( ' )$ .

So,  $( * ( \square ) ( ) ) =$   
 $= ( ' ' * ( ) ( ) ) * ( ( ' ) ( ' ) ( ' ) )$   
 $= ( ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) )$ .

2. Let  $\in ( )$ . We suppose that

$$( * ( \square ) ( ) ) = ( ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) ) ,$$

$$( * ( \square ) ( ) ) = ( ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) )$$

for all  $\in ( )$ .

Now, choose an arbitrary  $\in ( )$ . Then  $( * ( \square ) ( ) ) =$   
 $= ( ( ' ' * ( \square ) ( ) ) ( * ( \square ) ( ) ) )$   
 $= ( ( ' ' ' ' * ( ) ( ' ) ) *$   
 $* ( \rho ) ( ' ' * ( ) ( ) )$   
 $(( ' ' ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) ) )$ .

Now,  $' ' ' ' * ( ) ( ' ) =$   
 $= ' ( ' ' ' ) * ( ) ( ' )$   
 $= ( ) ( ' )$   
 $= ( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ' ) )$   
 $= ( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ' ) )$   
 $= ( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ' ) )$

and  $' ' ' ' * ( ) ( ) =$   
 $= ( ' ' ' ' * ( ) ( ) ) ( ' ' ' ' * ( ) ( ) )$   
 $= ( ' ' ' ' * ( ) ( ) ) ( ' ' ' ' * ( ) ( ) )$   
 $= ( ' ' ' ' * ( ) ( ) ) ( ' ' ' ' * ( ) ( ) )$ .

So,  $( * ( \square ) ( ) ) = ( ' ' ' ' * ( ) ( ) ) *$   
 $* ( ( * ( ) ( ' ) ) * ( \rho ) ( ' ' * ( ) ( ) ) )$   
 $(( ' ' ' ' * ( ) ( ) ) * ( \rho ) ( * ( ) ( ) ) )$ .

We will show that  $( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ) )$   
 $( * ( ) ( ' ) ) ( ' ' ' ' * ( ) ( ) ) ( \rho )$   
 $( \rho ) * ( ) ( ' ) : ( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ) )$   
 $( * ( ) ( ' ) ) ( ' ' ' ' * ( ) ( ) ) ( \rho )$   
 $( \rho ) ( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ) )$   
 $(( ' ' ' ' * ( ) ( ) ) ( * ( ) ( ) ) )$   
 $= * ( ) ( ) ( ( ) ( ) )$   
 $= * ( ( ) ' * ( ) ( ) ) ( * ( ( ) ' * ( ) ( ) ) )$   
 $= * ( ( ) ' ( ) ' * ( ) ( ) ) ( * ( ) ( ) )$



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