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**COUNTABLE INDUCTIVE DEFINITIONS IN AST  
A. TZOUVARAS**

Abstract: We transfer the notions of inductive definition, fixed point and inductive class in Alternative Set Theory and show that every  $\Sigma$ -semiset is a fixed point and every  $\Sigma$ -class is inductive.

Key words: Inductive definition, fixed point, inductive class, Alternative Set Theory,  $\Sigma$ -class.

Classification: 02K10, 02B99

The main source of reference is [MO]. In Section 1 we adapt key definitions and cite basic facts from [MO]. In Section 2 we prove the results which seem to be specific for the context of AST. In Section 3 we show that all  $\Sigma$ -classes are inductive.

We assume the reader's familiarity with all basic concepts of the Alternative Set Theory as exposed e.g. in [V].

§ 1. Adapted definitions and facts. Let  $\varphi(Z)$  be a normal formula of  $FL_V$ ,  $Z$  being among the class variables of  $\varphi$ . We say that  $\varphi$  is positive in  $Z$ , or simply  $\varphi$  is positive (if  $Z$  is the only class variable of  $\varphi$ ), if  $\varphi$  belongs to the collection  $\Phi_p(Z)$  of formulas defined as follows:

$\Phi_p(Z)$  is the smallest class of formulas such that:

- (i) If  $Z$  does not occur in  $\varphi$ , then  $\varphi \in \Phi_p(Z)$ .
- (ii) If  $t$  is a constant or variable, then  $t \in Z$  is in  $\Phi_p(Z)$ .

(iii) If  $\varphi, \psi$  are in  $\Phi_p(Z)$  and  $x$  is any (set) variable, then  $\varphi \wedge \psi, \varphi \vee \psi, (\exists x)\varphi, (\forall x)\varphi$  are in  $\Phi_p(Z)$ .

The main property of positive formulas is monotonicity.

Lemma 1.1. If  $\varphi(Z)$  is positive in  $Z$ , then for any classes  $X, Y, X \subseteq Y$  and  $\varphi(X)$  imply  $\varphi(Y)$ .

Proof. By induction through the steps of the previous definition.

If the only class variable of  $\varphi(x, Z)$  is  $Z$ , the only set variable is  $x$  and  $\varphi$  is positive in  $Z$ , then we can adjoin to  $\varphi$  an operator  $\Gamma_\varphi$  sending the class  $X$  to the class

$$\Gamma_\varphi(X) = \{x; \varphi(x, X)\}.$$

$\Gamma_\varphi$  is monotone, i.e.

$$X \subseteq Y \rightarrow \Gamma_\varphi(X) \subseteq \Gamma_\varphi(Y).$$

( $\Gamma_\varphi$  is an informal object and we use it just to simplify the notation in some cases.) If  $\varphi$  contains only set-definable class parameters, then  $\Gamma_\varphi$  sends every set-definable class to a set-definable class.

Lemma 1.2. If  $\varphi(x, Z)$  and  $\psi(x, Z)$  are positive, then  $\varphi(x, \Gamma_\psi(Z))$  is positive in  $Z$ .

Proof. By induction on the length of  $\varphi$ .

Given a positive  $\varphi(x, Z)$  we define an increasing sequence of classes  $(I_\varphi^n)_{n \in \mathbb{N}}$  as follows:

$$I_\varphi^0 = \emptyset$$

$$I_\varphi^{n+1} = \Gamma_\varphi(I_\varphi^n) = \{x; \varphi(x, I_\varphi^n)\}.$$

This is a typical inductive definition which could probably be continued beyond the finite ordinals.

Here, however, we are interested in countable inductions, that is inductions which terminate in  $\omega$  steps. These are defined by positive formulas  $\varphi$  such that  $\Gamma_{\varphi}(\bigcup_n I_{\varphi}^n) = \bigcup_n I_{\varphi}^n$ .

Let us put

$$I_{\varphi} = \bigcup \{I_{\varphi}^n; n \in \mathbb{N}\}$$

for every positive  $\varphi$ .

We say that  $I_{\varphi}$  is a fixed point if  $\varphi$  contains no class parameters and  $\Gamma_{\varphi}(I_{\varphi}) = I_{\varphi}$ .

If  $\varphi(Z)$  is positive in  $Z$  and  $X$  is a class, we say that the parameter  $X$  is positive in  $\varphi(X)$ .

If all the parameters  $X_1, X_2, \dots, X_k$  of  $\varphi$  are positive and  $\Gamma_{\varphi}(I_{\varphi}) = I_{\varphi}$ , then we say that  $I_{\varphi}$  is a fixed point in  $X_1, X_2, \dots, X_k$ .

A class  $X$  is inductive (inductive in  $X_1, \dots, X_k$ ) if for some fixed point (fixed point in  $X_1, \dots, X_k$ )  $I_{\varphi}$  and some set parameter  $a$ ,

$$x \in X \leftrightarrow \langle x, a \rangle \in I_{\varphi} \leftrightarrow x \in I_{\varphi}^n \{a\}.$$

**Lemma 1.3.** (i) Every fixed point is inductive. (ii) Every inductive class is a  $\Sigma$ -class. (iii) Every set-definable class is a fixed point.

**Proof.** (i) Let  $I_{\varphi}$  be a fixed point in  $X_1, \dots, X_k$  and let  $a$  be a parameter. Put

$$\psi(x, y, Z) \equiv \varphi(x, Z \setminus \{a\}) \wedge y = a.$$

If  $\sigma(x, Z) \equiv \langle x, a \rangle \in Z$ , then clearly  $\sigma$  is positive and  $Z \setminus \{a\} = \Gamma_{\sigma}(Z)$ . By 1.2  $\psi$  is positive and we can see inductively that  $I_{\psi}^n = I_{\varphi}^n \times \{a\}$ , whence  $I_{\psi} = I_{\varphi} \times \{a\}$ ,  $I_{\psi}$  is a fixed point and  $I_{\varphi} = I_{\psi} \setminus \{a\}$ .

(ii) Let  $I_{\varphi}$  be a fixed point. If  $I_{\varphi}^n$  is set-definable then

obviously  $I_{\varphi}^{n+1} = \{x; \varphi(x, I^n)\}$  is set-definable. Similarly, for any  $a$ ,  $I_{\varphi}^n\{a\} = \bigcup_n I_{\varphi}^n\{a\}$  and  $I_{\varphi}^n\{a\}$  are set-definable. (This is no longer true, however, for classes inductive in other classes.)

(iii) If  $X$  is set-definable, put

$$\varphi(x, Z) \equiv x \in X.$$

Then  $I_{\varphi}^n = I_{\varphi} = X$  for  $n \geq 1$ .

Since we are interested in countable inductions, we have to deal exclusively with positive formulas leading to such inductions.

Let  $\varphi(Z)$  be positive and let  $(Y_n)_{n \in \mathbb{N}}$  be an increasing sequence of set-definable classes. We say that  $\varphi$  is stationary in  $Z$  w.r.t.  $(Y_n)_{n \in \mathbb{N}}$ , if

$$\varphi\left(\bigcup_n Y_n\right) \leftrightarrow (\exists n \in \mathbb{N}) \varphi(Y_n).$$

We say that  $\varphi$  is stationary in  $Z$  if it is stationary w.r.t. any such sequence.

Lemma 1.4. Let  $\varphi(Z)$  be positive and stationary. Then for every increasing sequence of inductive (or, more generally,  $\Sigma^1$ -) classes  $(D_n)_n$ ,

$$\varphi\left(\bigcup_n D_n\right) \leftrightarrow (\exists n \in \mathbb{N}) \varphi(D_n).$$

Proof. Suppose  $D_n, n \in \mathbb{N}$ , are  $\Sigma^1$ -classes and  $D_n = \bigcup_m D_n^m$ . Let  $(E_k)_{k \in \mathbb{N}}$  be an enumeration of all  $D_n^m, m, n \in \mathbb{N}$ . Define two functions  $H_1, H_2$  from  $\mathbb{N}$  to  $\mathbb{N}$  by recursion as follows:

$H_1(0) = H_2(0) = 0$ , and  $H_1(k+1) =$  least  $m$  such that there is an  $n$  such that

$$E_{k+1} \cup D_{H_1(k)}^{H_2(k)} \subseteq D_m^n,$$

$H_2(k+1) =$  least  $n$  such that

$$E_{k+1} \cup D_{H_1(k)}^{H_2(k)} \subseteq D_{H_1(k+1)}^n$$

The definition makes sense for all  $k \in \text{FN}$  because  $(D_n)_n$  is increasing. Clearly  $D_{H_1(n)}^{H_2(n)} \subseteq D_{H_1(n+1)}^{H_2(n+1)}$  and  $\bigcup_n D_n = \bigcup_n D_{H_1(n)}^{H_2(n)}$ .

Since  $\varphi$  is positive and stationary, we get

$$\begin{aligned} \varphi(\bigcup_n D_n) &\leftrightarrow \varphi(\bigcup_n D_{H_1(n)}^{H_2(n)}) \leftrightarrow (\exists n) \varphi(D_{H_1(n)}^{H_2(n)}) \rightarrow (\exists n) \varphi(D_{H_1(n)}) \rightarrow \\ &\rightarrow (\exists k) \varphi(D_k). \end{aligned}$$

The other direction follows from positivity.

It is evident that if  $\varphi(x, Z)$  is positive and stationary in  $Z$  and does not contain class parameters, then  $I_\varphi$  is a fixed point. A fixed point  $I_\varphi$  (or an inductive class  $I_\varphi''\{a\}$ ) is called stationary if the defining formula is stationary.

We shall see later that stationary formulas form a sufficiently large part of all positive formulas.

The following is a version of the Transitivity Theorem (cf. [MO], 1C.3).

Theorem 1.5. Let  $\varphi(x, y, Z_1, \dots, Z_k, Z)$  be a formula positive and stationary in all its class variables. If  $X_1, \dots, X_k$  are stationary inductive classes,  $\varphi_0(x, y, Z) \equiv \varphi(x, y, X_1, \dots, X_k, Z)$ ,  $I_{\varphi_0}$  is a fixed point in  $X_1, \dots, X_k$  and  $X = I_{\varphi_0}''\{a\}$  for some  $a$ , then  $X$  is stationary inductive.

Proof. To simplify the argument suppose  $k = 1$ . The treatment for  $k > 1$  is quite the same. Since  $X_1$  is stationary inductive, there is a positive stationary formula  $\varphi_1(x_1, y_1, Z)$  and a constant  $b_1$  such that  $X_1 = I_{\varphi_1}''\{b_1\}$ . We shall combine the two inductions defined by  $\varphi_1$  and  $\varphi$  into a single induction defined by a positive and stationary formula  $\varphi(x, y, x_1, y_1, t, Z)$ . Consider arbitrary constants  $x^*, y^*, x_1^*, y_1^*$ , and put:

$$\begin{aligned} \wp(x, y_1, y_1, t, Z) \equiv \\ [t = 1 \wedge \wp_1(x_1, y_1, \{ \langle x'_1, y'_1 \rangle; \langle x^*, y^*, x'_1, y'_1, 1 \rangle \in Z \})] \vee \\ [t = 2 \wedge \wp(x, y, \{ x'_1; \langle x, y, x'_1, b_1, 1 \rangle \in Z, \{ \langle x', y' \rangle; \langle x', y', x'_1, y'_1, 2 \rangle \in Z \})]. \end{aligned}$$

By 1.2  $\wp$  is positive and we claim that:

- i)  $\wp$  is stationary,
- ii)  $\langle x, y \rangle \in I_{\wp_0} \leftrightarrow \langle x, y, x'_1, y'_1, 2 \rangle \in I_{\wp}$ .

To see i) put

$$\begin{aligned} \{ \langle x'_1, y'_1 \rangle; \langle x^*, y^*, x'_1, y'_1, 1 \rangle \in Z \} &= Z^1, \\ \{ x'_1; \langle x, y, x'_1, b_1, 1 \rangle \in Z \} &= Z^2, \\ \{ \langle x', y' \rangle; \langle x', y', x'_1, y'_1, 2 \rangle \in Z \} &= Z^3. \end{aligned}$$

Then,

$$\wp(t, Z) \equiv [t = 1 \wedge \wp_1(Z^1)] \vee [t = 2 \wedge \wp(Z^2, Z^3)].$$

Let  $(Y_n)_{n \in \mathbb{N}}$  be an increasing sequence of set-definable classes. Then so are the sequences  $(Y_n^i)_{n \in \mathbb{N}}$ ,  $i = 1, 2, 3$ . Then,

$$\wp(t, \bigcup_n Y_n) \equiv [t = 1 \wedge \wp_1(\bigcup_n Y_n^1)] \vee [t = 2 \wedge \wp(\bigcup_n Y_n^2, \bigcup_n Y_n^3)],$$

and since  $\wp_1, \wp$  are stationary in all variables, we get

$$\wp(t, \bigcup_n Y_n) \leftrightarrow (\exists k, m, n \in \mathbb{N}) [t = 1 \wedge \wp_1(Y_k^1)] \vee [t = 2 \wedge \wp(Y_m^2, Y_n^3)],$$

whence, by monotonicity,

$$\wp(t, \bigcup_n Y_n) \leftrightarrow (\exists n \in \mathbb{N}) \wp(t, Y_n).$$

The essential part of the theorem is, of course, claim ii) and this is just the content of the Combination Lemma 1C.2 of [MO]. From ii) we have

$$x \in X \leftrightarrow \langle x, a \rangle \in I_{\wp_0} \leftrightarrow \langle x, a, x'_1, y'_1, 2 \rangle \in I_{\wp},$$

hence  $X$  is stationary inductive.

Remarks. 1) If  $\wp_0$  in the preceding theorem does not contain class variables, then the condition that  $I_{\wp_0}$  be a fixed point is obviously satisfied, hence for any formula  $\wp(x, X_1, \dots, X_k)$  with  $X_1, \dots, X_k$  as in the theorem, the class  $\{x; \wp(x, X_1, \dots, X_k)\}$  is stationary inductive.

2) It follows from the previous remark that, given  $\varphi(x, X_1, \dots, X_k, Z)$  with  $X_1, \dots, X_k$  as in the theorem,  $I_{\varphi}^1$  is stationary inductive and we see inductively that all  $I_{\varphi}^n$  are inductive. Thus, by 1.4  $\varphi(x, X_1, \dots, X_k, I_{\varphi}) \leftrightarrow x \in I_{\varphi}$ , that is,  $I_{\varphi}$  is a fixed point in  $X_1, \dots, X_k$ . This shows that the requirement for  $I_{\varphi_0}$  in the preceding theorem to be a fixed point in  $X_1, \dots, X_k$  is superfluous.

The reason that we use positive formulas instead of merely monotone ("up hereditary" in the terminology of Mlček, [M]) is that the former admit a canonical form. Namely the following holds:

**Theorem 1.6.** Let  $\varphi(Z)$  be a positive formula. Then there is a quantifier-free set-formula  $\Theta(\bar{z}, u) \equiv \Theta(z_1, \dots, z_n, u)$  and a string  $\bar{Q} \equiv Q_1 \dots Q_n$  of quantifiers such that

$$\varphi(Z) \equiv (Q_1 z_1) \dots (Q_n z_n) (\forall u) [\Theta(z_1, \dots, z_n, u) \vee u \in Z],$$

for all  $Z \neq V$ , or briefly,

$$(*) \quad \varphi(Z) \equiv (\bar{Q}\bar{z})(\forall u) (\Theta(\bar{z}, u) \vee u \in Z),$$

for all  $Z \neq V$ .

Proof. [MQ], 4B.1.

**Remark.** The restriction  $Z \neq V$  is for the case that  $\varphi(V)$  is false, since the right hand side of (\*) is always true for  $Z = V$ . However, given  $\varphi$ , we can put  $\psi(x, Z) \equiv \varphi(x, Z) \vee (\forall y)(y \in Z)$ .

If  $\varphi$  is positive and stationary then so is  $\psi$ ,  $I_{\varphi} = I_{\psi}$  and, in addition,  $(\forall x)\psi(x, V)$  is true. Thus studying fixed points we may always assume that (\*) holds for all  $Z$ .



§ 2. Some results on fixed points. In the proof of the next lemma we just use the fact that for every set-definable class X and for every  $\Sigma$ -class  $\bigcup_n Y_n$ ,

$$X \subseteq \bigcup_n Y_n \leftrightarrow (\exists n)(X \subseteq Y_n).$$

Lemma 2.1. Let  $\varphi(x, Z)$  be positive. If  $\varphi$  has a canonical form

$$(\bar{Q}\bar{z})(\forall u)(\Theta(x, \bar{z}, u) \vee u \in Z),$$

where  $\bar{Q} = \emptyset$ , or  $\forall^k$ , or  $\exists^k$ , or  $\exists^k \forall^k$ , then  $\varphi$  is stationary.

Proof. Put  $R = \{ \langle x, \bar{z}, u \rangle; \neg \Theta(x, \bar{z}, u) \}$ , and let  $R''(x, \bar{z}) = \{ u; \langle x, \bar{z}, u \rangle \in R \}$ . Then

$$(1) \quad \varphi(x, Z) \equiv (\bar{Q}\bar{z})(R''(x, \bar{z}) \subseteq Z).$$

Let us show for example the case  $\bar{Q} = \forall^k$ . The rest is shown similarly. Let  $(Y_n)_n$  be an increasing sequence of set definable classes. Then:

$$\begin{aligned} \varphi(x, \bigcup_n Y_n) &\leftrightarrow (\forall z_1 \dots z_k)(R''(x, z_1, \dots, z_k) \subseteq \bigcup_n Y_n) \leftrightarrow \\ &\cup \{ R''(x, z_1, \dots, z_k); z_1, \dots, z_k \in V \} \subseteq \bigcup_n Y_n \leftrightarrow \\ &(\exists n) [ \cup \{ R''(x, z_1, \dots, z_k); z_1, \dots, z_k \in V \} \subseteq Y_n ] \leftrightarrow \\ &(\exists n)(\forall z_1 \dots z_k)(R''(x, z_1, \dots, z_k) \subseteq Y_n) \leftrightarrow (\exists n)\varphi(x, Y_n). \end{aligned}$$

Let us call an existential quantifier in the prefix  $(\bar{Q}\bar{z})$  of the canonical form  $(*)$  inessential if there is some set-definable Skolem-function for it. The following is obvious:

Lemma 2.2. If the prefix  $(\bar{Q}\bar{z})$  of the canonical form of  $\varphi$  after the extraction of all inessential quantifiers is as in Lemma 2.1, then  $\varphi$  is stationary.

The last two lemmas imply that the simplest positive non-stationary formula cannot be less complicated than the formula  $(\forall z_1)(\exists z_2)(R''(x, z_1, z_2) \subseteq Z)$ , where  $(\exists z_2)$  is not inessential.

However we do not know whether there exist non-stationary

positive formulas.

The next result restricts further the possible  $\varphi$  for which there is not a fixed point.

Lemma 2.3. If  $(u_n)_{n \in \mathbb{N}}$  is an increasing sequence of sets, then every positive formula is stationary w.r.t.  $(u_n)_{n \in \mathbb{N}}$ .

Proof. Let  $\varphi(x, Z)$  be positive. Taking a canonical form of  $\varphi$  and defining  $R$  as in 2.1, we have

$$\varphi(x, Z) \equiv (\bar{Q}\bar{z})(R(x, \bar{z}) \subseteq Z).$$

Then,  $\varphi(x, \bigcup_n u_n) \equiv (\bar{Q}\bar{z})(R(x, \bar{z}) \subseteq \bigcup_n u_n) \leftrightarrow (\bar{Q}\bar{z})(\exists n)(R(x, \bar{z}) \subseteq u_n)$ .

Put  $R(x, \bar{z}) \subseteq u_n \equiv \psi(x, \bar{z}, u_n)$ .

$\psi$  is positive in  $u_n$  and it suffices to show that for every set-formula  $\psi(x, \bar{z}, w)$ , positive in  $w$  and every increasing sequence  $(u_n)_{n \in \mathbb{N}}$ ,

$$(\bar{Q}\bar{z})(\exists n) \psi(x, \bar{z}, u_n) \leftrightarrow (\exists n) \bar{Q}z \psi(x, \bar{z}, u_n).$$

Suppose  $\psi$  and  $(u_n)_{n \in \mathbb{N}}$  are given. Prolong the sequence  $(u_n)_{n \in \mathbb{N}}$  to a set  $\{u_\beta; \beta \leq \alpha\}$  such that  $u_\beta \subseteq u_\gamma$  for  $\beta \leq \gamma$ . If the string  $\bar{Q}$  has length  $m$ , the above equivalence is shown in  $m$  steps by pulling at the step  $\alpha$  the quantifier  $(\exists n)$  to the front of the quantifier  $Q_{m-k+1}$ . If  $Q_{m-k+1} = \exists$  this is trivially possible. Thus it suffices to show that

$$(+)$$

$$(\forall z)(\exists n) \delta(x, z, u_n) \leftrightarrow (\exists n)(\forall z) \delta(x, z, u_n)$$

for  $\delta$  positive in  $u_n$ .

Let  $F: V \rightarrow N$  be a function such that

$$F(z) = \beta \leftrightarrow \delta(x, z, u_\beta) \wedge (\forall \gamma < \beta) \neg \delta(x, z, u_\gamma).$$

The direction " $\leftarrow$ " of (+) is obvious.

Now, if the left hand-side of (+) is true, then  $F$  is (set-) defined on  $V$  and  $F''V \subseteq \mathbb{N}$ . Therefore  $F''V$  is finite. It follows that  $(\forall z)(\exists n) \delta(x, z, u_n) \leftrightarrow (\exists n)(\forall z) [\delta(x, z, u_1) \vee \dots \vee \delta(x, z, u_n)]$ .

Since  $\sigma$  is positive and  $(u_n)_{n \in \mathbb{N}}$  is increasing, we get

$$\sigma(x, z, u_1) \vee \dots \vee \sigma(x, z, u_n) \leftrightarrow \sigma(x, z, u_n)$$

and the proof is complete.

A positive formula  $\varphi(x, Z)$  will be called reversible, if there is some positive  $\psi(x, Z)$  such that  $\Gamma_\psi(\Gamma_\varphi(Z)) = Z$  for every class  $Z$ .  $\psi$  is called a reverse of  $\varphi$ .

For example the formula

$$\psi(x, Z) \equiv (\exists y)(x \in y \wedge y \in Z),$$

with operator  $\Gamma_\psi = \cup$  is a reverse of the formula

$$\varphi(x, Z) \equiv x \in Z,$$

with operator  $P$  (the power-set operator), since  $\cup P(Z) = Z$ .

Lemma 2.4. If  $I_\varphi$  is a fixed point and  $\varphi(x, Z)$  is a stationary in  $(I_\varphi^n)_{n \in \mathbb{N}}$  and reversible positive formula, then  $\Gamma_\psi(I_\varphi)$  is a fixed point.

Proof. Let  $\sigma(x, Z)$  be a reverse of  $\psi$  and put

$$\rho(x, Z) \equiv \psi(x, \Gamma_\sigma^{-1}(Z)).$$

By 1.2  $\rho$  is positive. We show that  $I_\rho^n = \Gamma_\psi(I_\varphi^n)$  for  $n \geq 1$ .

First notice that  $\Gamma_\sigma(\emptyset) = \emptyset$  because  $\Gamma_\sigma^{-1}(\Gamma_\psi(\emptyset)) = \emptyset$  and  $\emptyset \subseteq \Gamma_\psi(\emptyset)$  and  $\Gamma_\sigma^{-1}(\emptyset) \subseteq \Gamma_\sigma^{-1}(\Gamma_\psi(\emptyset))$ . Hence  $I_\rho^1 = \{x; \psi(x, \Gamma_\sigma^{-1}(\emptyset))\} = \{x; \psi(x, \Gamma_\varphi^{-1}(\emptyset))\} = \Gamma_\psi(I_\varphi^1)$ .

Suppose  $I_\rho^n = \Gamma_\psi(I_\varphi^n)$ . Then

$$\begin{aligned} I_\rho^{n+1} &= \Gamma_\rho(I_\rho^n) = \Gamma_\psi \Gamma_\sigma^{-1} \Gamma_\sigma^{-1}(\Gamma_\psi(I_\varphi^n)) = \Gamma_\psi \Gamma_\sigma^{-1}(\Gamma_\psi(I_\varphi^n)) = \\ &= \Gamma_\psi(I_\varphi^{n+1}). \end{aligned}$$

Since  $\psi$  is stationary, we have

$$I_\rho = \bigcup_n I_\rho^n = \bigcup_n \Gamma_\psi(I_\varphi^n) = \Gamma_\psi(\bigcup_n I_\varphi^n) = \Gamma_\psi(I_\varphi).$$

It remains to show that  $\Gamma_\rho(I_\rho) = I_\rho$ .

$$\Gamma_\rho(I_\rho) = \Gamma_\psi \Gamma_\sigma^{-1} \Gamma_\sigma^{-1}(I_\rho) = \Gamma_\psi \Gamma_\sigma^{-1} \Gamma_\sigma^{-1}(\Gamma_\psi(I_\varphi)) = \Gamma_\psi \Gamma_\sigma^{-1}(I_\varphi) = \Gamma_\psi(I_\varphi) = I_\rho$$

(If moreover  $\varphi, \sigma$  are stationary, then  $\rho$  is stationary.)

Corollary 2.5. If  $I_\varphi$  is a fixed point and  $F$  is a 1-1 set-definable function with  $I_\varphi \subseteq \text{dom}(F)$ , then  $F''I_\varphi$  is a fixed point. If  $I_\varphi$  is stationary,  $F''I_\varphi$  is stationary.

Proof. Consider the formula

$\psi(x, Z) \equiv (\exists y \in Z)(F(y) = x)$ . Clearly  $\Gamma'_\psi(Z) = F''Z$  for  $Z \subseteq \text{dom}(F)$ ,  $\psi$  is stationary and the formula

$\delta(x, Z) \equiv (\exists y \in Z)(F(x) = y)$  which is stationary, is a reverse for  $\psi$ . The conclusion follows from 2.4.

Corollary 2.6 Every countable class is a stationary fixed point.

Proof. Since for any countable class  $X$  there is a 1-1 set-function  $f$  such that  $f''FN = X$ , it suffices by 2.5 to prove that  $FN$  is a stationary fixed point.

Let  $<$  be the ordering of natural numbers and put

$$\varphi(x, Z) \equiv (\forall u)(u < x \rightarrow u \in Z).$$

Clearly  $\varphi$  is positive, stationary and  $I_\varphi^n = n$ .

We shall now prove that every countable union of sets ( $\Sigma$ -semisets) or cosets is a fixed point. First a lemma.

Lemma 2.7. Let  $(u_n)_{n \in FN}$  be a sequence of sets. Then there is an increasing sequence  $(v_n)_{n \in FN}$  such that  $\bigcup_n u_n = \bigcup_n v_n$  and the sequence of natural numbers  $\{v_{n+1} - v_n\}$  is either increasing or decreasing ( $\{u\}$  is the unique  $\alpha \in N$  such that  $u \hat{\approx} \alpha$ ).

Proof. Suppose, without loss of generality, that  $(u_n)_{n \in FN}$  is increasing. Then either

$$(1) (\exists n)(\forall m > n)(\exists k > m)(|u_m - u_n| \neq |u_k - u_m|),$$

or the negation of (1) is true:

$$(2) (\forall n)(\exists m > n)(\forall k > m)(|u_k - u_m| < |u_m - u_n|).$$

In the first case we can find a subsequence  $(u_{n_k})_k = (v_k)_k$ , such that  $|v_{k+1} - v_k|$  is increasing and in the second case we find a subsequence  $(v_k)_k$  such that  $|v_{k+1} - v_k|$  is decreasing.

**Theorem 2.8.** For every sequence  $(u_n)_{n \in \mathbb{N}}$ ,  $\bigcup_n u_n$  is a stationary fixed point.

**Proof. Case 1.** Suppose there is an increasing sequence  $(v_n)_n$  with  $|v_{n+1} - v_n|$  increasing such that  $\bigcup_n u_n = \bigcup_n v_n$ . Extend  $(v_n)_n$  to a set  $\{v_\beta; \beta \leq \alpha\}$  with  $|v_{\beta+1} - v_\beta|$  increasing.

For every  $\beta$ ,  $2 \leq \beta \leq \alpha$ , let  $g_\beta$  be the surjection from  $v_\beta - v_{\beta-1}$  onto  $v_{\beta-1} - v_{\beta-2}$  which is least in the usual set-definable ordering of  $V$ . Let also  $g_1$  be the least surjection from  $v_1$  onto  $v_0$ . The correspondence  $\beta \mapsto g_\beta$  is set-definable and put  $f_\beta = \bigcup \{g_\sigma; \sigma \leq \beta\}$ .

Then  $f_\beta$  is a function from  $v_\beta$  onto  $v_{\beta-1}$ . Let  $f = f_\alpha$ . It is easily seen that  $f \upharpoonright v_\beta = f_\beta$  and  $f^{-1} \upharpoonright v_\beta = v_{\beta+1}$ . (Some Venn-diagrams illustrate best the situation.)

If we put

$$\varphi(x, Z) \equiv x \in v_0 \vee (\exists y \in Z)(f(x) = y)$$

then  $\varphi$  is stationary and  $I_\varphi^{n+1} = v_n$ , hence  $\bigcup_n u_n = \bigcup_n v_n = I_\varphi$ .

**Case 2.** Let again  $\bigcup_n u_n = \bigcup_n v_n$  where  $|v_{n+1} - v_n|$  is decreasing. Extend as before  $(v_n)_n$  to a set  $\{v_\beta; \beta \leq \alpha\}$ . Let  $g_\beta: v_{\beta-1} - v_{\beta-2} \rightarrow v_\beta - v_{\beta-1}$  be least surjections for  $2 \leq \beta \leq \alpha$ , while  $g_1$  is the identity on  $v_0$ . Put  $f_\beta = \bigcup \{g_\sigma; \sigma \leq \beta\}$ ; then  $f_\beta$  maps  $v_{\beta-1}$  onto  $v_\beta - (v_1 - v_0)$ . Put  $f = f_\alpha$ . Then  $f \upharpoonright v_\beta = f_{\beta+1}$  for  $\beta < \alpha$ . Consider the formula

$$\varphi(x, Z) \equiv x \in v_1 \vee (\exists y \in Z)(f(y) = x).$$

Again  $\varphi$  is stationary and  $I_\varphi^n = v_n$  for  $n \geq 1$ .

A completely analogous result holds for sequences of cosets

$(V - u_n)_{n \in \mathbb{N}}$  or, more generally,  $(X - u_n)_{n \in \mathbb{N}}$ , where  $X$  is set-definable. It is evident that Lemma 2.7 is equally true if we substitute "decreasing" for "increasing" and  $\cap$  for  $\cup$ . Then, given a set-definable  $X$  and a sequence  $(u_n)_n$ , such that

$$\dots \subseteq u_1 \subseteq u_0 \subseteq X,$$

with  $|u_n - u_{n+1}|$  either increasing or decreasing, it is easy to construct a set-definable function  $f$  with  $\text{dom}(f) \subseteq X$  and such that either  $f$  or  $f^{-1}$  defines  $X - u_{n+1}$  by means of  $X - u_n$ . Therefore:

Lemma 2.9. If  $X$  is set-definable and  $(u_n)_n$  is any sequence of sets, then  $\bigcup_n (X - u_n)$  is a stationary fixed point.

§ 3. All  $\Sigma$ -classes are inductive. Let  $\text{Fix}$ ,  $\text{Ind}$ ,  $\Sigma$  denote respectively the (codable) classes of fixed points, inductive classes and  $\Sigma$ -classes. By 1.3

$$\text{Fix} \subseteq \text{Ind} \subseteq \Sigma.$$

We shall prove in this section that

Theorem 3.1.  $\Sigma = \text{Ind}$ .

This will be done through a number of lemmas.

Lemma 3.2. If  $\text{Sd}_V$  has a code  $\langle K, S \rangle$  such that the class  $S$  is stationary inductive, then  $\Sigma = \text{Ind}$ .

Proof. Let  $\langle K, S \rangle$  be a code of  $\text{Sd}_V$  such that  $S$  is stationary inductive and let  $X = \bigcup_n X_n$  be a  $\Sigma$ -class. Then evidently there is a countable  $Y \subseteq K$  such that

$$\{X_n; n \in \mathbb{N}\} = \{S^{\ulcorner y \urcorner}; y \in Y\},$$

hence we get

$$x \in X \leftrightarrow (\exists y \in Y)(x \in S^{\ulcorner y \urcorner}).$$

The formula  $(\exists y \in Y)(x \in S''\{y\})$  contains the class-parameters  $Y, S$  in positive stationary positions and  $Y$  (by 2.6) as well as  $S$  (by assumption) are stationary inductive. It follows from the Transitivity Theorem 1.5 that  $X$  is stationary. Thus  $\Sigma \subseteq \text{ind}$ .

Take a Gödelization of the language  $\text{FSL}_V$ , i.e. of all the finite set-formulas, as a mapping  $G: \text{FSL}_V \rightarrow V$  defined as follows:

- 1)  $G(x_n) = \langle 0, n \rangle$ , for the set variables  $x_n, n \in \text{FN}$ .
- 2)  $G(x) = \langle 1, x \rangle$ , for the set-constants  $x \in V$ .
- 3)  $G(t=s) = \langle 2, G(t), G(s) \rangle$ , for constants or variables  $t, s$ .
- 4)  $G(t \in s) = \langle 3, G(t), G(s) \rangle$ , " " " " " "
- 5)  $G(\neg \varphi) = \langle 4, G(\varphi) \rangle$ .
- 6)  $G(\varphi \wedge \psi) = \langle 5, G(\varphi), G(\psi) \rangle$ .
- 7)  $G(\exists x_n \varphi) = \langle 6, n, G(\varphi) \rangle$ .

We say that  $G(\varphi)$  is the Gödel-set of  $\varphi$ .

Lemma 3.3. The class  $\text{Fml} = \{x; x \text{ is the Gödel-set of a set formula}\}$  is stationary inductive.

Proof. Let us denote by  $\text{AFml}$  the class of (Gödel-sets of) atomic formulas. Then, from the definition of  $G$  we have

$$\begin{aligned} x \in \text{AFml} &\leftrightarrow (\exists m, n \in \text{FN})(\exists y, z)[x = \langle 2, \langle 0, m \rangle, \langle 0, n \rangle \rangle \vee \\ &x = \langle 3, \langle 0, m \rangle, \langle 0, n \rangle \rangle \vee x = \langle 2, \langle 0, m \rangle, \langle 1, y \rangle \rangle \vee x = \langle 3, \langle 0, m \rangle, \langle 1, y \rangle \rangle \vee \\ &x = \langle 2, \langle 1, y \rangle, \langle 1, x \rangle \rangle \vee x = \langle 3, \langle 1, y \rangle, \langle 1, x \rangle \rangle]. \end{aligned}$$

It is clear that the defining formula is positive and stationary in  $\text{FN}$  and  $\text{FN}$  is stationary inductive, thus  $\text{AFml}$  is stationary inductive by 1.5.

Next,

$$\begin{aligned} x \in \text{Fml} &\leftrightarrow (\exists f)(\exists n \in \text{FN})(\text{dom}(f) = n+1 \wedge f(n) = x \wedge \\ &(\forall k < n)[f(k) \in \text{AFml} \vee (\exists l, m < k)(f(k) = \langle 4, f(l) \rangle \vee f(k) = \\ &= \langle 5, f(l), f(m) \rangle) \vee (\exists i \in \text{FN})(\exists j < k)(f(k) = \langle 6, i, f(j) \rangle)]]. \end{aligned}$$

Again the defining formula is positive and stationary in

FN, AFml while the latter are stationary inductive, hence Fml is stationary inductive.

Consider now the predicate of satisfaction  $Sat(x,g)$  expressing the fact: "x is the Gödel-set of a set-formula  $\varphi$ , g is a sequence of valuations for the variables  $x_n, n \in FN$ , and  $\varphi$  is true substituting  $g(i)$  for its free variable  $x_i$ ".

We fix a number  $\alpha_0 \in N-FN$  and let

$$A = \{g; \text{dom}(g) = \alpha_0 \wedge (\exists n \in FN)(\forall k > n)(g(k) = 0)\}.$$

Clearly A is stationary inductive. For  $g \in A$  write  $g \models \varphi$  for the fact that  $\varphi$  is true w.r.t. the valuation g.

The following is a version of 5.3.2 of [MO].

Lemma 3.4. The class  $Sat = \{\langle x, g \rangle; Sat(x, g)\}$  is stationary inductive.

Proof. Define the predicate  $Val(x, g, t)$  by:

$$Val(x, g, t) \equiv (x \text{ is the Gödel-set of a formula } \varphi) \wedge g \in A \wedge \wedge [(t = 0 \wedge g \models \varphi) \vee (t = 1 \wedge g \models \neg \varphi)].$$

Then obviously

$$Sat(x, g) \leftrightarrow Val(x, g, 0),$$

and it suffices to prove that there is a formula  $\varphi$  positive and stationary in all its class parameters, the latter being stationary inductive, such that  $I_\varphi$  is a fixed point and

$$Val(x, g, t) \leftrightarrow \langle x, g, t \rangle \in I_\varphi.$$

Let  $AVal(x, g, t) \equiv x \in AFml \wedge Val(x, g, t)$ .

Then:

$$AVal(x, g, t) \leftrightarrow g \in A \wedge (\exists m, n \in FN)(\exists y, z)$$

$$\{x = \langle 2, \langle 0, m \rangle, \langle 0, n \rangle \rangle \wedge [(t = 0 \wedge g(m) = g(n)) \vee (t = 1 \wedge g(m) \neq g(n))]\} \vee$$

$$\{x = \langle 3, \langle 0, m \rangle, \langle 0, n \rangle \rangle \wedge [(t = 0 \wedge g(m) \in g(n)) \vee (t = 1 \wedge g(m) \notin g(n))]\} \vee$$

$$\{x = \langle 2, \langle 0, m \rangle, \langle 1, y \rangle \rangle \wedge [(t = 0 \wedge g(m) = y) \vee (t = 1 \wedge g(m) \neq y)]\} \vee$$



$$\{x = \langle 3, \langle 0, m \rangle, \langle 1, y \rangle \rangle \wedge [(t=0 \wedge g(m) \in y) \vee (t=1 \wedge g(m) \notin y)]\} \vee$$

$$\{x = \langle 2, \langle 1, y \rangle, \langle 1, z \rangle \rangle \wedge [(t=0 \wedge y = z) \vee (t=1 \wedge y \neq z)]\} \vee$$

$$\{x = \langle 3, \langle 1, y \rangle, \langle 1, z \rangle \rangle \wedge [(t=0 \wedge y \in z) \vee (t=1 \wedge y \notin z)]\}$$

Then the class  $AVal = \{ \langle x, g, t \rangle; AVal(x, g, t) \}$  is stationary inductive since  $AVal(x, g, t)$  is positive stationary in  $A$ ,  $FN$ .

Define the function  $F$  on  $V^{\alpha_0} \times \alpha_0 \times V$  as follows:

$$F(g, \beta, u) = (g - \{ \langle \beta, g(\beta) \rangle \}) \cup \{ \langle \beta, u \rangle \}$$

i.e.,  $F(g, \beta, u)$  is the function resulting from  $g$  if we replace its value at  $\beta$  by  $u$ . Now define the required formula  $\wp$  as follows:

$$\wp(x, g, t, Z) \equiv x \in Fml \wedge g \in A \wedge \{ \langle x, g, t \rangle \in AVal \vee$$

$$(\exists y) [ (x = \langle 4, y \rangle \wedge t = 0 \wedge \langle y, g, 1 \rangle \in Z) \vee$$

$$(x = \langle 4, y \rangle \wedge t = 1 \wedge \langle y, g, 0 \rangle \in Z) ] \vee$$

$$(\exists y, z) [ (x = \langle 5, y, z \rangle \wedge t = 0 \wedge \langle y, g, 0 \rangle \in Z \wedge \langle z, g, 0 \rangle \in Z) \vee$$

$$(x = \langle 5, y, z \rangle \wedge t = 1 \wedge \langle y, g, 1 \rangle \in Z \wedge \langle z, g, 1 \rangle \in Z)$$

$$(\exists k \in FN) (\exists y) [ (x = \langle 6, k, y \rangle \wedge t = 0 \wedge (\exists z) \langle y, F(g, k, z), 0 \rangle \in Z) \vee$$

$$(x = \langle 6, k, y \rangle \wedge t = 1 \wedge (\forall z) \langle y, F(g, k, z), 1 \rangle \in Z) ] \}$$

We can summarize  $\wp$  as follows:

$$\wp(x, g, t, Z) \equiv x \in Fml \wedge g \in A \wedge [ \langle x, g, t \rangle \in AVal \vee$$

$$(\exists k \in FN) (\exists \bar{y}) (\forall \bar{z}) \sigma(t, x, g, \bar{y}, \bar{z}, Z) ],$$

where  $\sigma$  is positive in  $Z$  and contains only inessential existential quantifiers. It follows from 2.1 and 2.2 that  $\wp$  is stationary in  $Z$ . That  $\wp$  is positive stationary in  $Fml$ ,  $A$ ,  $AVal$ ,  $FN$  is evident. Also all these class-parameters are stationary inductive as we proved earlier. We must also prove that induction by  $\wp$  closes in  $\omega$  steps but this is clear from the remarks following Th. 1.5.

It remains to see that

$$Val(x, g, t) \leftrightarrow \langle x, g, t \rangle \in I_{\wp}^n.$$

Direction " $\rightarrow$ " is shown by induction on the length of the formula  $x$ , while by induction on  $n$  we show that

$$\langle x, g, t \rangle \in I_{\wp}^n \rightarrow Val(x, g, t).$$

Proof of Theorem 3.1. The pair  $\langle \text{Fml}, \text{Sat} \rangle$  is a code for  $\text{Sd}_V$  since for every  $X \in \text{Sd}_V$  such that  $X = \{x; \varphi(x)\}$ , we have  $X = \{g; \text{Sat}(x, g)\} = \text{Sat}^{\ulcorner x \urcorner}$ , where  $x$  is the Gödel-set of  $\varphi$ . From 3.4 it follows that the code is inductive, and by 3.2  $\Sigma = \text{Ind}$ .

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