

Aplikace matematiky

Zdeněk Režný; Evžen Kindler

Algorithms. 31. PERMUT. Simple algorithm generating all permutations

Aplikace matematiky, Vol. 18 (1973), No. 3, 211–212

Persistent URL: <http://dml.cz/dmlcz/103471>

Terms of use:

© Institute of Mathematics AS CR, 1973

Institute of Mathematics of the Czech Academy of Sciences provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these *Terms of use*.



This document has been digitized, optimized for electronic delivery and stamped with digital signature within the project *DML-CZ: The Czech Digital Mathematics Library* <http://dml.cz>

ALGORITMY

31. PERMUT

SIMPLE ALGORITHM GENERATING ALL PERMUTATIONS

Ing. ZDENĚK REŽNÝ, CSc., Dr. EVŽEN KINDLER, CSc.,
 Biofyzikální ústav Fakulty všeobecného lékařství Karlovy university,
 Praha 2, Salmovská 3

The procedure *PERMUT* generates all permutations of the sequence $a_m, a_{m+1}, \dots, a_{m+n-1}$. In every step of its work, it generates a permutation, assigned for the variables $a[m], \dots, a[m+n-1]$. The first permutation is the original one, entering into the algorithm; after generating all the permutations, the procedure *PERMUT* restitutes the original order of values in the array a . Any permutation can be processed by the procedure *P* immediately after it has been generated. Besides the permutations, the number of their inversions is generated and assigned for the variable r^*). If we do not need it, we can simplify the algorithm by omitting all statements where the variable r occurs.

```

procedure PERMUT( $a, m, n, p$ );
value  $m, n$ ;
array  $a$ ; integer  $m, n$ ; procedure  $P$ ;
begin
    integer  $i, k, r$ ; integer array  $j[1 : n - 1]$ ; real  $b$ ;
    for  $i := 1$  step 1 until  $n - 1$  do  $j[i] := i + m; r := 0$ ;
     $G : P; i := n - 1$ ;
     $H : \text{if } i \leq 0 \text{ then go to } D$ ;
    if  $j[i] = m$  then go to  $F$ ;
     $b := a[j[i]]$ ;  $a[j[i]] := a[j[i] - 1]$ ;  $j[i] := j[i] - 1$ ;  $a[j[i]] := b$ ;
     $r := r + 1$ ; go to  $G$ ;
     $F : j[i] := j[i] + i$ ;  $b := a[m]$ ;
    for  $k := m$  step 1 until  $m + i - 1$  do  $a[k] := a[k + 1]$ ;
     $a[m + i] := b$ ;  $r := r - i$ ;  $i := i - 1$ ; go to  $H$ ;
     $D :$ 
end PERMUT;
    
```

*) The number of inversions in a permutation $a_{j_1}, a_{j_2}, \dots, a_{j_n}$ is defined commonly as the number of all pairs of indices (i, k) , $1 \leq i, k \leq n$, where $i < k, j_i > j_k$.

The array a may be of course of the type integer, in which case the variable b is to be declared integer too.

Example. Calling $PERMUT(a, 1, 4, p)$ where $a_i = i$ and P prints the generated permutation, followed by the number of its inversions in parentheses:

1234(0), 1243(1), 1423(2), 4123(3), 1324(1), 1342(2), 1432(3), 4132(4),
3124(2), 3142(3), 3412(4), 4312(5), 2134(1), 2143(2), 2413(3), 4213(4),
2314(2), 2341(3), 2431(4), 4231(5), 3214(3), 3241(4), 3421(5), 4321(6).

An interesting property of the algorithm may be seen which holds not only for the given example, but generally: the sequence of generated permutations has certain symmetrical attributes, namely that the sum of the pair of the numbers of the inversions of the i -th and $(n! - i + 1)$ -th permutations is equal to $\binom{n}{2}$, and that the order of the i -th permutation is exactly inverse to the order of the $(n! - i + 1)$ -th one ($i = 1, \dots, n!$).

The algorithm has been programmed and tested in the small computer ODRA 1013 [1] in the programming language MOST [2].

References

- [1] Černý, V., Pár J.: Programmer's Manual on Automatic Computer ODRA 1013 (in Czech). Kancelářské stroje n. p., Hradec Králové 1967.
- [2] Szczepkiewicz, J.: Programming in the autocode MOST 1 (in Polish). ELWRO Publication 03—VI—1, Wrocław.