

Václav J. Havel

3-configurations with simple edge basis and their corresponding quasigroup identities

Archivum Mathematicum, Vol. 29 (1993), No. 3-4, 187--195

Persistent URL: <http://dml.cz/dmlcz/107482>

Terms of use:

© Masaryk University, 1993

Institute of Mathematics of the Academy of Sciences of the Czech Republic provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these *Terms of use*.



This paper has been digitized, optimized for electronic delivery and stamped with digital signature within the project *DML-CZ: The Czech Digital Mathematics Library* <http://project.dml.cz>

3-CONFIGURATIONS WITH SIMPLE EDGE BASIS AND THEIR CORRESPONDING QUASIGROUP IDENTITIES

V. J. HAVEL

ABSTRACT. There is described a procedure which determines the quasigroup identity corresponding to a given 3-coloured 3-configuration with a simple edge basis.

PART 1 3-CONFIGURATIONS

1.1. MAIN NOTATIONS

Under a *(3)-configuration* we shall understand a finite incidence structure $(\mathcal{V}, \mathcal{E}, \mathbb{I})$ such that $|\{y \in \mathcal{E} \mid x \mathbb{I} y\}| \leq 3$ for all $x \in \mathcal{V}$. If, moreover, $|\{y \in \mathcal{E} \mid x \mathbb{I} y\}| = 3$ for all $x \in \mathcal{V}$ and $|\{x \in \mathcal{V} \mid x \mathbb{I} y\}| \geq 2$ for all $y \in \mathcal{E}$, then we get a *3-configuration*. If one (3)-configuration is an incidence substructure of another (3)-configuration we speak of a *subconfiguration*. Elements of \mathcal{V} will be called *vertices* and elements of \mathcal{E} *edges* of a given (3)-configuration $(\mathcal{V}, \mathcal{E}, \mathbb{I})$. After all, such (3)-configuration are hypergraphs with vertices of degree ≤ 3 and with the property that two distinct vertices cannot lie simultaneously on two distinct edges.

If $x \mathbb{I} y$ then x, y are said to be *adjacent*. Two distinct vertices adjacent to the same edge; respectively two or three mutually distinct edges adjacent to the same vertex, are said to be *neighbouring*. A (3)-configuration is called *connected* if for any distinct vertices a, b there is a finite sequence of edges e_1, e_2, \dots, e_k such that $a \mathbb{I} e_1, b \mathbb{I} e_k$ and that e_i, e_{i+1} are neighbouring for all $i \in \{1, \dots, k-1\}$. Throughout the paper every 3-configuration will be supposed to be connected.

Let there be given a 3-configuration $\mathfrak{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$. We shall define a halfoperation on \mathcal{E} : for every couple (x, y) of neighbouring edges x, y their product $x \cdot y$ is equal to the remaining edge which forms together with x and y a neighbouring triple. We obtain the *edge halfgroupoid* (\mathcal{E}, \cdot) associated to \mathfrak{C} .

1991 *Mathematics Subject Classification*: Primary 05B30, Secondary 20N05.

Key words and phrases: 3-configuration, pillar (over an edge set), 3-web, closure condition, quasigroup identity.

Received July 1, 1992.

1.2. BASES OF 3-CONFIGURATIONS

Let $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ be a 3-configuration. If \mathcal{X} is a non-void set of edges then denote by $[\mathcal{X}]$ the set of such vertices which are adjacent to least two edges from \mathcal{X} .

Now start with a non-void set \mathcal{A} of edges, construct a sequence $\mathcal{A}_0 := \mathcal{A}$, $\mathcal{A}_1 := [\mathcal{A}_0]$, $\mathcal{A}_2 := [\mathcal{A}_1]$, $\mathcal{A}_3 := [\mathcal{A}_2]$, ... and form the unions $\mathcal{E}_{\mathcal{A}} := \bigcup_i^{\infty} \mathcal{A}_i$, $\mathcal{V}_{\mathcal{A}} := \bigcup_i^{\infty} \mathcal{A}_i$. Then $(\mathcal{V}_{\mathcal{A}}, \mathcal{E}_{\mathcal{A}}, \mathbb{I})$ is a subconfiguration in \mathcal{C} generated by \mathcal{A} ; we

denote it by $\langle \mathcal{A} \rangle$. We shall call the sequence $(\mathcal{A}_i)_i^{\infty}$ the *pillar* over \mathcal{A} . Further let the *stage* of elements of \mathcal{A}_i , respectively \mathcal{A} be defined as 0, respectively 1. If the stages of elements of $\mathcal{A}_0 \cup \mathcal{A}_1 \cup \dots \cup \mathcal{A}_i$ are already known then define the stage of elements of $\mathcal{A}_{i+1} \setminus \mathcal{A}_i$ to be equal to $i+2$. The edge, respectively vertex a of $\langle \mathcal{A} \rangle$ is said to be *terminal* if there is no edge, respectively no vertex of greater stage neighbouring to a . Under a *constituent* we shall understand a term expression (by means of \cdot) of edges over \mathcal{A} within the framework of the pillar. The set \mathcal{A} will be called *independent* if $\langle \mathcal{A} \setminus \{a\} \rangle \neq \langle \mathcal{A} \rangle$ for all $a \in \mathcal{A}$. The set \mathcal{A} is called an *edge basis* of \mathcal{C} if it is independent and if $\langle \mathcal{A} \rangle = \mathcal{C}$. The set of all element bases of a given 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ cannot be void: we can namely delete, step by step, all dependent elements from the full edge set \mathcal{E} . If \mathcal{B} is an edge basis and \mathcal{T} the set of all terminal elements (vertices, as well as edges, if they exist) then

$\mathcal{V} + \mathcal{B} = \mathcal{E} + \mathcal{T}$ (a well-known condition, cf. [1], pp. 50-51). Different edge bases of \mathcal{C} can we have different number of edges. Bases which lead to just one terminal vertex or edge will be denoted as *simple*. Thus for simple bases we have

$$\mathcal{T} = \mathcal{V} + \mathcal{B} - \mathcal{E} = 1, \text{ i.e. } \mathcal{B} = \mathcal{E} - \mathcal{V} + 1.$$

If \mathcal{B} is a simple basis of \mathcal{C} then in the corresponding sequence $(\mathcal{B}_i)_i^{\infty}$ the terminal vertex t lies on two edges e, e' with maximal stage (equal to stage of t minus one) and still on the third edge e'' . This gives the equation $e \cdot e' = e''$ and consequently the equation between corresponding terms over \mathcal{B} in the halfgroupoid (\mathcal{B}, \cdot) . Analogously, if there is the terminal edge t then it has maximal stage and determines an equality between two corresponding terms over \mathcal{B} in (\mathcal{E}, \cdot) .

We give two examples: First let us investigate the 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ on Fig. 1 with $\mathcal{V} = 10$, $\mathcal{E} = 12$, $\mathcal{B} = \{x, y, z\}$.

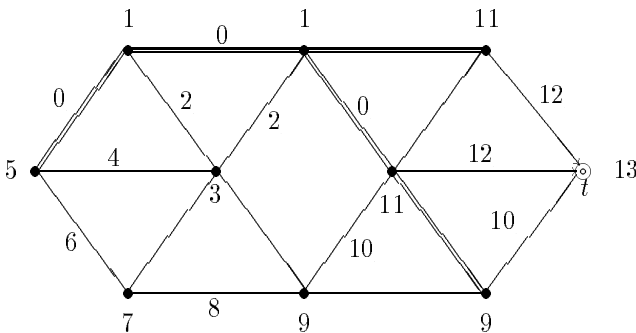
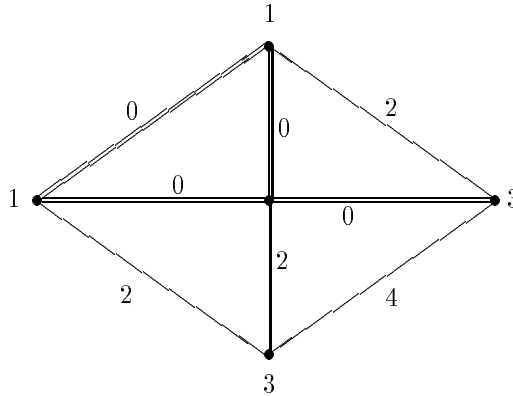


Fig.1

Here: $0 \dots x; 0 \dots y; 0 \dots z; 2 \dots xy; 2 \dots yz; 4 \dots (xy)(yz); 6 \dots x((xy)(yz));$
 $8 \dots (x((xy)(yz)))(yz); 10 \dots ((x((xy)(yz)))(yz))(xy); 10 \dots ((x((xy)(yz)))(yz))z;$
 $12 \dots (((x((xy)(yz)))(yz))(xy)); 12 \dots (((x((xy)(yz)))(yz))(xy))y.$

This basis is simple and determines just one terminal vertex t of stage 13, which results as the vertex adjacent to two neighbouring edges of stages 12. The remaining edge through t is one of edges of stage 10. The corresponding term equality is thus $((((x((xy)(yz)))(yz))(xy))y) \cdot (((x((xy)(yz)))(yz))xy)z = ((x((xy)(yz)))(yz))z.$

The second example deals with a 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ on Fig. 2 with $\mathcal{V} = 5, \mathcal{E} = 7, \mathcal{B} = \{x, y, z\}$. There is just one terminal element t . It has stage 4 and determines the equality among terms $(xy)(yz), (xy)z.$



$0 \dots x, 0 \dots y, 0 \dots z, 2 \dots xy, 2 \dots yz, 2 \dots xz, 4 \dots (xy)z \dots (xy)(yz)$

Fig. 2

Now let \mathcal{B} be a simple edge basis of a given 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ and $t = t$ the term equality over \mathcal{B} in (\mathcal{E}, \cdot) corresponding to the terminal element (vertex or edge).

We shall find some properties of the equality $t = t$:

- (i) The set \mathcal{B} must contain at least three edges:
 In fact, by the existence of just one edge in \mathcal{B} or just two edges in \mathcal{B} , respectively, the pillar over \mathcal{B} should have no vertices of stage 1 respectively 3 contrary to further size of the pillar.
- (ii) Every $x \in \mathcal{B}$ occurs at least twice in $t = t$:
 By exactly one occurrence (in the subterm $u \cdot x$ of the pillar or in $t = x$, respectively $t = x$) we should have on x only one vertex or we should express $x \in \mathcal{B}$ by the edges from \mathcal{B} (which contradicts the independence of edges of the basis).
- (iii) No constituent in $(\mathcal{B}_i)_i^\infty$ can be of the form $u(uv)$ with u and v as constituents in $(\mathcal{B}_i)_i^\infty$:

This follows at once from the definition of the pillar over \mathcal{B} .

- (iv) Equations $t = uv$, $t = uv$, respectively $t = u$, $t = uv$, respectively $t = uv$, $t = u$ with constituents in $(\mathcal{B}_i)_i^\infty$ are not possible:

The equality $uv = uv$ or $u = uv$, respectively, should violate the fundamental properties of the 3-configuration ($uv = uv \Rightarrow v = v$, but on the other side, v and v should be distinct edges as constituents in the pillar; moreover, $u = uv$ contradicts the fact that uv must be the third edge neighbour to u and v).

- (v) Let t and t be not formed by exactly one edge of \mathcal{B} and let $x \in \mathcal{B}$ enters in $v = ux$, where u, v are constituents of the pillar. Then x enters still in some term $v' = u'x \neq v$ where also u', v' are constituents.

We can define a closure condition in a 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ derived from a 3-configuration $\mathcal{C}_* = (\mathcal{V}_*, \mathcal{E}_*, \mathbb{I}_*)$ having a simple basis \mathcal{B} with terminal element t and with the *main subconfiguration* $\hat{\mathcal{C}}_* = (\mathcal{V}_*, \mathcal{E}_* \setminus \{t\}, \mathbb{I}_*)$. The corresponding *closure condition* claims:

For every incidence structure homomorphism $\sigma : \hat{\mathcal{C}}_* \rightarrow \mathcal{C}$ there exists its prolongation onto a homomorphism $\bar{\sigma} : \mathcal{C}_* \rightarrow \mathcal{C}$.

In the second part of our considerations we shall adapt this definition for 3-web configurations and 3-webs. As we are familiar with the subject, only less is known about closure conditions in general 3-configurations (or, more largely, in webs of arbitrary number of parallel line pencils).

PART 2 CLOSURE CONDITIONS IN 3-WEBS

2.1. COLOURED CONFIGURATIONS

We say that a (3)-configuration is 3-coloured if for every vertex x an injective map of the set of all edges through x into the colour set $\{1, 2, 3\}$ is given such that for all neighbouring vertices x, y the colour assigned to the common adjacent edge is the same. The edges with colour 1, 2, or 3, respectively will be called *vertical*, *horizontal* or *skew*, respectively.

A 3-coloured (3)-configuration \mathcal{C} is a *subconfiguration* of a coloured (3)-configuration \mathcal{C} if it is an incidence substructure preserving the colour of every edge from \mathcal{C} .

Now we shall define a 3-web independently on 3-configurations though it is possible to take a (finite) 3-web as a special 3-configuration.

A *3-web* is defined as an incidence structure $(\mathcal{P}, \mathcal{L}, \mathbb{I})$, $\mathcal{P} \leq 2$, together with a decomposition of the set \mathcal{L} onto mutually disjoint sets $\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3$ such that for every $x \in \mathcal{P}$ and every $i \in \{1, 2, 3\}$ there is just one $y \in \mathcal{L}_i$ such that $x \mathbb{I} y$ and for every $x \in \mathcal{L}_i$ and every $y \in \mathcal{L}_j$ for distinct $i, j \in \{1, 2, 3\}$ there is just one $z \in \mathcal{P}$ such that $z \mathbb{I} x$ and $z \mathbb{I} y$. Elements of \mathcal{P} are *points*, elements of \mathcal{L} *lines*, $\mathcal{L}_1, \mathcal{L}_2$ and \mathcal{L}_3 are *pencils* (of *parallel lines*). Lines from $\mathcal{L}_1, \mathcal{L}_2$ or \mathcal{L}_3 , respectively, are said to be *vertical*, *horizontal* or *skew*, respectively. If $z \mathbb{I} x$ and $z \mathbb{I} y$ is true for distinct lines x, y then z is called the *intersection point* of both lines.

A *homomorphism* of a 3-coloured (3)-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ into a 3-web $\mathcal{W} = (\mathcal{P}, \mathcal{L}, \mathbb{I})$ is an incidence structure homomorphism which maps every

vertical, horizontal or skew edge, respectively, onto a vertical, horizontal or skew line, respectively.

On Fig. 3 we present a non-3-colourable 3-configuration $(\mathcal{V}, \mathcal{E}, \mathbb{I})$,

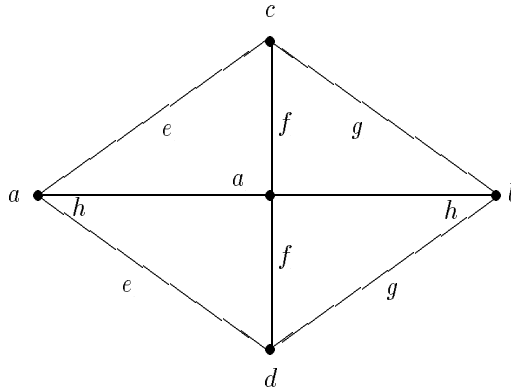


Fig. 3

with $\mathcal{V} = \{a, b, c, d\}$, $\mathcal{E} = \{e, f, g, h\}$ having one edge adjacent to just three vertices. If the colours of e, f, g, h are $i, j, k, l \in \{1, 2, 3\}$ then $i \neq j \neq k \neq l$ so that every possibility for l leads to a contradiction.

Let $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ be a 3-coloured 3-configuration and $\mathcal{E}_h, \mathcal{E}_v, \mathcal{E}_c$ sets of all its horizontal, vertical or complementary edges, respectively. Then define a half-operation $\binom{k}{ij} : \mathcal{E}_i \times \mathcal{E}_j \rightarrow \mathcal{E}_k$ such that $\binom{k}{ij}(x_i, x_j) = x_k$ is defined if and only if $x_i \in \mathcal{E}_i, x_j \in \mathcal{E}_j, x_k \in \mathcal{E}_k$ are neighbouring edges ($\binom{k}{ij}$ is an arbitrary permutation of the set $\{1, 2, 3\}$). So we obtain six *parastrophic* halfoperations associated to \mathcal{C} which from a *three-sorted edge halfquasigroup* $(\mathcal{E}_h, \mathcal{E}_v, \mathcal{E}_c, \{\binom{k}{ij} \mid \binom{k}{ij} \in \mathcal{S}\})$ associated to \mathcal{C} .

Let \mathcal{B} be a basis of a simple 3-coloured 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ with the terminal element t . We express the edges of \mathcal{C} as constituents in the pillar $(\mathcal{B}_i)_i^\infty$ over \mathcal{B} with use of halfoperations $\binom{k}{ij}$. The existence of t leads to an equality between constituents of greatest stage.

From the successive construction of constituents in $(\mathcal{B}_i)_i^\infty$ we obtain the validity of the *Position Property*: In every occurrence of the edge e as a constituent $\binom{k_1}{i_1 j_1}(\cdot, \cdot)$ or as a member of other constituents $\binom{k_2}{i_2 j_2}(e, \cdot), \binom{k_3}{i_3 j_3}(\cdot, e)$ the equalities $k = i = j$ must hold.

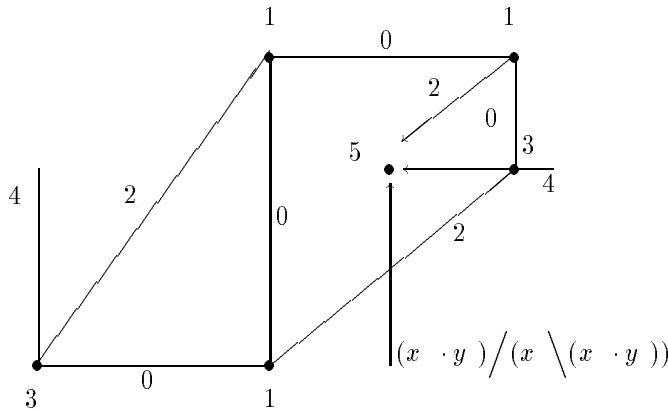
Now we see that Position Property implies that constituents of the form $\binom{k_1}{i_1 j_1}(u, \binom{k_2}{i_2 j_2}(u, v))$ with $i \neq j$ do not occur. Property (iii) affirms that constituents of the form $\binom{k_1}{i_1 j_1}(u, \binom{k_2}{i_2 j_2}(u, v))$ do not occur. Property (iv) excludes the equality $t = t$ with $t = \binom{k}{ij}(u, v), t = \binom{k}{ij}(u, v)$ whereas the case $t = \binom{k}{ij}(t, v)$ with $i \neq 3$ is excluded again by Position Property.

If Position Property holds (v) can be reduced onto the occurrence of a constituent $\binom{k}{ij}(x, u')$, $u' \neq u$ when a constituent $\binom{k}{ij}(x, u)$ occurs (with t and t' different from a single $x \in \mathcal{B}$).

When a 3-sorted quasigroup $\mathbb{Q} = (Q_1, Q_2, Q_3, \{ \binom{k}{ij} | (i, j, k) \in \mathcal{S} \})$ is given with $Q_1 = Q_2 = Q_3 > 1$ then let $t = t'$ be a quasigroup identity for \mathbb{Q} with a finite set \mathcal{X} , $|\mathcal{X}| \leq 3$ of variables (the formal construction of terms $t = t'$ is supposed as known, see e.g.[1], pp. 30-31). The subterms of $t = t'$ fulfill necessarily this *Position Property*: In every occurrence of a subterm $\tau = \binom{k_1}{i_1 j_1}(\cdot, \cdot)$ and subterms $\binom{k_2}{i_2 j_2}(\tau, \cdot)$, $\binom{k_3}{i_3 j_3}(\cdot, \tau)$ the equalities $k = i = j$ must hold.

We assume that following conditions are fulfilled: every $x \in \mathcal{X}$ occurs in $t = t'$ at least twice; the case $t = \binom{k}{ij}(u, v)$, $t' = \binom{k}{ij}(u, v')$ is excluded; if t, t' are not single variables and a subterm $\binom{k}{ij}(x, u)$ occurs for $x \in \mathcal{X}$ then also a further subterm $\binom{k}{ij}(x, u')$ with $u' \neq u$ occurs; subterms are reduced by fundamental quasigroup identities $\binom{j}{ik}(\xi, \binom{k}{ij}(\xi, \eta)) = \eta$ where $\binom{k}{ij}(\xi, \xi) = \xi \Leftrightarrow \binom{k}{ij}(\xi_i, \xi_j) = \xi_k$ for all permutations $\binom{k}{ij}$. Under these assumptions the subterm buildup of the identity coincides with the edge generation of a convenient 3-coloured 3-configuration over a simple basis corresponding to \mathcal{X} . We leave the details aside.

As an example we shall investigate a 3-coloured 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ on Fig. 4 with a simple basis $\mathcal{B} = \{x, x, y, y\}$, where x, x are vertical edges, y, y horizontal edges and the terminal element is the vertex b . The edges of stages 4 determine the remaining edge through b . This leads to the equality $((x \cdot y)/y) \cdot (x \setminus (x \cdot y)) = x \cdot y$ where $\binom{k}{ij} = \cdot$, $\binom{k}{ij} = /$, $\binom{k}{ij} = \setminus$. We can write equivalently $(x \cdot y)/y = (x \cdot y)/(x \setminus (x \cdot y))$.



0 ... x , 0 ... x , 0 ... y , 0 ... y , 1 ... a' , 2 ... $x \cdot y$, 2 ... $x \cdot y$,
 2 ... $x \cdot y$, 3 ... a , 3 ... a'' , 4 ... $(x \cdot y)/y$, 4 ... $x \setminus (x \cdot y)$, 5 ... b

Fig. 4

The stages of single edges and vertices are written as labels on Fig. 4 as well as

the corresponding constituents. We can readily confirm that properties (i)-(v) are fulfilled.

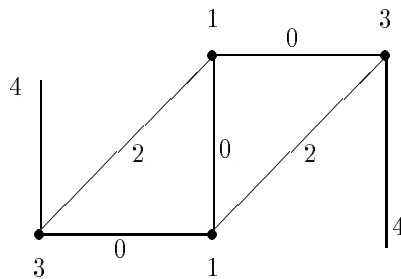
We say that a 3-web $\mathcal{W} = (\mathcal{P}, \mathcal{L}, \mathbb{I})$ satisfies the closure condition $\Gamma(\mathcal{C}, \mathcal{B})$ derived from a 3-coloured 3-configuration $\mathcal{C} = (\mathcal{V}, \mathcal{E}, \mathbb{I})$ with a simple basis \mathcal{B} if every homomorphism $\sigma : \hat{\mathcal{C}} \rightarrow \mathcal{W}$ can be prolonged onto a homomorphism $\bar{\sigma} : \bar{\mathcal{C}} \rightarrow \mathcal{W}$. Here $\hat{\mathcal{C}} = (\mathcal{V} \setminus \{t\}, \mathcal{E}, \mathbb{I})$ or $\bar{\mathcal{C}} = (\mathcal{V}, \mathcal{E} \setminus \{t\}, \mathbb{I})$ where t is the terminal vertex or edge.

To a given 3-web $\mathcal{W} = (\mathcal{P}, \mathcal{L}, \mathbb{I})$ with pencils $\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3$ we associate the *coordinatizing 3-sorted quasigroup* $(\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3, \{ \binom{k}{ij} | (\binom{k}{ij}, \binom{k}{ij}) \in \mathcal{S} \})$ which is defined by $\binom{k}{ij} : \mathcal{L}_i \times \mathcal{L}_j \rightarrow \mathcal{L}_k, (x_i, x_j) \mapsto x_k$ where x_i, x_j, x_k are concurrent lines.

Let a given 3-web \mathcal{W} with pencils $\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3$ satisfy the closure condition derived from a 3-coloured 3-configuration \mathcal{C} with a simple basis \mathcal{B} . Then the corresponding term equality in \mathcal{C} over \mathcal{B} can be interpreted as an identity with respect to the coordinatizing 3-sorted quasigroup \mathbb{Q} and \mathcal{B} as the set of variables. Thus \mathcal{W} determines the corresponding identity over \mathbb{Q} .

We shall return to the 3-configuration \mathcal{C} on Fig. 4 and suppose that a 3-web \mathcal{W} satisfies a closure condition $\Gamma(\mathcal{C}, \mathcal{B})$. Then the coordinatizing 3-sorted quasigroup satisfies the identity $(x \cdot y) / y = (x \cdot y) / (x \setminus (x \cdot y))$ with variables x, x' of "position" 1 and variables y, y' of "position" 2.

If we take a homomorphism $\varrho : \hat{\mathcal{C}} \rightarrow \mathcal{W}$ with $\varrho(b) = \varrho(a') = \varrho(a'')$ we get a subconfiguration $\bar{\mathcal{C}}$ of \mathcal{C} on Fig. 5 with a simple basis $\bar{\mathcal{B}} = \{x, y, y'\}$. The closure condition $\Gamma(\bar{\mathcal{C}}, \bar{\mathcal{B}})$ determines the identity $(x \cdot y) / y = (x \cdot y) / y'$. Thus in \mathcal{W} from $\Gamma(\mathcal{C}, \mathcal{B})$ it follows that $\Gamma(\bar{\mathcal{C}}, \bar{\mathcal{B}})$.



$$0 \dots x, 0 \dots y, 0 \dots y', 2 \dots x \cdot y, 2 \dots x \cdot y', 4 \dots (x \cdot y) / y, 4 \dots (x \cdot y) / y'$$

Fig. 5

with $x' \cdot y' = x \cdot y$, $y' = y$. Thus $y' = x' \setminus (x \cdot y)$ and, after substitution into (ii'), we get $(x' \cdot y' / (x' \setminus (x \cdot y))) = (x' \cdot y) / y'$, i.e. $(x' \cdot y) / (x' \setminus (x \cdot y')) = (x \cdot y) / y$.

Writing here $x = x$, $x' = x$ we get (i).

REFERENCES

- [1] Belousov, V. D., *Configurations in algebraic nets*, Kishinev, 1979. (in Russian)
- [2] Krapež, A., Taylor, M. A., *Bases of web configurations*, Publications de l'Institut mathématique (nouvelle série) **38**(52), 21-30.
- [3] Brožíková, E., *On universal quasigroup identities*, Math. Bohem. **17** (1992), 20-32.

V. J. HAVEL
 DEPARTMENT OF MATHEMATICS
 TECHNICAL UNIVERSITY
 KRAVÍ HORA 21
 602 00 BRNO, CZECH REPUBLIC